

## Exercise 6: Containment

### Task 1: Containing Choice

The goal in this exercise is to prove Lemma 6.5.

- a) Show the equivalence stated in the lemma.
- b) Construct a  $k$ -bit  $\text{MUX}_M$  implementation out of two  $(k - 1)$ -bit  $\text{MUX}_M$  implementations and a CMUX. (Hint: To show correctness, make a case distinction on the  $k^{\text{th}}$  control bit, which is fed to the CMUX.)
- c) What is the size of the resulting  $\text{MUX}_M$  implementation when applying the construction from b) recursively?

### Task 2: Copy and Conquer

Masking registers allow to mask internal metastability, resulting in, e.g., the sequence  $0 \dots 0M1 \dots 1$  when reading sequentially from a mask-0 register. The key property for this exercise is that there is only a single M read from the register. We consider a function  $f: \{0, 1\}^n \rightarrow \{0, 1\}$  in this exercise.

- a) Suppose inputs are stored in masking registers, which we read  $2n + 1$  times, each time making a separate copy  $x^{(i)}$ ,  $i \in \{1, \dots, 2n + 1\}$ , of the input  $x$ . If  $f_M(x) \neq M$ , what can you say about the collection of  $2n + 1$  outputs generated from feeding each  $x^{(i)}$  to a copy of a (non-containing) circuit implementing  $f$ ?
- b) Come up with a small circuit that sorts its  $n$  inputs according to the total order  $0 \leq M \leq 1$ . (Hint: Figure out a solution sorting two values and then plug it into a sorting network to get the general circuit. You don't have to (re)invent sorting networks, you may just point to a reference.)
- c) Combine a) and b) to derive a circuit implementing  $f_M$  from any (non-containing) circuit implementing  $f$ ! Your solution should only be by a factor of  $n^{\mathcal{O}(1)}$  larger than to the non-containing solution.

### Task 3\*: Clocked Circuits

- a) How would a model for clocked circuits based on the same worst-case assumptions look like? (Hint: Reading up on it is fine.)
- b) Standard registers, when being read, will output M if they're internally metastable and 0 or 1, respectively, when they're stable. Show that they add no power in terms of the functions that can be computed! (Hint: Unroll the circuit, i.e., perform the multi-round computation in a single round with a larger circuit.)
- c) In Task 2, you saw that masking registers allow for more efficient metastability-containing circuits. Show that they are also computationally more powerful, i.e., they can compute functions that cannot be computed with masking registers! (Hint: You already used this in Task 2!)