



Artificial Intelligence in Theorem Proving

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VTSA 2019

- Computer used to automate reasoning in a logic
- Traditionally part of artificial intelligence
 - (not machine learning)
- Field of research since the fifties
- Applications: program verification, mathematical deduction, ...
- Theorem proving logics, precision, automation, ... very varied.

- 1940s: Algorithmic proof search (λ-calculus)
- 1960s: de Bruijn's Automath
- 1970s: Small Certifiers (LCF)
- 1990s: Resolution (Superposition)
- 2000s: Large proofs and theories
- 2010s: Machine Learning for Reasoning?

Part I

Theorem proving systems

Machine learning problems

Lemma relevance

Deep learning for theorem proving

Part II

Guided Automated Reasoning Lemma mining Unsupervised methods Longer proofs

A Proof Assistant is a

- a computer program
- to assist a mathematician
- in the production of a proof
- that is mechanically checked

What does a Proof Assistant do?

- Keep track of theories, definitions, assumptions
- Interaction proof editing
- Proof checking
- Automation proof search

What does it implement? (And how?)

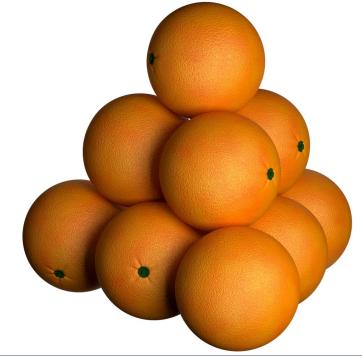
- a formal logical system intended as foundation for mathematics
- decision procedures

The Kepler Conjecture (year 1611)

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The most compact way of stacking balls of the same size in space is a pyramid.

$$V = \frac{\pi}{\sqrt{18}} \approx 74\%$$



Proved in 1998

- Tom Hales, 300 page proof using computer programs
- Submitted to the Annals of Mathematics

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- Tom Hales, 300 page proof using computer programs
- Submitted to the Annals of Mathematics
- 99% correct... but we cannot verify the programs

1039 equalities and inequalities

For example:

$$\frac{\frac{-x_{1}x_{3}-x_{2}x_{4}+x_{1}x_{5}+x_{3}x_{6}-x_{5}x_{6}+}{+x_{2}(-x_{2}+x_{1}+x_{3}-x_{4}+x_{5}+x_{6})}}{\sqrt{4x_{2}\binom{x_{2}x_{4}(-x_{2}+x_{1}+x_{3}-x_{4}+x_{5}+x_{6})+}{+x_{1}x_{5}(x_{2}-x_{1}+x_{3}+x_{4}-x_{5}+x_{6})+}{-x_{1}x_{3}(x_{2}-x_{1}-x_{3}+x_{4}+x_{5}-x_{6})-}{-x_{1}x_{3}x_{4}-x_{2}x_{3}x_{5}-x_{2}x_{1}x_{6}-x_{4}x_{5}x_{6})}} < \tan(\frac{\pi}{2} - 0.74)$$

Solution? Formalized Proof!

- Formalize the proof using Proof Assistants
- Implement the computer code in the system
- Prove the code correct
- Run the programs inside the Proof Assistant

Flyspeck Project

- Project results published 2017
- Many Proof Assistants and contributors

Superscalar; Dual integer pipeline; Faster floating-point, ...

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 $\frac{4159835}{3145727} = 1.333820... \qquad \frac{4159835}{3145727} \stackrel{\rm P5}{=} 1.333739...$

Superscalar; Dual integer pipeline; Faster floating-point, ...

 $\frac{4159835}{3145727} = 1.333820... \qquad \frac{4159835}{3145727} \stackrel{\rm P5}{=} 1.333739...$

FPU division lookup table: for certain inputs division result off

Replacement

- Few customers cared, still cost of \$475 million
- Testing and model checking insufficient:
 - Since then Intel and AMD processors formally verified (*)
 - HOL Light and ACL2 (along other techniques)

```
theorem sqrt2_not_rational:
    "sqrt (real 2) ∉ ℚ"
proof
```

qed

```
theorem sqrt2_not_rational:
    "sqrt (real 2) ∉ ℚ"
proof
    assume "sqrt (real 2) ∈ ℚ"
```

thus False qed

```
theorem sqrt2_not_rational:
    "sqrt (real 2) ∉ Q"
proof
    assume "sqrt (real 2) ∈ Q"
    then obtain m n :: nat where
    n_nonzero: "n ≠ 0" and sqrt_rat: "¦sqrt (real 2)¦ = real m / real
    and lowest terms: "gcd m n = 1" ...
```

thus False qed

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```

```
have eq: m^2 = 2 * n^2"
hence "2 dvd m<sup>2</sup>"
have dvd m: "2 dvd m"
```

```
hence "2 dvd n<sup>2</sup>" ..
have "2 dvd n"
have "2 dvd gcd m n"
thus False
qed
```

```
theorem sqrt2 not rational:
  "sqrt (real 2) ∉ Q"
proof
  assume "sqrt (real 2) \in \mathbb{Q}"
  then obtain m n :: nat where
    n nonzero: "n ≠ 0" and sqrt rat: "¦sqrt (real 2)¦ = real m / real
    and lowest terms: "acd m n = 1"
  from n nonzero and sqrt rat have "real m = \{sqrt (real 2)\} * real n
  then have "real (m^2) = (sqrt (real 2))^2 * real <math>(n^2)"
    by (auto simp add: power2 eq square)
  also have "(sqrt (real 2))<sup>2</sup> = real 2" by simp
  also have "... * real (m^2) = real (2 * n^2)" by simp
  finally have eq: m^2 = 2 * n^2...
  hence "2 dvd m<sup>2</sup>"
  with two is prime have dvd m: "2 dvd m" by (rule prime dvd power tw
  then obtain k where "m = 2^* k"
  with eq have "2 * n^2 = 2^2 * k^2" by (auto simp add: power2 eq square
  hence "n^2 = 2 * k^2" by simp
  hence "2 dvd n<sup>2</sup>" .
  with two is prime have "2 dvd n" by (rule prime dvd power two)
  with dvd m have "2 dvd qcd m n" by (rule qcd greatest)
  with lowest terms have "2 dvd 1" by simp
  thus False \overline{b}v arith
qed
```

Proof Assistant (2/2)

- Keep track of theories, definitions, assumptions
 - set up a theory that describes mathematical concepts (or models a computer system)
 - express logical properties of the objects
- Interaction proof editing
 - typically interactive
 - specified theory and proofs can be edited
 - provides information about required proof obligations
 - allows further refinement of the proof
 - often manually providing a direction in which to proceed.
- Automation proof search
 - various strategies
 - decision procedures
- Proof checking
 - checking of complete proofs
 - sometimes providing certificates of correctness
- Why should we trust it?
 - small core

Decidability!

- Validity of formulas is undecidable
- (for non-trivial logical systems)

Automated Theorem Provers

- Specific domains
- Adjust your problem
- Answers: Valid (Theorem with proof)
- Or: Countersatisfiable (Possibly with counter-model)

Proof Assistants

- Generally applicable
- Direct modelling of problems
- Interactive

(Many already covered in the courses in past few days)

ATPs (tomorrow)

- Built in automation (model elimination, resolution)
- Vampire, Eprover, SPASS, ...
- Applications: Robbin's conjecture, Programs, and AIM

Computer Science

- Modelling and specifying systems
- Proving properties of systems
- Proving software correct

Mathematics

- Defining concepts and theories
- Proving (mostly verifying) proofs
- (currently less common)

Theorems and programs that use ITP

Theorems

- Kepler Conjecture
- 4 color theorem
- Feit-Thomson theorem (2012)

Software

- Processors and Chips
- Security Protocols
- Project Cristal (Comp-Cert)
- L4-Verified
- Java Bytecode

Freek Wiedijk's list of 100 theorems

HOL Light	86
Isabelle	81
MetaMath	71
Coq	69
Mizar	69
any	94

Coverage by other tools

- much less as single steps
- (actually hard to compare)

[Wiedijk'15]

Complicated Proofs (Math, Computer Science)

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 - formal logical system

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Human proofs

Proof skeletons

Filling in the gaps: most of the work

- Small intermediate steps
- Sometimes also hard ones

Complicated Proofs (Math, Computer Science)

Proof Assistant

- a computer program to assist a mathematician
 - keep track of theories, definitions, assumptions, check individual steps, provide decision procedures
- in the production of a proof
- that is mechanically checked
 - formal logical system

- Proof skeletons
- Filling in the gaps: most of the work
 - Small intermediate steps
 - General Purpose Automation!
 - Sometimes also hard ones
 - Selected domains

What is Machine Learning?

Tuning a big number of parameters

Algorithms that improve their performance based on data

- Face detection
- Recommender systems
- Speech recognition
- Stock prediction
- Spam detection
- Molecule modeling
- Automated translation

· ...

Tasks involving logical inference

	Natural	language	question	answering
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- Knowledge base completion
- Automated translation

Games

AlphaGo (Zero) problems similar to proving

- Node evaluation
- Policy decisions

[Sukhbaatar+2015]

[Socher+2013]

[Wu+2016]

[Silver+2016]

High-level AI guidance

- premise selection: select the right lemmas to prove a new fact
- based on suitable features (characterizations) of the formulas
- and on learning lemma-relevance from many related proofs
- tactic selection

Mid-level AI guidance

- learn good ATP strategies/tactics/heuristics for classes of problems
- learning lemma and concept re-use
- learn conjecturing

Low-level AI guidance

- guide (almost) every inference step by previous knowledge
- good proof-state characterization and fast relevance

Is my conjecture true?

Problems for Machine Learning

- Is my conjecture true?
- Is a statement is useful?

 $a^n + b^n = c^n$

Problems for Machine Learning

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 - For a conjecture

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- What are the dependencies of statement? (premise selection)

$$a^n + b^n = c^n$$

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- Should a theorem be named? How?

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Problems for Machine Learning

- Is my conjecture true?
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- Should a theorem be named? How?
- What should the next proof step be?
 - Tactic? Instantiation?

$$a^n + b^n = c^n$$

- Is my conjecture true?
- Is a statement is useful?
 - For a conjecture
- What are the dependencies of statement? (premise selection)
- Should a theorem be named? How?
- What should the next proof step be?
 - Tactic? Instantiation?
- What new problem is likely to be true?
 - Intermediate statement for a conjecture

$$a^n + b^n = c^n$$

Intuition

Given:

- set of theorems T (together with proofs)
- conjecture c

Find: minimal subset of T that can be used to prove c

More formally

 $\underset{t\subseteq T}{\arg\min\{|t| \mid t\vdash c\}}$

(or \emptyset if not provable)

Note: implicit assumption on a proving system. ATP in practice.

Multi-label classification

Input: set of samples S, where samples are triples s, F(s), L(s)

- *s* is the sample ID
- F(s) is the set of features of s
- L(s) is the set of labels of s

Output: function f : features \rightarrow labels

Predicts n labels (sorted by relevance) for set of features

Sample features

Sample add_comm (a + b = b + a) characterized by:

- F(add_comm) = { "+", "=", "num" }
- L(add_comm) = {num_induct, add_0, add_suc, add_def}

Not exactly the usual machine learning problem

Labels correspond to premises and samples to theorems

Very often same

Labels correspond to premises and samples to theorems

Very often same

Similar theorems are likely to be useful in the proof

Also likely to have similar premises

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Theorems sharing logical features are similar

Theorems sharing rare features are very similar

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Temporal order

- Recently considered theorems and premises are important
- Also in evaluation

Needs efficient learning and prediction

- Frequent major data updates
- Automation cannot wait more than 10 seconds, often less

Multi-label classifier output

• Often asked for 1000 or more most relevant lemmas

Easy to get many interesting features

- Complicated feature relations
- PCA / LSA / …?

Premise Selection

- Syntactic methods
 - Neighbours using various metrics
 - Recursive

SInE, MePo

- Naive Bayes, k-Nearest Neighbours
- Linear / Logistic Regression
 - Needs feature and theorem space reduction
 - Kernel-based multi-output ranking
- Decision Trees (Random Forests)
- Neural Networks
 - Winnow, Perceptron
 - DeepMath

SNoW, MaLARea

k-Nearest Neighbours:

- finds a fixed number (k) of proved facts nearest to the conjecture c
- weight the dependencies each such fact f by the distance between f and c
- relevance is the sum of weights across the k nearest neighbors

Naive Bayes:

- probability of f being needed to prove c
- based on the previous use of f in proving conjectures similar to c
- assumes independence of features to use the Bayes theorem

MePo: (Meng–Paulson)

- score of a fact is r/(r + i), where r is the number of relevant features and i the number of irrelevant features
- iteratively select all top-scoring facts and add their features to the set of relevant features.

Combination

$$s(a,b) = \sum_{f \in F(a) \cap F(b)} 1$$

$$s(a,b) = \sum_{f \in F(a) \cap F(b)} w(f)$$

$$s(a,b) = \sum_{f \in F(a) \cap F(b)} w(f)^{\tau_1}$$

$$s(a,b) = \sum_{f \in F(a) \cap F(b)} w(f)^{\tau_1}$$

Relevance of fact a for goal g

$$\begin{pmatrix} \sum_{b \in N \mid a \in D(b)} \frac{s(b,g)}{\mid D(b) \mid} \end{pmatrix}$$

$$s(a,b) = \sum_{f \in F(a) \cap F(b)} w(f)^{\tau_1}$$

Relevance of fact a for goal g

$$\left(\sum_{b \in N \mid a \in D(b)} \frac{s(b,g)}{|D(b)|} \right) + \begin{cases} s(a,g) & \text{if } a \in N \\ 0 & \text{otherwise} \end{cases}$$

$$s(a,b) = \sum_{f \in F(a) \cap F(b)} w(f)^{\tau_1}$$

Relevance of fact a for goal g

$$\left(\tau_2 \sum_{b \in N \mid a \in D(b)} \frac{s(b,g)}{|D(b)|}\right) + \begin{cases} s(a,g) & \text{if } a \in N \\ 0 & \text{otherwise} \end{cases}$$

k-NN (2/2)

```
let knn eval csvms (svm ths, svm wght) deps maxth no adv =
 let neighbours = Array.init maxth (fun j -> (j, 0.)) in
 let ans = Array.copy neighbours in
  (* for each symbol, increase the importance of the theorems
     which contain the symbol by a given symbol weight *)
 List.iter (fun svm ->
     let ths = svm ths svm and weight = svm wght svm in
     List.iter (fun th ->
         if th < maxth then map snd neighbours th ((+.) (weight ** 6.0))) ths) csyms;
 Array.fast sort sortfun neighbours;
 let no recommends = ref 0 in
 let add ans k i o =
   if snd (ans.(i)) <= 0. then begin
     incr no recommends;
     map_snd ans i (fun _ -> float_of_int (age k) +. o))
   end else map snd ans i ((+.) o) in
  (* Additionally stop when given no_recommends reached *)
 Array.iteri (fun k (nn, o) ->
   add ans k nn o:
   let ds = deps nn in
   let ol = 2.7 *. o /. (float of int (List.length ds)) in
   List.iter (fun d -> if d < maxth then add ans k d ol) ds;
 ) neighbours;
 Array.fast sort sortfun ans:
```

$$P(f \text{ is relevant for proving } g)$$

$$= P(f \text{ is relevant } | g' \text{s features})$$

$$= P(f \text{ is relevant } | f_1, \dots, f_n)$$

$$\propto P(f \text{ is relevant}) \prod_{i=1}^n P(f_i | f \text{ is relevant})$$

$$\propto \#f \text{ is a proof dependency} \cdot \prod_{i=1}^n \frac{\#f_i \text{ appears when } f \text{ is a proof dependency}}{\#f \text{ is a proof dependency}}$$

extended features $\overline{F}(a)$ of a fact *a* features of *a* and of the facts that were proved using *a*

More precise estimation of the relevance of ϕ to prove $\gamma {:}$

$$P(a \text{ is used in } \psi \text{'s proof})$$

$$\cdot \prod_{f \in F(\gamma) \cap \overline{F}(a)} P(\psi \text{ has feature } f \mid a \text{ is used in } \psi \text{'s proof})$$

$$\cdot \prod_{f \in F(\gamma) - \overline{F}(a)} P(\psi \text{ has feature } f \mid a \text{ is not used in } \psi \text{'s proof})$$

$$\cdot \prod_{f \in \overline{F}(a) - F(\gamma)} P(\psi \text{ does not have feature } f \mid a \text{ is used in } \psi \text{'s proof})$$

extended features $\overline{F}(a)$ of a fact a

features of a and of the facts that were proved using a (only one iteration)

More precise estimation of the relevance of ϕ to prove γ :

$$\begin{split} & P(\textbf{a} \text{ is used in } \psi \text{'s proof}) \\ & \cdot \prod_{f \in F(\gamma) \cap \overline{F}(a)} P(\psi \text{ has feature } f \mid \textbf{a} \text{ is used in } \psi \text{'s proof}) \\ & \cdot \prod_{f \in F(\gamma) - \overline{F}(a)} P(\psi \text{ has feature } f \mid \textbf{a} \text{ is not used in } \psi \text{'s proof}) \\ & \cdot \prod_{f \in \overline{F}(a) - F(\gamma)} P(\psi \text{ does not have feature } f \mid \textbf{a} \text{ is used in } \psi \text{'s proof}) \end{split}$$

Update two functions (tables):

- t(a): number of times a fact a was dependency
- *s*(*a*, *f*):

number of times a fact \boldsymbol{a} was dependency of a fact described by feature \boldsymbol{f} Then:

$$P(a \text{ is used in a proof of (any) } \psi) = \frac{t(a)}{K}$$

$$P(\psi \text{ has feature } f \mid a \text{ is used in } \psi \text{'s proof}) = \frac{s(a, f)}{t(a)}$$

$$P(\psi \text{ does not have feature } f \mid a \text{ is used in } \psi \text{'s proof}) = 1 - \frac{s(a, f)}{t(a)}$$

$$\approx 1 - \frac{s(a, f) - 1}{t(a)}$$

Naive Bayes "in practice"

```
double NaiveBayes::score(sample t i, set<feature t> symh) const {
 // number of times current theorem was used as dependency
 const long n
                    = tfrea[i]:
 const auto sfreqh = sfreq[i];
 double s = 30 * log(n);
 for (const auto sv : sfreqh) {
   // sv.first ranges over all features of theorems depending on i
   // sv.second is the number of times sv.first appears among theorems
   // depending on i
   double sfreqv = sv.second;
   // if sv.first exists in guery features
   if (symh.erase(sv.first) == 1)
     s += tfidf.get(sv.first) * log (5 * sfreqv / n);
   else
     s += tfidf.get(sv.first) * 0.2 * log (1 + (1 - sfreqv) / n);
  }
 // for all query features that did not appear in features of dependencies
 // of current theorem
 for (const auto f : symh) s -= tfidf.get(f) * 18;
```

return s;

Basic algorithm

If symbol s is d-relevant and appears in axiom a, then a and all symbols in a become d + 1-relevant.

Problem: Common Symbols

- Simple *relevance* usually selects all axioms
- Because of common symbols, such as subclass or subsumes subclass (beverage, liquid).
 subclass (chair, furniture).

Solution: Trigger based selection "appears" is changed to "triggers"

But how to know if *s* is common? Approximate by number of occurrences in the current problem

SInE: Tolerance

- Only symbols with *t*-times more occurrences than the least common symbol trigger an axiom
- For $t = \infty$ this is the same as relevance
 - 1: subclass(X,Y) ∧ subclass(Y,Z) → subclass(X,Z) subclass(petrol,liquid) ¬subclass(stone,liquid)
 - 2: subclass(beverage,liquid)
 - subclass(beer,beverage) subclass(guinness,beer) subclass(pilsner,beer)

? subclass(beer,liquid)

Occ.	Symbols
7	subclass
3	liquid, beer
2	beverage
1	petrol, stone, guinness, pilsner

[Hoder]

SInE in E

Implementation: GSInE in e_axfilter

- Parameterizable filters
 - Different generality measures (frequency count, generosity, benevolence)
 - Different limits (absolute/relative size, # of iterations)
 - Different seeds (conjecture/hypotheses)
- Efficient implementation
 - E data types and libraries
 - Indexing (symbol \rightarrow formula, formula \rightarrow symbol)
- Multi-filter support
 - Parse & index once (amortize costs)
 - Apply different independent filters

Primary use

- Initial over-approximation (efficiently reduce HUGE input files to manageable size)
- Secondary use: Filtering for individual E strategies

Premises: Classification

- Dimensions in the input
- Matrix QR decomposition

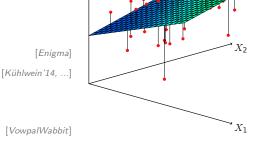
Probabilities: Logistic

Non-linearity

- Kernels
- Multi-output Ranking

State space reduction

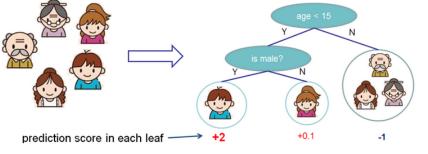
- Random projections
- Decomposition



Y

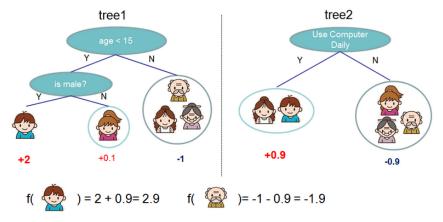
Input: age, gender, occupation, ...

Does the person like computer games



[Chen, Guestrin]

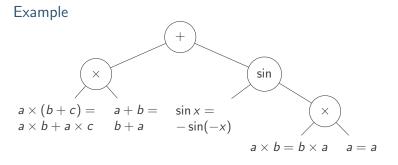
Decision Trees (2/2)

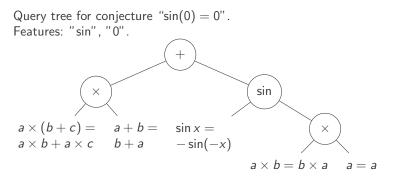


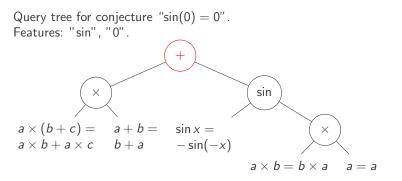
[Chen, Guestrin]

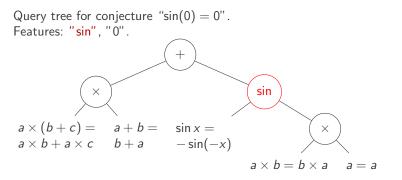
Definition

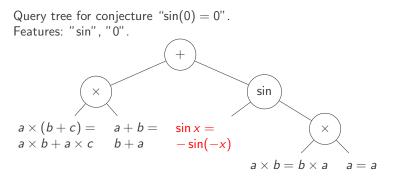
- each leaf stores a set of samples
- each branch stores a feature *f* and two subtrees, where:
 - the left subtree contains only samples having f
 - the right subtree contains only samples not having f

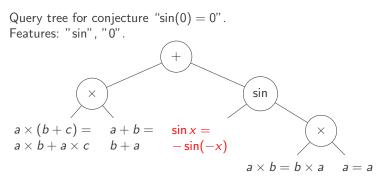




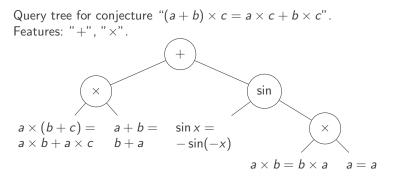


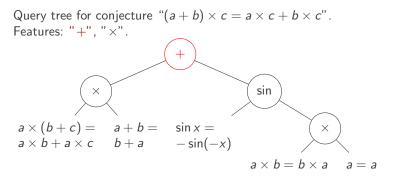


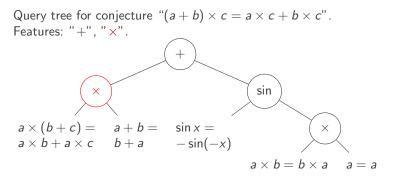


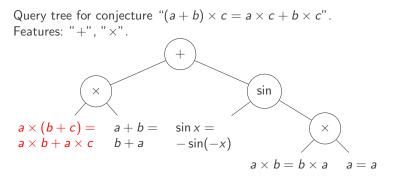


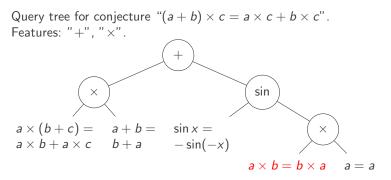
The overall result will be the premises of $\sin x = -\sin(-x)$.



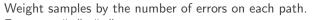








 $a \times b = b \times a$ is not considered!



Features: "+", "×". 0 + 1 0×1 $a \times (b+c) = a+b = \sin x = 1$ $a \times b+a \times c \quad b+a \quad -\sin(-x)$ $a \times b = b \times a \quad a = a$

Agrawal et al.

- Take *n* random features from samples and choose feature with lowest Gini impurity (probability of mis-labeling)
- Problem: Gini impurity calculation slow
- Choose feature that divides samples most evenly $(|S_f| pprox |S_{
 eg f}|)$

Online / Offline forests

tree is updated or completely rebuilt

[Agraval, Saffari]

Approach for premise selection

- when a branch learns new samples, check whether the branch feature is still an optimal splitting feature wrt. the new data
- if yes, update subtrees with new data
- if no, rebuild tree

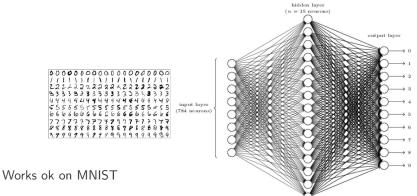
learning takes 21 min for the Mizar dataset...

Neural Networks (Introduction in 2 slides)

Neural Networks (Introduction in 2 slides)

Recognize a handwritten character

Measure: recognition rate



Modelling of Neurophysiological Networks (1950s – 1960s)

- Simple networks of individual perceptrons, with basic learning
- Severe limitations
 [Minsky, Papert]

Paralled Distributed Processing (1990s)

- rejuvenated interest
- But statistical algorithms were comparably powerful (SVM)

Deep Learning (2010s)

- Data-oriented algorithms
- Data and processing were a limitation before

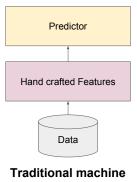
[Rumelhart, MacClelland]

Perceptrons implement linear separators, but:

- Every continuous function modeled with three layers (= 1 hidden)
- Every function can be modeled with four layers
- But the layers are assumed to be arbitrarily large!

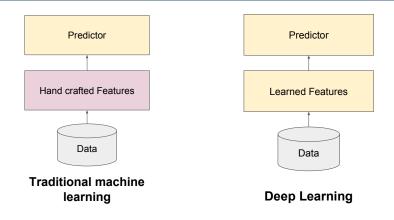
(Results recently formalized)

Deep Learning vs Shallow Learning

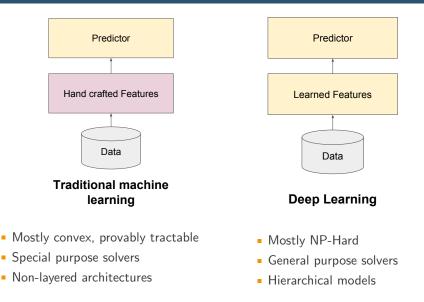


learning

Deep Learning vs Shallow Learning

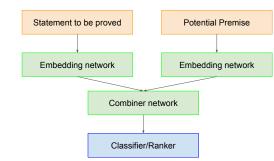


Deep Learning vs Shallow Learning



Simple classifier on top of concatenated embeddings

- different model of premise selection
- trained to estimate usefulness
- positive and negative examples



Architecture

Deep Learning for Mizar Lemma Selection

- No hand-engineered features
- Comparison of various neural architectures
- Semantic-aware definition embeddings
- Complementary to previous approaches
- Can be ensembled

The Mizar Mathematical Library (MML) is a corpus of formal mathematical proofs, containing 57,917 theorems from a wide varity of mathematical subjects. We worked with a version converted to first order logic in the TPTP format.

```
:: t99_jordan: Jordan curve theorem in Mizar
for C being Simple_closed_curve holds C is Jordan;
```

```
:: Translation to first order logic
fof(t99_jordan, axiom, (! [A] : ( (v1_topreal2(A) & m1_subset_1(A,
k1_zfmisc_1(u1_struct_0(k15_euclid(2))))) => v1_jordan1(A)) ) ).
```

Figure 1: (top) The final statement of the Mizar formalization of the Jordan curve theorem. (bottom) The translation to first-order logic, using name mangling to ensure uniqueness across the entire corpus.

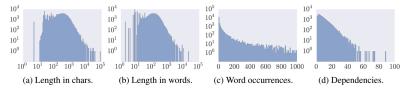


Figure 2: Histograms of statement lengths, occurrences of each word, and statement dependencies in the Mizar corpus translated to first order logic. The wide length distribution poses difficulties for RNN models and batching, and many rarely occurring words make it important to take definitions of words into account.

DeepMath: Problem, Metric, Model

Definition (Premise selection problem). *Given a large set of premises* \mathcal{P} , an ATP system A with given resource limits, and a new conjecture C, predict those premises from \mathcal{P} that will most likely lead to an automatically constructed proof of C by A.

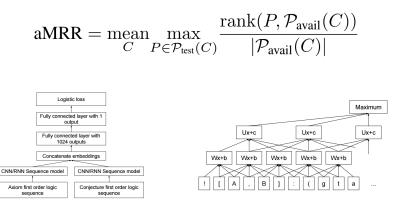


Figure 3: (left) Our network structure. The input sequences are either character-level (section 5.1) or word-level (section 5.2). We use separate models to embed conjecture and axiom, and a logistic layer to predict whether the axiom is useful for proving the conjecture. (right) A convolutional model.

Recurrent Neural Networks (RNN)

process sequences by feeding back the output into the next input

Long-Short Term Memory (LSTM)

add forgetting to RNNs

DeepMath: Architectures

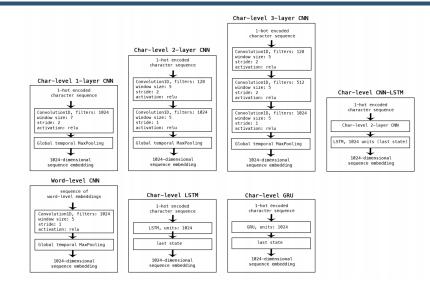


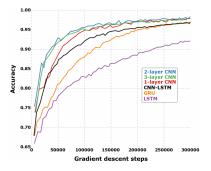
Figure 4: Specification of the different embedder networks.

Cutoff	k-NN Baseline (%)	char-CNN (%)	word-CNN (%)	def-CNN-LSTM (%)	def-CNN (%)	def+char-CNN (%)
16	674 (24.6)	687 (25.1)	709 (25.9)	644 (23.5)	734 (26.8)	835 (30.5)
32	1081 (39.4)	1028 (37.5)	1063 (38.8)	924 (33.7)	1093 (39.9)	1218 (44.4)
64	1399 (51)	1295 (47.2)	1355 (49.4)	1196 (43.6)	1381 (50.4)	1470 (53.6)
128	1612 (58.8)	1534 (55.9)	1552 (56.6)	1401 (51.1)	1617 (59)	1695 (61.8)
256	1709 (62.3)	1656 (60.4)	1635 (59.6)	1519 (55.4)	1708 (62.3)	1780 (64.9)
512	1762 (64.3)	1711 (62.4)	1712 (62.4)	1593 (58.1)	1780 (64.9)	1830 (66.7)
1024	1786 (65.1)	1762 (64.3)	1755 (64)	1647 (60.1)	1822 (66.4)	1862 (67.9)

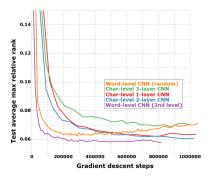
Table 1: Results of ATP premise selection experiments with hard negative mining on a test set of 2,742 theorems.

- E-prover proved theorem percentages
- Union of all methods: 80.9%
- Union of deep network methods: 78.4%

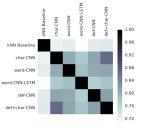
DeepMath: Accuracy



(a) Training accuracy for different character-level models without hard negative mining. Recurrent models seem underperform, while pure convolutional models yield the best results. For each architecture, we trained three models with different random initialization seeds. Only the best runs are shown on this graph; we did not see much variance between runs on the same architecture.



(b) Test average max relative rank for different models without hard negative mining. The best is a word-level CNN using definition embeddings from a character-level 2-layer CNN. An identical wordembedding model with random starting embedding overfits after only 250,000 iterations and underperforms the best character-level model.



(a) Jaccard similarities between proved sets of conjectures across models. Each of the neural network model prediction are more like each other than those of the k-NN baseline.

Model	Test min average relative rank
char-CNN	0.0585
word-CNN	0.06
def-CNN-LSTM	0.0605
def-CNN	0.0575

(b) Best sustained test results obtained by the above models. Lower values are better. This was monitored continuously during training on a holdout set with 400 theorems, using all true positive premises and 128 randomly selected negatives. In this setup, the lowest attainable max average relative rank with perfect predictions is 0.051.

Hard Negatives

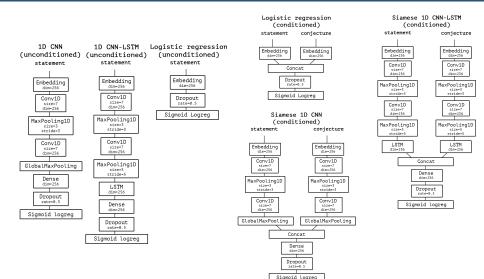
HOLStep Dataset

- Intermediate steps of the Kepler proof
- Only relevant proofs of reasonable size
- Annotate steps as useful and unused
 - Same number of positive and negative
- Tokenization and normalization of statements

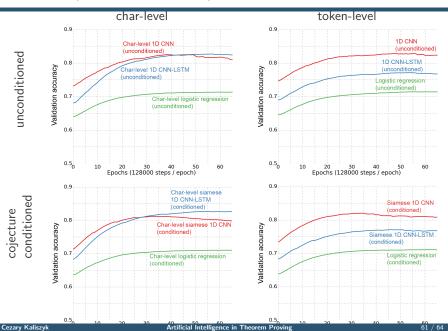
5		Train	Test	Positive	Negative	
	Examples	2013046	196030	1104538	1104538	-
	Avg. length	503.18	440.20	535.52	459.66	
	Avg. tokens	87.01	80.62	95.48	77.40	
	Conjectures	9999	1411	-	-	
	Avg. deps	29.58	22.82	-	-	

Statistics

Considered Models



Baselines (Training Profiles)



Proof by contradiction

- Assume that the conjecture does not hold
- Derive that axioms and negated conjecture imply \perp

Saturation

- Convert problem to CNF
- Enumerate the consequences of the available clauses
- Goal: get to the empty clause

Redundancies

Simplify or eliminate some clauses (contract)

Today

- Theorem proving systems
- Machine learning problems
- Lemma relevance
- Deep learning for theorem proving

Tomorrow

- Guided Automated Reasoning
- More human-like proof
- Logical translations
- Unsupervised methods