Parameterized Algorithms using Matroids
Lecture 0: Matroid Basics

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Kruskal’s Greedy Algorithm for MWST

Let $G = (V, E)$ be a connected undirected graph and let $w : E \to \mathbb{R}^{\geq 0}$ be a weight function on the edges. Kruskal’s so-called greedy algorithm is as follows. The algorithm consists of selecting successively edges $e_1, e_2, \ldots, e_r$. If edges $e_1, e_2, \ldots, e_k$ has been selected, then an edge $e \in E$ is selected so that:

1. $e \notin \{e_1, \ldots, e_k\}$ and $\{e, e_1, \ldots, e_k\}$ is a forest.
2. $w(e)$ is as small as possible among all edges $e$ satisfying (1).

We take $e_{k+1} := e$. If no $e$ satisfying (1) exists then $\{e_1, \ldots, e_k\}$ is a spanning tree.
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It is obviously not true that such a greedy approach would lead to an optimal solution for any combinatorial optimization problem.
Consider **Maximum Weight Matching** problem.

- Application of the greedy algorithm gives *cd* and *ab*.
- However, *cd* and *ab* do not form a matching of maximum weight.
It is obviously not true that such a greedy approach would lead to an optimal solution for any combinatorial optimization problem.

It turns out that the structures for which the greedy algorithm does lead to an optimal solution, are the matroids.
Definition

A pair $M = (E, \mathcal{I})$, where $E$ is a ground set and $\mathcal{I}$ is a family of subsets (called independent sets) of $E$, is a matroid if it satisfies the following conditions:

(I1) $\emptyset \in \mathcal{I}$ or $\mathcal{I} \neq \emptyset$.

(I2) If $A' \subseteq A$ and $A \in \mathcal{I}$ then $A' \in \mathcal{I}$.

(I3) If $A, B \in \mathcal{I}$ and $|A| < |B|$, then $\exists e \in (B \setminus A)$ such that $A \cup \{e\} \in \mathcal{I}$.

The axiom (I2) is also called the hereditary property and a pair $M = (E, \mathcal{I})$ satisfying (I1) and (I2) is called hereditary family or set-family.
Matroids

Definition

A pair \( M = (E, \mathcal{I}) \), where \( E \) is a ground set and \( \mathcal{I} \) is a family of subsets (called independent sets) of \( E \), is a matroid if it satisfies the following conditions:

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Rank and Basis

**Definition**

A pair $M = (E, \mathcal{I})$, where $E$ is a ground set and $\mathcal{I}$ is a family of subsets (called independent sets) of $E$, is a *matroid* if it satisfies the following conditions:

(I1) $\varnothing \in \mathcal{I}$ or $\mathcal{I} \neq \emptyset$.

(I2) If $A' \subseteq A$ and $A \in \mathcal{I}$ then $A' \in \mathcal{I}$.

(I3) If $A, B \in \mathcal{I}$ and $|A| < |B|$, then $\exists \ e \in (B \setminus A)$ such that $A \cup \{e\} \in \mathcal{I}$.

An inclusion wise maximal set of $\mathcal{I}$ is called a *basis* of the matroid. Using axiom (I3) it is easy to show that all the bases of a matroid have the same size. This size is called the *rank* of the matroid $M$, and is denoted by $\text{rank}(M)$. 
Matroids and Greedy Algorithms

Let $M = (E, \mathcal{I})$ be a set family and let $w : E \to \mathbb{R}_{\geq 0}$ be a weight function on the elements.

**Objective:** Find a set $Y \in \mathcal{I}$ that maximizes
$$w(Y) = \sum_{y \in Y} w(y).$$

The **greedy algorithm** consists of selecting successively $y_1, \ldots, y_r$ as follows. If $y_1, \ldots, y_k$ has been selected, then choose $y \in E$ so that:

1. $y \notin \{y_1, \ldots, y_k\}$ and $\{y, y_1, \ldots, y_k\} \in \mathcal{I}$.
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We stop if no $y$ satisfying (1) exists.
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We stop if no $y$ satisfying (1) exists.
Matroids and Greedy Algorithms

**Theorem:** A set family $M = (E, \mathcal{I})$ is a matroid

if and only if

the greedy algorithm leads to a set $Y$ in $\mathcal{I}$ of maximum weight $w(Y)$, for each weight function $w : E \rightarrow \mathbb{R}_{\geq 0}$. 
Examples Of Matroids
Uniform Matroid

A pair \( M = (E, \mathcal{I}) \) over an \( n \)-element ground set \( E \), is called a uniform matroid if the family of independent sets is given by

\[
\mathcal{I} = \left\{ A \subseteq E \mid |A| \leq k \right\},
\]

where \( k \) is some constant. This matroid is also denoted as \( U_{n,k} \).

Eg: \( E = \{1, 2, 3, 4, 5\} \) and \( k = 2 \) then

\[
\mathcal{I} = \left\{ \emptyset, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{1, 2\}, \{1, 3\}, \{1, 4\}, \{1, 5\}, \{2, 3\}, \{2, 4\}, \{2, 5\}, \{3, 4\}, \{3, 5\}, \{4, 5\} \right\}
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Uniform Matroid

A pair $M = (E, \mathcal{I})$ over an $n$-element ground set $E$, is called a *uniform matroid* if the family of independent sets is given by

$$\mathcal{I} = \left\{ A \subseteq E \mid |A| \leq k \right\},$$

where $k$ is some constant. This matroid is also denoted as $U_{n,k}$. Eg: $E = \{1, 2, 3, 4, 5\}$ and $k = 2$ then

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Partition Matroid

A partition matroid \( M = (E, \mathcal{I}) \) is defined by a ground set \( E \) being partitioned into (disjoint) sets \( E_1, \ldots, E_\ell \) and by \( \ell \) non-negative integers \( k_1, \ldots, k_\ell \). A set \( X \subseteq E \) is independent if and only if \( |X \cap E_i| \leq k_i \) for all \( i \in \{1, \ldots, \ell\} \). That is,

\[
\mathcal{I} = \left\{ X \subseteq E \mid |X \cap E_i| \leq k_i, \ i \in \{1, \ldots, \ell\} \right\}.
\]

- If \( X, Y \in \mathcal{I} \) and \( |Y| > |X| \), there must exist \( i \) such that \( |Y \cap E_i| > |X \cap E_i| \) and this means that adding any element \( e \) in \( E_i \cap (Y \setminus X) \) to \( X \) will maintain independence.
- \( M \) in general would not be a matroid if \( E_i \) were not disjoint. Eg: \( E_1 = \{1, 2\} \) and \( E_2 = \{2, 3\} \) and \( k_1 = 1 \) and \( k_2 = 1 \) then both \( Y = \{1, 3\} \) and \( X = \{2\} \) have at most one element of each \( E_i \) but one can’t find an element of \( Y \) to add to \( X \).
Partition Matroid

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Graphic Matroid

Given a graph $G$, a graphic matroid is defined as $M = (E, \mathcal{I})$ where and

- $E = E(G)$ – edges of $G$ are elements of the matroid
- $\mathcal{I} = \{ F \subseteq E(G) : F \text{ is a forest in the graph } G \}$
Given a graph $G$, a co-graphic matroid is defined as $M = (E, I)$ where

- $E = E(G)$ – edges of $G$ are elements of the matroid

- $I = \left\{ S \subseteq E(G) : G \setminus S \text{ is connected} \right\}$
Transversal Matroid

Let $G$ be a bipartite graph with the vertex set $V(G)$ being partitioned as

A

B

\[ A \quad \text{B} \]
Transversal Matroid

Let $G$ be a bipartite graph with the vertex set $V(G)$ being partitioned as $A$ and $B$. The *transversal matroid* $M = (E, \mathcal{I})$ of $G$ has $E = A$ as its ground set,

$$\mathcal{I} = \left\{ X \mid X \subseteq A, \text{ there is a matching that covers } X \right\}$$
Gammoids

Let $D = (V, A)$ be a directed graph, and let $S \subseteq V$ be a subset of vertices. A subset $X \subseteq V$ is said to be linked to $S$ if there are $|X|$ vertex disjoint paths going from $S$ to $X$.

The paths are disjoint, not only internally disjoint. Furthermore, zero-length paths are also allowed if $X \cap S = \emptyset$.

Given a digraph $D = (V, A)$ and subsets $S \subseteq V$ and $T \subseteq V$, a gammoid is a matroid $M = (E, \mathcal{I})$ with $E = T$ and

$$\mathcal{I} = \left\{ X \mid X \subseteq T \text{ and } X \text{ is linked to } S \right\}$$
Gammoids

Let $D = (V, A)$ be a directed graph, and let $S \subseteq V$ be a subset of vertices. A subset $X \subseteq V$ is *said to be linked to* $S$ if there are $|X|$ vertex disjoint paths going from $S$ to $X$.

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Gammoid: Example
Strict Gammoids

Given a digraph $D = (V, A)$ and subset $S \subseteq V$, a strict gammoid is a matroid $M = (E, \mathcal{I})$ with $E = V$ and

$$\mathcal{I} = \left\{ X \mid X \subseteq V \text{ and } X \text{ is linked to } S \right\}$$
An Alternate Definition of Matroids
Let $E$ be a finite set and $\mathcal{B}$ be a family of subsets of $E$ that satisfies:

(B1) $\mathcal{B} \neq \emptyset$.
(B2) If $B_1, B_2 \in \mathcal{B}$ then $|B_1| = |B_2|$.
(B3) If $B_1, B_2 \in \mathcal{B}$ and there is an element $x \in (B_1 \setminus B_2)$ then there is an element $y \in (B_2 \setminus B_1)$ so that $B_1 \setminus \{x\} \cup \{y\} \in \mathcal{B}$. 

Basis Family
(B3) If $B_1, B_2 \in \mathcal{B}$ and there is an element $x \in (B_1 \setminus B_2)$ then there is an element $y \in (B_2 \setminus B_1)$ so that $B_1 \setminus \{x\} \cup \{y\} \in \mathcal{B}$. 
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Given $\mathcal{B}$, we define

$$\mathcal{I} = \mathcal{I}(\mathcal{B}) = \left\{ I \mid I \subseteq B \text{ for some } B \in \mathcal{B} \right\}$$
Let $E$ be a finite set and $\mathcal{B}$ be the family of subsets of $E$ that satisfies:

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Family of Bases

Let $M = (E, \mathcal{I})$ be a matroid and $\mathcal{B}$ be the family of bases of $M$ – a family of maximal independent sets. Then $\mathcal{B}$ satisfies (B1), (B2) and (B3). That is,

(B1) $\mathcal{B} \neq \emptyset$.
(B2) If $B_1, B_2 \in \mathcal{B}$ then $|B_1| = |B_2|$.
(B3) If $B_1, B_2 \in \mathcal{B}$ and there is an element $x \in (B_1 \setminus B_2)$ then there is an element $y \in (B_2 \setminus B_1)$ so that $B_1 \setminus \{x\} \cup \{y\} \in \mathcal{B}$. 
Proof for B3

Let $M = (E, \mathcal{I})$ be a matroid and $\mathcal{B}$ be the family of bases of $M$ – a family of maximal independent sets.

\[
|B_1 - x| < |B_2|
\]
An Alternate Axiom System

**Theorem:** Let $E$ be a finite set and $\mathcal{B}$ be the family of subsets of $E$ that satisfies (B1), (B2) and (B3) then $M = (E, \mathcal{I})$ is a matroid and $\mathcal{B}$ is the family of bases of this matroid. Recall, that

$$\mathcal{I} = \mathcal{I}(\mathcal{B}) = \left\{ I \mid I \subseteq B \text{ for some } B \in \mathcal{B} \right\}.$$
New Matroids from Old
Deletion and Contraction

Let $M = (E, \mathcal{I})$ be a matroid, and let $X$ be a subset of $E$.
Deleting $X$ from $M$ gives a matroid $M \setminus X = (E \setminus X, \mathcal{I}')$ such that $S \subseteq E \setminus X$ is independent in $M \setminus X$ if and only if $S \in \mathcal{I}$.

$$\mathcal{I}' = \{ S \mid S \subseteq E \setminus X, S \in \mathcal{I} \}$$

Contracting $X$ from $M$ gives a matroid $M/X = (E \setminus X, \mathcal{I}')$ such that $S \subseteq E \setminus X$ is independent in $M/X$ if and only if $S \cup X \in \mathcal{I}$.

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Deletion and Contraction

- **Deletion**: The bases of $M \setminus X$ are those bases of $M$ that do not contain $X$.

- **Contraction**: The bases of $M/X$ are those bases of $M$ that do contain $X$ with $X$ then removed from each such basis.
Let $M_1 = (E_1, \mathcal{I}_1)$, $M_2 = (E_2, \mathcal{I}_2)$, \ldots, $M_t = (E_t, \mathcal{I}_t)$ be $t$ matroids with $E_i \cap E_j = \emptyset$ for all $1 \leq i \neq j \leq t$.

The direct sum $M_1 \oplus \cdots \oplus M_t$ is a matroid $M = (E, \mathcal{I})$ with $E := \bigcup_{i=1}^{t} E_i$ and $X \subseteq E$ is independent if and only if for all $i \leq t$, $X \cap E_i \in \mathcal{I}_i$.

$$\mathcal{I} = \left\{ X \mid X \subseteq E, \ (X \cap E_i) \in \mathcal{I}_i, \ i \in \{1, \ldots, t\} \right\}$$
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\[
\mathcal{I} = \left\{ X \mid X \subseteq E, (X \cap E_i) \in \mathcal{I}_i, \ i \in \{1, \ldots, t\} \right\}
\]
Truncation

The \textit{t-truncation} of a matroid $M = (E, \mathcal{I})$ is a matroid $M' = (E, \mathcal{I}')$ such that $S \subseteq E$ is independent in $M'$ if and only if $|S| \leq t$ and $S$ is independent in $M$ (that is $S \in \mathcal{I}$).
Dual

Let $M = (E, \mathcal{I})$ be a matroid, $\mathcal{B}$ be the family of its bases and $\mathcal{B}^* = \{ E \setminus B \mid B \in \mathcal{B} \}$.

The dual of a matroid $M$ is $M^* = (E, \mathcal{I}^*)$, where $\mathcal{I}^* = \{ X \mid X \subseteq B, \ B \in \mathcal{B}^* \}$.

That is, $\mathcal{B}^*$ is a family of bases of $M^*$. 
Let $M = (E, \mathcal{I})$ be a matroid, $\mathcal{B}$ be the family of its bases and

$$\mathcal{B}^* = \left\{ E \setminus B \mid B \in \mathcal{B} \right\}.$$

The dual of a matroid $M$ is $M^* = (E, \mathcal{I}^*)$, where

$$\mathcal{I}^* = \left\{ X \mid X \subseteq B, \ B \in \mathcal{B}^* \right\}.$$

That is, $\mathcal{B}^*$ is a family of bases of $M^*$. 
Matroid Representation
Remark

- Need compact representation to for the family of independent sets.
- Also to be able to test easily whether a set belongs to the family of independent sets.
Linear Matroid

Let $A$ be a matrix over an arbitrary field $\mathbb{F}$ and let $E$ be the set of columns of $A$. Given $A$ we define the matroid $M = (E, \mathcal{I})$ as follows.

A set $X \subseteq E$ is independent (that is $X \in \mathcal{I}$) if the corresponding columns are *linearly independent* over $\mathbb{F}$.

\[
A = \begin{bmatrix}
\ast & \ast & \ast & \cdots & \ast \\
\ast & \ast & \ast & \cdots & \ast \\
\ast & \ast & \ast & \cdots & \ast \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
\ast & \ast & \ast & \cdots & \ast \\
\end{bmatrix}
\]

$\ast$ are elements of $\mathbb{F}$

The matroids that can be defined by such a construction are called *linear matroids*. 
Linear Matroid

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$$A = \begin{bmatrix}
* & * & * & \cdots & * \\
* & * & * & \cdots & * \\
* & * & * & \cdots & * \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
* & * & * & \cdots & * \\
\end{bmatrix} \quad * \text{ are elements of } F$$

The matroids that can be defined by such a construction are called \textit{linear matroids}. 
Linear Matroids and Representable Matroids

If a matroid can be defined by a matrix $A$ over a field $F$, then we say that the matroid is *representable* over $F$. 
A matroid $M = (E, \mathcal{I})$ is representable over a field $\mathbb{F}$ if there exist vectors in $\mathbb{F}^\ell$ that correspond to the elements such that the linearly independent sets of vectors precisely correspond to independent sets of the matroid.

Let $E = \{e_1, \ldots, e_m\}$ and $\ell$ be a positive integer.

$$
\begin{bmatrix}
    e_1 & e_2 & e_3 & \cdots & e_m \\
    1 & * & * & * & \cdots & * \\
    2 & * & * & * & \cdots & * \\
    3 & * & * & * & \cdots & * \\
    \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
    \ell & * & * & * & \cdots & * \\
\end{bmatrix}_{\ell \times m}
$$

A matroid $M = (E, \mathcal{I})$ is called representable or linear if it is representable over some field $\mathbb{F}$. 
Linear Matroids and Representable Matroids

A matroid $M = (E, I)$ is representable over a field $\mathbb{F}$ if there exist vectors in $\mathbb{F}^\ell$ that correspond to the elements such that the linearly independent sets of vectors precisely correspond to independent sets of the matroid.

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e_1 & e_2 & e_3 & \cdots & e_m \\
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2 & * & * & * & \cdots & * \\
3 & * & * & * & \cdots & * \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
\ell & * & * & * & \cdots & * \\
\end{bmatrix}_{\ell \times m}
$$

A matroid $M = (E, I)$ is called \textit{representable} or \textit{linear} if it is representable over some field $\mathbb{F}$. 
Let $M = (E, \mathcal{I})$ be linear matroid and Let $E = \{e_1, \ldots, e_m\}$ and $d = \text{rank}(M)$.

We can always assume (using Gaussian Elimination) that the matrix has following form:

$$
\begin{bmatrix}
I_{d \times d} & D
\end{bmatrix}_{d \times m}
$$

Here $I_{d \times d}$ is a $d \times d$ identity matrix.
Let $M = (E, \mathcal{I})$ be linear matroid and Let $E = \{e_1, \ldots, e_m\}$ and $d = \text{rank}(M)$.
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Here $I_{d \times d}$ is a $d \times d$ identity matrix.
Direct Sum of Matroid

Let $M_1 = (E_1, \mathcal{I}_1)$, $M_2 = (E_2, \mathcal{I}_2)$, $\cdots$, $M_t = (E_t, \mathcal{I}_t)$ be $t$ matroids with $E_i \cap E_j = \emptyset$ for all $1 \leq i \neq j \leq t$. The direct sum $M_1 \oplus \cdots \oplus M_t$ is a matroid $M = (E, \mathcal{I})$ with $E := \bigcup_{i=1}^{t} E_i$ and $X \subseteq E$ is independent if and only if for all $i \leq t$, $X \cap E_i \in \mathcal{I}_i$.

Let $A_i$ be the representation matrix of $M_i = (E_i, \mathcal{I}_i)$ over the same field $\mathbb{F}$. Then,

$$A_M = \begin{pmatrix}
A_1 & 0 & 0 & \cdots & 0 \\
0 & A_2 & 0 & \cdots & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
0 & 0 & 0 & \cdots & A_t
\end{pmatrix}$$

is a representation matrix of $M_1 \oplus \cdots \oplus M_t$ over $\mathbb{F}$. 
Direct Sum of Matroid

Let $M_1 = (E_1, \mathcal{I}_1)$, $M_2 = (E_2, \mathcal{I}_2)$, $\cdots$, $M_t = (E_t, \mathcal{I}_t)$ be $t$ matroids with $E_i \cap E_j = \emptyset$ for all $1 \leq i \neq j \leq t$. The direct sum $M_1 \oplus \cdots \oplus M_t$ is a matroid $M = (E, \mathcal{I})$ with $E := \bigcup_{i=1}^{t} E_i$ and $X \subseteq E$ is independent if and only if for all $i \leq t$, $X \cap E_i \in \mathcal{I}_i$.

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$$A_M = \begin{pmatrix}
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0 & A_2 & 0 & \cdots & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
0 & 0 & 0 & \cdots & A_t
\end{pmatrix}$$

is a representation matrix of $M_1 \oplus \cdots \oplus M_t$ over $\mathbb{F}$. 
Deletion of a Matroid

Let $M = (E, \mathcal{I})$ be a matroid, and let $X$ be a subset of $E$. Deleting $X$ from $M$ gives a matroid $M \setminus X = (E \setminus X, \mathcal{I}')$ such that $S \subseteq E \setminus X$ is independent in $M \setminus X$ if and only if $S \in \mathcal{I}$.

$$\mathcal{I}' = \{ S \mid S \subseteq E \setminus X, S \in \mathcal{I} \}$$

Given a representation of $A_M$ of $M$, a representation of $M \setminus X$ can be obtained by deleting the columns corresponding to $X$.

$$A_M = \begin{bmatrix}
e_1 & e_2 & e_3 & \cdots & e_m \\
1 & * & * & * & \cdots & * \\
2 & * & * & * & \cdots & * \\
3 & * & * & * & \cdots & * \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
d & * & * & * & \cdots & * \\
\end{bmatrix}^{d \times m}$$
Deletion of a Matroid

Let $M = (E, \mathcal{I})$ be a matroid, and let $X$ be a subset of $E$. Deleting $X$ from $M$ gives a matroid $M \setminus X = (E \setminus X, \mathcal{I}')$ such that $S \subseteq E \setminus X$ is independent in $M \setminus X$ if and only if $S \in \mathcal{I}$.

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Given a representation of $A_M$ of $M$, a representation of $M \setminus X$ can be obtained by deleting the columns corresponding to $X$.

$$A_M = \begin{bmatrix}
  e_1 & e_2 & e_3 & \cdots & e_m \\
  \begin{array}{cccccc}
    1 & * & * & * & \cdots & * \\
    2 & * & * & * & \cdots & * \\
    3 & * & * & * & \cdots & * \\
    \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
    d & * & * & * & \cdots & * \\
  \end{array}
\end{bmatrix}_{d \times m}$$
Deletion of a Matroid

Let $X = \{e_2, e_3\}$.

$$A_M = \begin{bmatrix}
1 & * & * & * & \cdots & e_m \\
2 & * & * & * & \cdots & * \\
3 & * & * & * & \cdots & * \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
d & * & * & * & \cdots & * \\
\end{bmatrix}_{d \times m}$$

$$A_M = \begin{bmatrix}
e_1 & \cdots & e_m \\
1 & * & \cdots & * \\
2 & * & \cdots & * \\
3 & \cdots & * \\
\vdots & \vdots & \vdots & \vdots \\
d & * & \cdots & * \\
\end{bmatrix}_{d \times m}$$
Dual of a Matroid

Let $M = (E, \mathcal{I})$ be a matroid, $\mathcal{B}$ be the family of its bases and

$$\mathcal{B}^* = \left\{ E \setminus B \mid B \in \mathcal{B} \right\}.$$

The dual of a matroid $M$ is $M^* = (E, \mathcal{I}^*)$, where

$$\mathcal{I}^* = \left\{ X \mid X \subseteq B, \ B \in \mathcal{B}^* \right\}.$$

That is, $\mathcal{B}^*$ is a family of bases of $M^*$.

Let $A = [I_{d \times d} \mid D]$ represent the matroid $M$ then the matrix $A^* = [-D^T \mid I_{m-r \times m-r}]$ represents the dual matroid $M^*$. 
Dual of a Matroid: A concrete example

\[ A = \begin{bmatrix}
  a & b & c & d & e & f & g \\
  1 & 0 & 0 & 0 & 0 & 1 & 1 \\
  0 & 1 & 0 & 0 & 1 & 1 & 1 \\
  0 & 0 & 1 & 0 & 0 & 1 & 0 \\
  0 & 0 & 0 & 1 & 0 & 0 & 1
\end{bmatrix} \]

\{a, b, c, d\} is a basis of \( M \) then \( \{e, f, g\} \) is a basis of \( M^* \).

\[ A^* = \begin{bmatrix}
  a & b & c & d & e & f & g \\
  1 & 0 & 0 & 0 & 1 & 0 & 0 \\
  0 & 1 & 0 & 0 & 1 & 0 & 0 \\
  0 & 0 & 1 & 0 & 0 & 1 & 0
\end{bmatrix} \]

To find coordinates for columns \( a, b, c, d \), we will choose entries that make every row of \( A \) orthogonal to every row of \( A^* \).
### Dual of a Matroid: A concrete example

Let's consider the following matrix $A$:

$$
A = \begin{bmatrix}
1 & 0 & 0 & 1 & 1 & 1 & 1 \\
0 & 1 & 0 & 0 & 1 & -1 & -1 \\
0 & 0 & 1 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 1 \\
\end{bmatrix}
$$

Then, the dual matrix $A^*$ is:

$$
A^* = \begin{bmatrix}
-1 & -1 & 0 & 0 & 1 & 0 & 0 \\
-1 & 1 & -1 & 0 & 0 & 1 & 0 \\
-1 & 1 & 0 & -1 & 0 & 0 & 1 \\
\end{bmatrix}
$$

Here, $D$ is colored in green.
Uniform Matroid

Every uniform matroid is linear and can be represented over a finite field by a $k \times n$ matrix $A_M$ where the $A_M[i,j] = j^{i-1}$.

$$
\begin{bmatrix}
e_1 & e_2 & e_3 & \ldots & e_n \\
1 & 1 & 1 & \ldots & 1 \\
2 & 1 & 2 & 3 & \ldots & n \\
3 & 1 & 2^2 & 3^2 & \ldots & n^2 \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
k & 1 & 2^{k-1} & 3^{k-1} & \ldots & n^{k-1}
\end{bmatrix}_{k \times n}
$$

Observe that for $A_M$ to be representable over a finite field $\mathbb{F}$, we need that the determinant of any $k \times k$ submatrix of $A_M$ must not vanish over $\mathbb{F}$.

The determinant of any $k \times k$ submatrix of $A_M$ is upper bounded by $k! \times n^{k-1}$ (this follows from the Laplace expansion of determinants). Thus, choosing a field $\mathbb{F}$ of size larger than $k! \times n^{k-1}$ suffices.
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$$
\begin{bmatrix}
  e_1 & e_2 & e_3 & \cdots & e_n \\
  1 & 1 & 1 & \cdots & 1 \\
  2 & 1 & 2 & 3 & \cdots & n \\
  3 & 1 & 2^2 & 3^2 & \cdots & n^2 \\
  \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
  k & 1 & 2^{k-1} & 3^{k-1} & \cdots & n^{k-1}
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\[
\begin{bmatrix}
e_1 & e_2 & e_3 & \cdots & e_n \\
1 & 1 & 1 & \cdots & 1 \\
2 & 1 & 2 & 3 & \cdots & n \\
3 & 1 & 2^2 & 3^2 & \cdots & n^2 \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
k & 1 & 2^{k-1} & 3^{k-1} & \cdots & n^{k-1}
\end{bmatrix}_{k \times n}
\]

Observe that for $A_M$ to be representable over a finite field $\mathbb{F}$, we need that the determinant of any $k \times k$ submatrix of $A_M$ must not vanish over $\mathbb{F}$.

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Uniform Matroid

Every uniform matroid is linear and can be represented over a finite field by a $k \times n$ matrix $A_M$ where the $A_M[i, j] = j^{i-1}$.

\[
\begin{bmatrix}
e_1 & e_2 & e_3 & \cdots & e_n \\
1 & 1 & 1 & \cdots & 1 \\
2 & 1 & 2 & 3 & \cdots & n \\
3 & 1 & 2^2 & 3^2 & \cdots & n^2 \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
k & 1 & 2^{k-1} & 3^{k-1} & \cdots & n^{k-1}
\end{bmatrix}_{k \times n}
\]

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The determinant of any $k \times k$ submatrix of $A_M$ is upper bounded by $k! \times n^{k-1}$ (this follows from the Laplace expansion of determinants). Thus, choosing a field $\mathbb{F}$ of size larger than $k! \times n^{k-1}$ suffices.
Uniform Matroid: Size of the representation

So the size of the representation: \( O((k \log n) \times nk) \) bits.
Partition Matroid

A partition matroid $M = (E, \mathcal{I})$ is defined by a ground set $E$ being partitioned into (disjoint) sets $E_1, \ldots, E_\ell$ and by $\ell$ non-negative integers $k_1, \ldots, k_\ell$. A set $X \subseteq E$ is independent if and only if $|X \cap E_i| \leq k_i$ for all $i \in \{1, \ldots, \ell\}$.

Observe that a partition matroid is a direct sum of uniform matroids

$$U_{|E_1|, k_1} \oplus \cdots \oplus U_{|E_\ell|, k_\ell}$$

A uniform matroid $U_{n,k}$ is representable over a field $\mathbb{F}$ of size larger than $k! \times n^{k-1}$ and thus a partition matroid is representable.
Partition Matroid

A partition matroid $M = (E, \mathcal{I})$ is defined by a ground set $E$ being partitioned into (disjoint) sets $E_1, \ldots, E_\ell$ and by $\ell$ non-negative integers $k_1, \ldots, k_\ell$. A set $X \subseteq E$ is independent if and only if $|X \cap E_i| \leq k_i$ for all $i \in \{1, \ldots, \ell\}$. Observe that a partition matroid is a direct sum of uniform matroids

$$U|E_1|,k_1, \cdots ,U|E_\ell|,k_\ell, \quad U|E_1|,k_1 \oplus \cdots \oplus U|E_\ell|,k_\ell$$

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A partition matroid $M = (E, \mathcal{I})$ is defined by a ground set $E$ being partitioned into (disjoint) sets $E_1, \ldots, E_\ell$ and by $\ell$ non-negative integers $k_1, \ldots, k_\ell$. A set $X \subseteq E$ is independent if and only if $|X \cap E_i| \leq k_i$ for all $i \in \{1, \ldots, \ell\}$. Observe that a partition matroid is a direct sum of uniform matroids

$$U_{|E_1|, k_1} \oplus \cdots \oplus U_{|E_\ell|, k_\ell}, \quad U_{|E_1|, k_1} \oplus \cdots \oplus U_{|E_\ell|, k_\ell}$$

A uniform matroid $U_{n,k}$ is representable over a field $\mathbb{F}$ of size larger than $k! \times n^{k-1}$ and thus a partition matroid is representable.
The graphic matroid is representable over any field of size at least 2.

Consider the matrix $A_M$ with a row for each vertex $i \in V(G)$ and a column for each edge $e = ij \in E(G)$. In the column corresponding to $e = ij$, all entries are 0, except for a 1 in $i$ or $j$ (arbitrarily) and a $-1$ in the other.

$$
\begin{bmatrix}
e_1 & e_2 & e_3 & \cdots & e_m \\
1 & 1 & 0 & 1 & \cdots & 0 \\
2 & 0 & 0 & 0 & \cdots & 1 \\
3 & -1 & 1 & 0 & \cdots & 0 \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
n & 0 & -1 & -1 & \cdots & -1
\end{bmatrix} \quad n \times |E(G)|
$$

This is a representation over reals.
The graphic matroid is representable over any field of size at least 2.
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\[
\begin{bmatrix}
1 & 1 & 0 & 1 & \cdots & 0 \\
2 & 0 & 0 & 0 & \cdots & 1 \\
3 & -1 & 1 & 0 & \cdots & 0 \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
n & 0 & -1 & -1 & \cdots & -1
\end{bmatrix}
\]

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$$
\begin{bmatrix}
    e_1 & e_2 & e_3 & \cdots & e_m \\
    1 & 1 & 0 & 1 & \cdots & 0 \\
    2 & 0 & 0 & 0 & \cdots & 1 \\
    3 & -1 & 1 & 0 & \cdots & 0 \\
    \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
    n & 0 & -1 & -1 & \cdots & -1 \\
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\[
\begin{bmatrix}
e_1 & e_2 & e_3 & \cdots & e_m \\
1 & 1 & 0 & 1 & \cdots & 0 \\
2 & 0 & 0 & 0 & \cdots & 1 \\
3 & 1^{-1} & 1 & 0 & \cdots & 0 \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
n & 0 & 1^{-1} & 1^{-1} & \cdots & 1^{-1}
\end{bmatrix}_{n \times |E(G)|}
\]

To obtain a representation over a field $\mathbb{F}$, one simply needs to take the representation given above over reals and simply replace all $-1$ by the additive inverse of 1.
Transversal Matroid

For the bipartite graph with partition \( A \) and \( B \), form an incidence matrix \( A_M \) as follows. Label the rows by vertices of \( B \) and the columns by the vertices of \( A_M \) and define:

\[
a_{ij} = \begin{cases} 
z_{ij} & \text{if there is an edge between } a_i \text{ and } b_j, \\
0 & \text{otherwise.}
\end{cases}
\]

where \( z_{ij} \) are in-determinants. Think of them as independent variables.

\[
T = \begin{bmatrix}
a_1 & a_2 & \cdots & a_j & \cdots & a_\ell \\
b_1 & z_{11} & z_{12} & \cdots & z_{1j} & \cdots & z_{1\ell} \\
& \vdots & \vdots & \ddots & \vdots & \ddots & \vdots \\
& \vdots & \vdots & \ddots & z_{ij} & \cdots & z_{i\ell} \\
& \vdots & \vdots & \ddots & \vdots & \ddots & \vdots \\
b_n & z_{n1} & z_{n2} & \cdots & z_{nj} & \cdots & z_{n\ell}
\end{bmatrix}
\]
Transversal Matroid

For the bipartite graph with partition $A$ and $B$, form an incidence matrix $A_M$ as follows. Label the rows by vertices of $B$ and the columns by the vertices of $A_M$ and define:

$$a_{ij} = \begin{cases} 
z_{ij} & \text{if there is an edge between } a_i \text{ and } b_j, \\
0 & \text{otherwise}. 
\end{cases}$$

where $z_{ij}$ are in-determinants. Think of them as independent variables.

$$T = \begin{bmatrix}
a_1 & a_2 & \cdots & a_j & \cdots & a_\ell \\
b_1 & z_{11} & z_{12} & \cdots & z_{1j} & \cdots & z_{1\ell} \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
b_i & z_{i1} & z_{i2} & \cdots & z_{ij} & \cdots & z_{i\ell} \\
\vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\
b_n & z_{n1} & z_{n2} & \cdots & z_{nj} & \cdots & z_{n\ell}
\end{bmatrix}$$
Example of the Construction

\[
\begin{bmatrix}
  a_1 & a_2 & a_3 & a_4 & a_5 & a_6 \\
  b_1 & z_{11} & z_{12} & z_{13} & 0 & z_{15} & 0 \\
  b_2 & 0 & z_{22} & z_{23} & z_{24} & z_{25} & 0 \\
  b_3 & 0 & 0 & 0 & 0 & z_{35} & z_{36}
\end{bmatrix}
\]
Example of the Construction

\[
\begin{pmatrix}
    a_1 & a_2 & a_3 & a_4 & a_5 & a_6 \\
    b_1 & z_{11} & z_{12} & z_{13} & 0 & z_{15} & 0 \\
    b_2 & 0 & z_{22} & z_{23} & z_{24} & z_{25} & 0 \\
    b_3 & 0 & 0 & 0 & 0 & z_{35} & z_{36}
\end{pmatrix}
\]
Permutation expansion of Determinants

**Theorem:** Let

\[ A = (a_{ij})_{n \times n} \]

be a \( n \times n \) matrix with entries in \( \mathbb{F} \). Then

\[ \det(A) = \sum_{\pi \in S_n} \text{sgn}(\pi) \prod_{i=1}^{n} a_{i\pi(i)}. \]
Proof that Transversal Matroid is Representable over $F[\vec{z}]$

Forward direction: (Board for Picture)

- Suppose some subset $X = \{a_1, \ldots, a_q\}$ is independent.
- Then there is a matching that saturates $X$. Let $Y = \{b_1, b_2, \ldots, b_q\}$ be the endpoints of this matching and $a_i b_i$ are the matching edges.
- Consider $T[Y, X]$ – a submatrix with rows in $Y$ and columns in $X$. Consider the determinant of $T[Y, X]$ then we have a term

$$\prod_{i=1}^{q} z_{ii}$$

which can not be cancelled by any other term! So these columns are linearly independent.
Proof that Transversal Matroid is Representable over $F[\vec{z}]$

Reverse direction: (Board for Picture)

- Suppose some subset $X = \{a_1, \ldots, a_q\}$ of columns is independent in $T$.
- Then there is a submatrix of $T[\star, X]$ that has non-zero determinant – say $T[Y, X]$.
- Consider the determinant of $T[Y, X]$

  $$0 \neq \det(T[Y, X]) = \sum_{\pi \in S(Y)} \sgn(\pi) \prod_{i=1}^{q} z_{i\pi(i)}.$$  

  - This implies that we have a term

    $$\prod_{i=1}^{q} z_{i\pi(i)} \neq 0$$

    and this gives us that there is a matching that saturates $X$ in and hence $X$ is independent.
Proof that Transversal Matroid is Representable over $F[\vec{z}]$ 

Reverse direction: (Board for Picture)

- This implies that we have a term
  \[
  \prod_{i=1}^{q} z_{i\pi(i)} \neq 0
  \]
  and this gives us that there is a matching that saturates $X$ in and hence $X$ is independent.

- For this direction we do not use $z_{ij}$, the very fact that $X$ forms independent set of column is enough to argue that there is a matching that saturates $X$. 

-
Removing $z_{ij}$

To remove the $z_{ij}$ we do the following.

Uniformly at random assign $z_{ij}$ from values in finite field $\mathbb{F}$ of size $P$.

What should be the upper bound on $P$? What is the probability that the randomly obtained $T$ is a representation matrix for the transversal matroid.
Removing $z_{ij}$

To remove the $z_{ij}$ we do the following.

*Uniformly at random assign $z_{ij}$ from values in finite field $\mathbb{F}$ of size $P$.*

What should be the upper bound on $P$? What is the probability that the randomly obtained $T$ is a representation matrix for the transversal matroid.
**Theorem:** Let $p(x_1, x_2, \ldots, x_n)$ be a non-zero polynomial of degree $d$ over some field $\mathbb{F}$ and let $S$ be an $N$ element subset of $\mathbb{F}$. If each $x_i$ is independently assigned a value from $S$ with uniform probability, then $p(x_1, x_2, \ldots, x_n) = 0$ with probability at most $\frac{d}{N}$.

- We think $\det(T[Y, X])$ as polynomial in $z_{ij}$’s of degree at most $n = |A|$.
- Probability that $\det(T[Y, X]) = 0$ is less than $\frac{n}{P}$. There are at most $2^n$ independent sets in $A$ and thus by union bound probability that not all of them are independent in the matroid represented by $T$ is at most $\frac{2^n n}{P}$.
Using Zippel-Schwartz Lemma

- We think $\det(T[Y,X])$ as polynomial in $z_{ij}$’s of degree at most $n = |A|$.
- Probability that $\det(T[Y,X]) = 0$ is less than $\frac{n}{\mathbb{P}}$. There are at most $2^n$ independent sets in $A$ and thus by union bound probability that not all of them are independent in the matroid represented by $T$ is at most $\frac{2^n n}{\mathbb{P}}$.
- Thus probability that $T$ is the representation is at least $1 - \frac{2^n n}{\mathbb{P}}$. Take $\mathbb{P}$ to be some field with at least $2^n n 2^n$ elements :-).
- Size of this representation with be like $n^{O(1)}$ bits!
Strict Gammoids: Idea of Proof for Representation

- Show that this is a dual of a transversal matroid.
- Transversal Matroids are representable.
- Dual of a Representable Matroid is Representable.
- Conclude that Strict Gammoids are representable.
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Transversal Matroid and Strict Gammoids

Let $D = (V, A)$ be a directed graph and $S \subseteq V = \{v_1, \ldots, v_n\}$. Let $M = (E, \mathcal{I})$ be the corresponding strict gammoid.

- For transversal matroid we need a bipartite graph. Let $G^*$ be a bipartite graph with bipartition $U$ and $W$. $U = \{u_i \mid v_i \in V\}$ and $W = \{w_i \mid w_i \in V \setminus S\}$.
- For each edge $v_i \in V \setminus S$ we have an edge $u_iw_i$ and for every edge $\overrightarrow{v_iv_j}$ there is an edge $u_iw_j$.

Want to show:

Let $V_0 \in \mathcal{I}$ and $|V_0| = |S|$. $V_0$ is linked to $S$ if and only if there is a matching that saturates $U \setminus U_0$ in $G^*$. 
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Proof Forward Direction

Assume first that there are $|S|$ disjoint paths going from $S$ to $V_0$. Consider the matching as follows:

- $w_i \in W$ is matched to $u_j$ if $\overrightarrow{v_jv_i}$ is an arc on some path; else
- $w_i$ is matched to $u_i$.

- This means that $u_i$ is matched if one of the paths reaches $v_i$ and continues further on, or if none of the paths reaches $v_i$.
- Thus the unmatched $u_i$ corresponds to the end points of the paths, as required.
Proof Reverse Direction

- There is a matching that saturates $U \setminus U_0$.
- Now for every vertex corresponding to $S$ that is not in $V_0$, grow maximal path using matching edges.
- This corresponds to paths from $S$ to $V_0$ in $D$. 
Find the representation of strict-gammoids over a large field $\mathbb{F}$.

Delete the columns corresponding to $V \setminus T$.

Now reduce the size of the numbers using modular arithmetic over a prime with at most $O((|S| + |T|)^{O(1)})$ bits.

One can show: that the size of representation now will be $O((|S| + |T|)^{O(1)})$. 
Final Slide

Thank You!
Any Questions?