

Topics in Computational Social Choice Theory

Lecture 4: Existence of envy-free cake divisions

Nidhi Rathi



Last Lecture: Introduction to Cake Cutting

- The resource: **Cake [0,1]** (heterogeneous and divisible)
- Set of **agents**: {1,2, ..., n}
- **Piece** of a cake: finite union of subintervals of [0,1]
- Valuation function v_i : Agent *i* values piece X at $v_i(X) \ge 0$



Fairness Notions

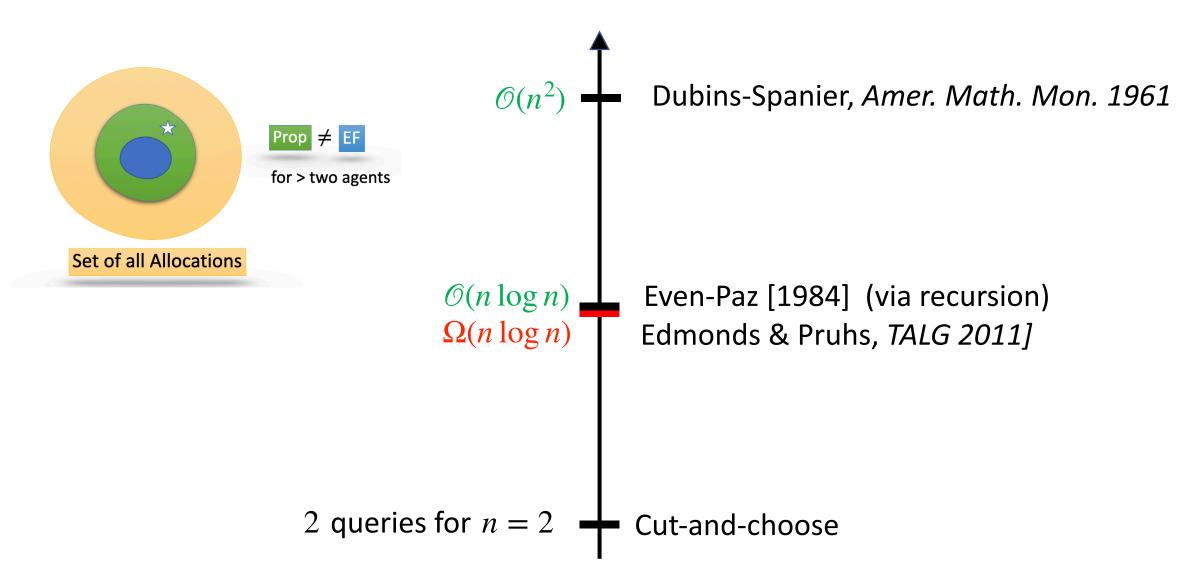
Allocation:

A partition $A = (A_1, A_2, ..., A_n)$ of the cake [0,1] where piece A_i belongs to agent i



- **<u>Proportionality</u>**: for each agent $i \in [n]$, we have $v_i(A_i) \ge 1/n$ [Steinhaus, 1948]
- Envy-freeness: for every pair $i, j \in [n]$ of agents, we have $v_i(A_i) \ge v_i(A_j)$ [Foley 1967]

Query Complexity of Proportionality



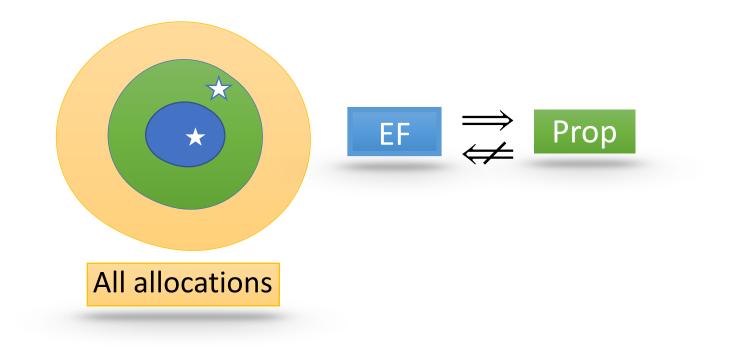
Existence of Envy-free Cake Divisions

- Computing an envy-free cake division:
- Cut-and-choose: between two agents using 2 queries
- **Selfridge-Conway:** among three agents using 8 queries

non-contiguous pieces

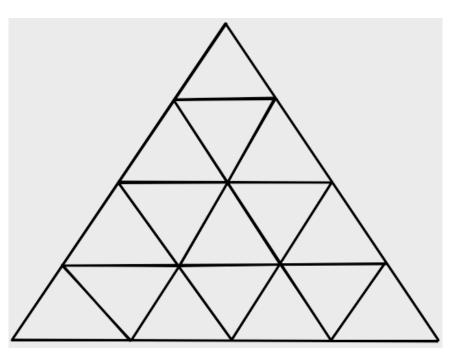
What about $n \ge 4$ agents?

Existence of Envy-free Cake Divisions



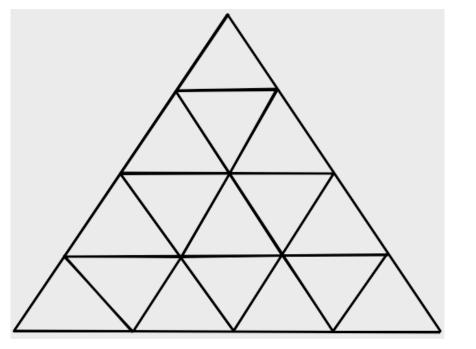
Stromquist [1980], Su [1999]connected piecesEnvy-free cake division exist for any number of agents

A beautiful lemma that, on the face of it, has nothing to do with cake division



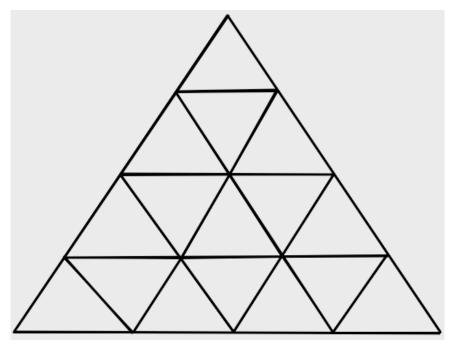


1) A **triangle** that is *subdivided* into smaller triangles (Formal terms: simplex and its triangulation)



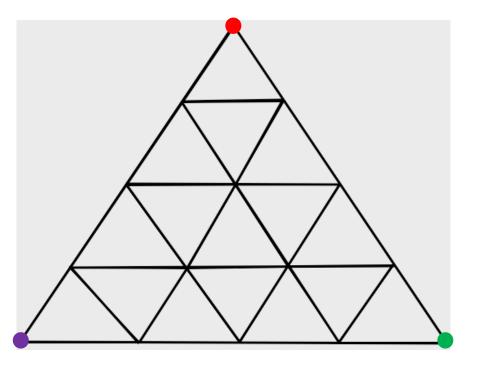


1) A **triangle** that is *subdivided* into baby triangles (Formal terms: simplex and its triangulation)



Ingredients:

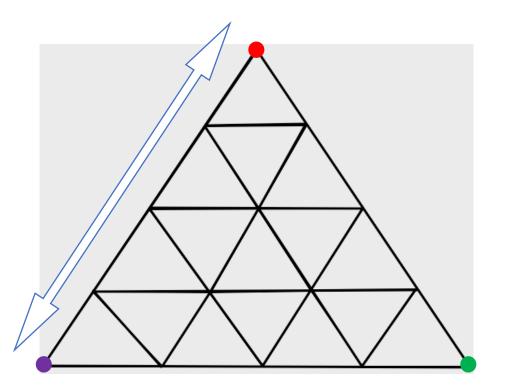
- 1) A **triangle** that is *subdivided* into baby triangles (Formal terms: simplex and its triangulation)
- 2) Sperner coloring
 - *Main vertices* have *distinct* colors



Ingredients:

1) A **triangle** that is *subdivided* into baby triangles (Formal terms: simplex and its triangulation)

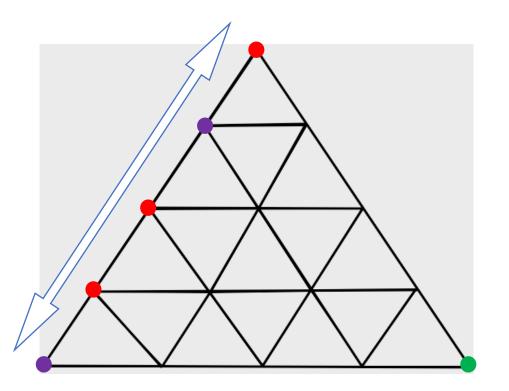
- *Main vertices* have *distinct* colors
- *Boundary vertices inherit* colors of the adjacent main vertices



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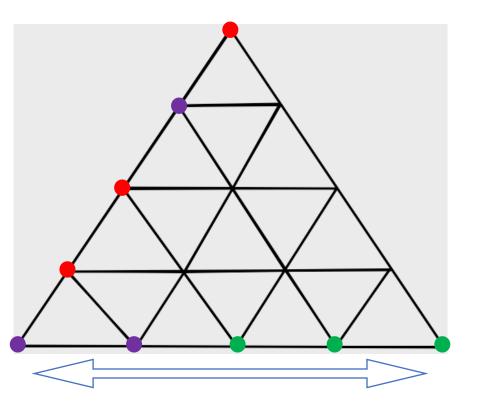
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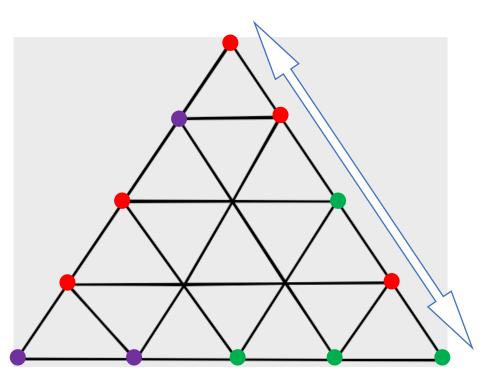
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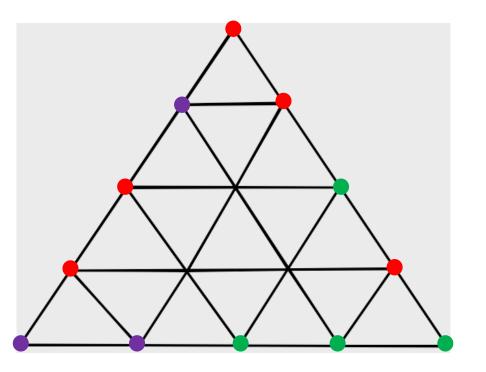
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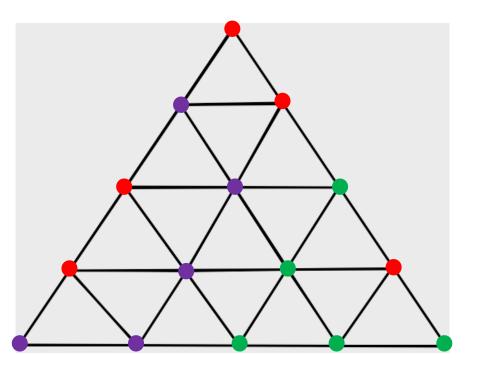
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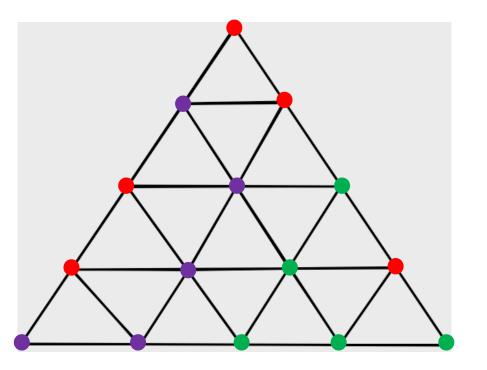


Ingredients:

1) A **triangle** that is *subdivided* into baby triangles (Formal terms: simplex and its triangulation)

2) Sperner labeling

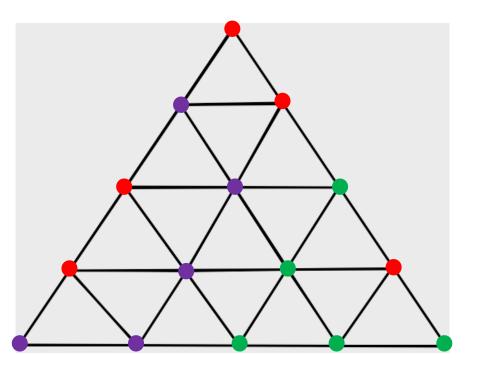
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Ingredients:

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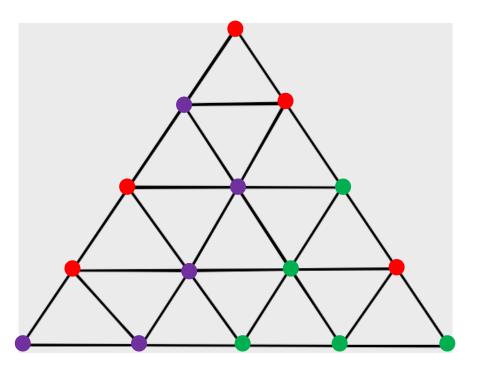


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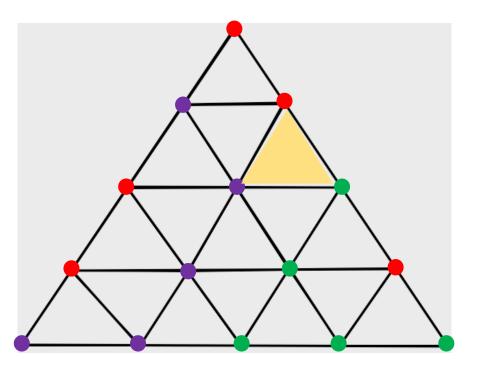


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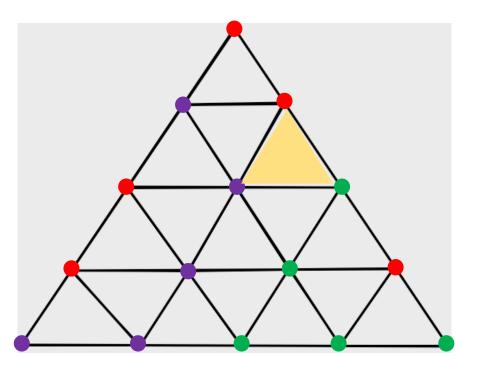


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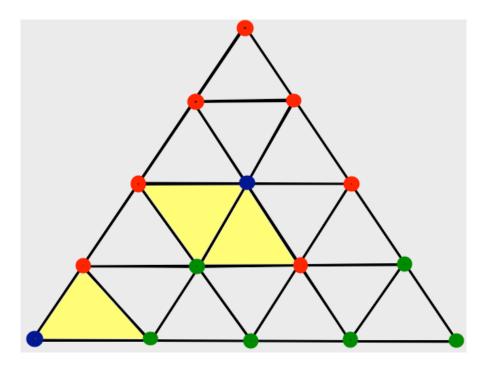
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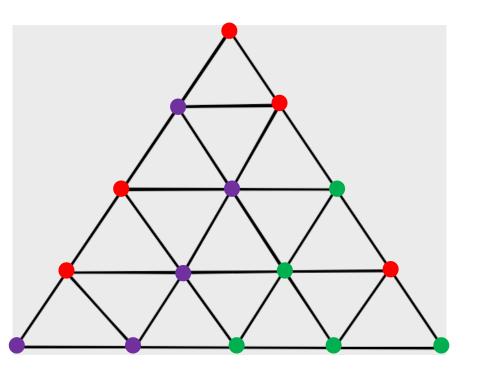
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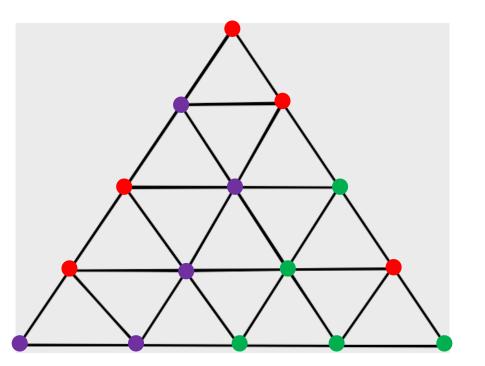
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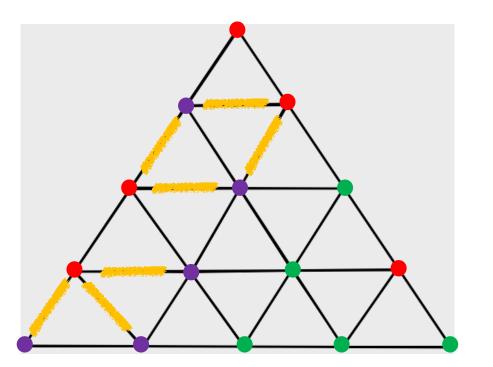
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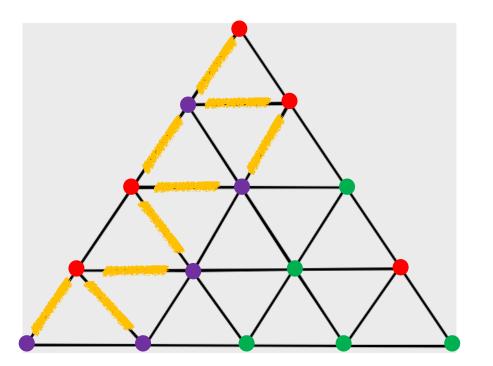
- Entire main triangle: HOUSE
- Baby triangles: ROOMS
- • DOOR



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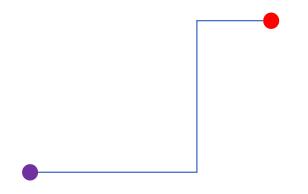


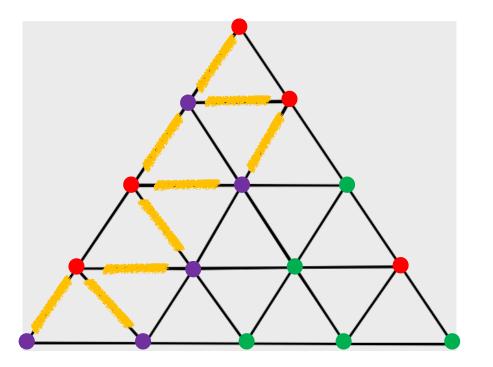
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Observation 1:

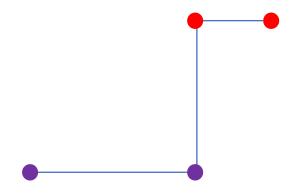
Number of doors on the boundary is **ODD**

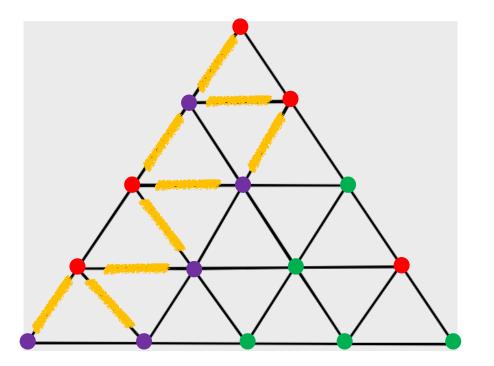




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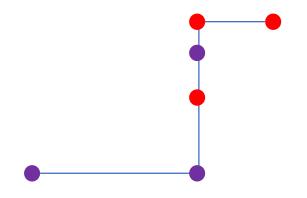
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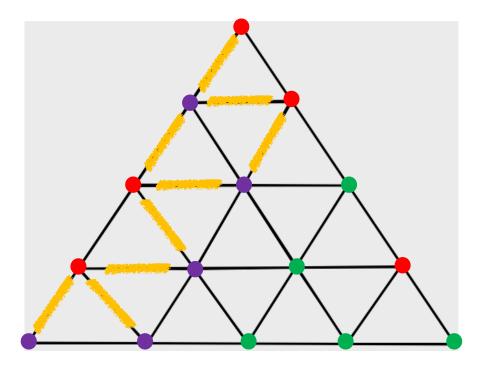




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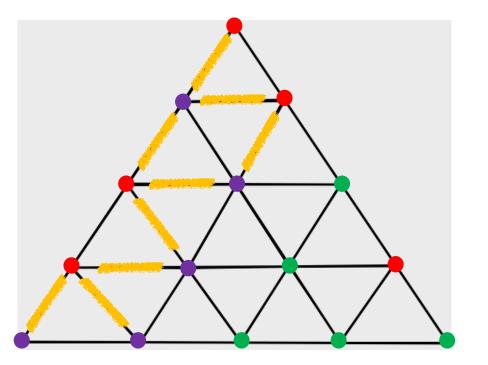
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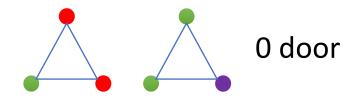
Observation 2:

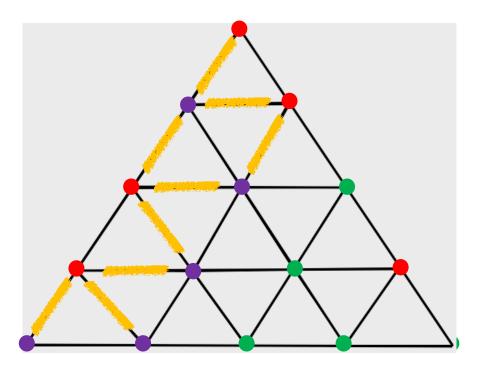
A room can have 0, 1, or 2 doors



Observation 2:

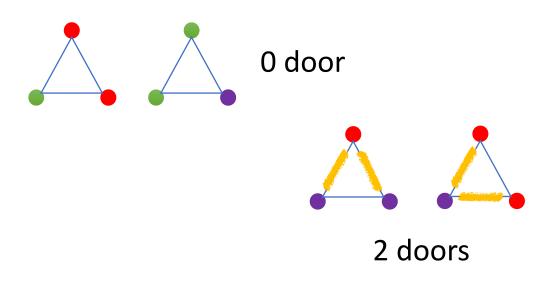
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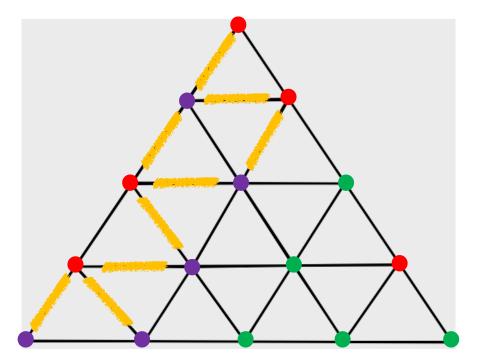




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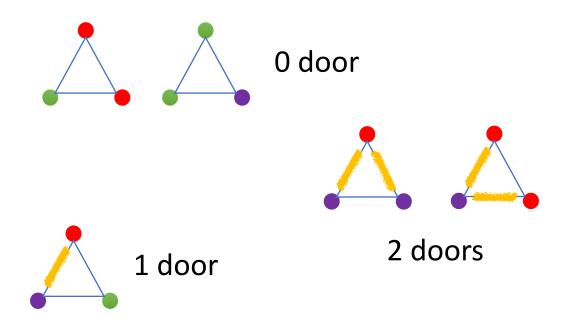
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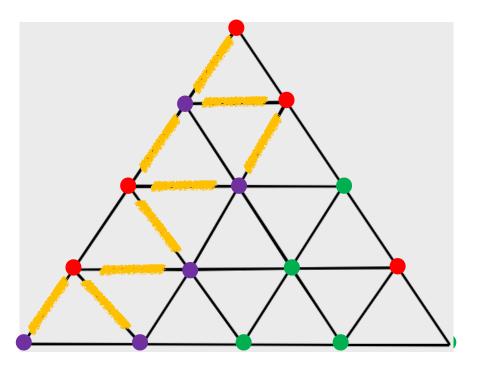




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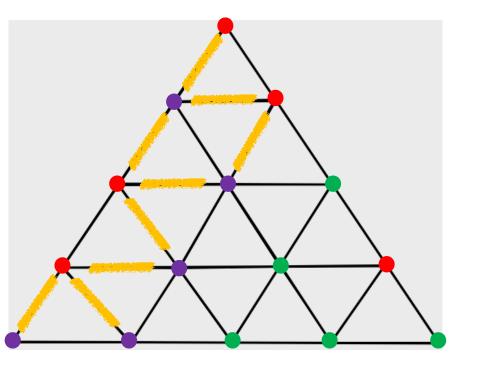
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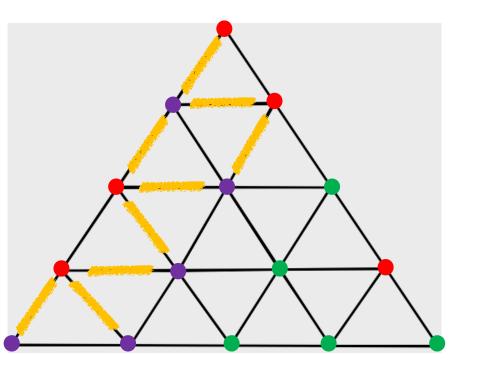
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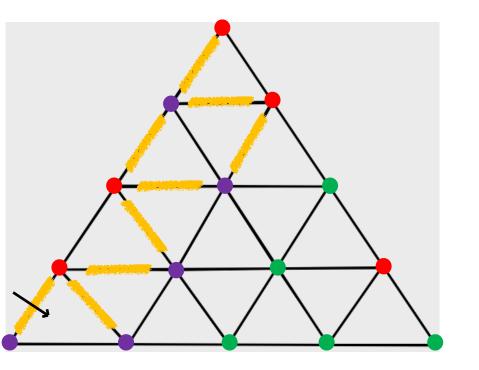
So, we enter the house through a door!



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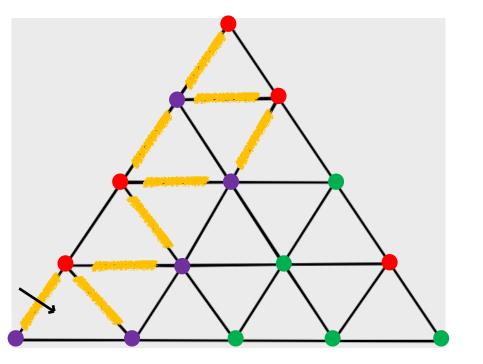
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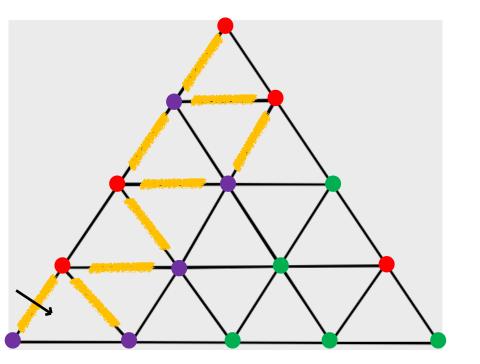
The room we entered can have either 1 or 2 doors



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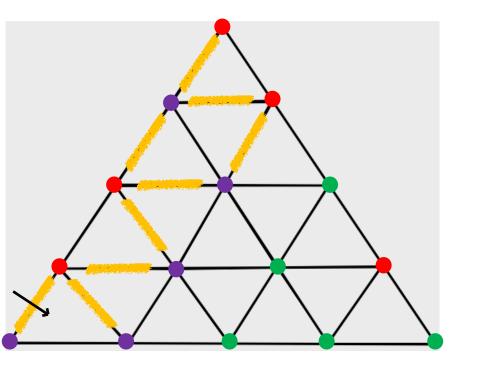
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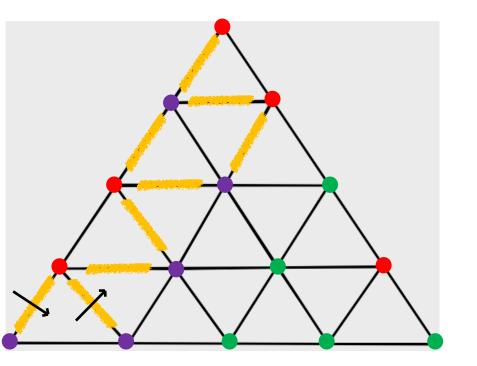
- 1 door : sperner solution
- 2 doors: leave the room using the other door and enter a new room



Enter the house through a door.

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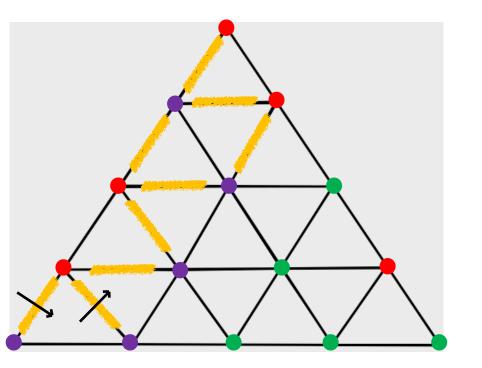


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Keep walking!

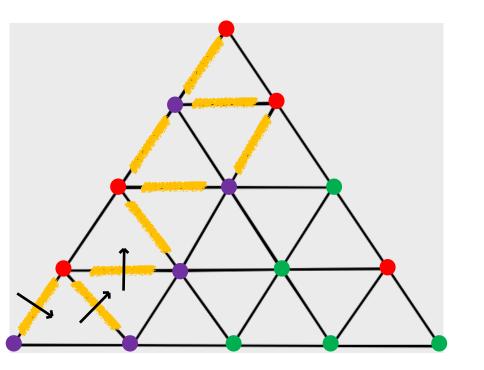


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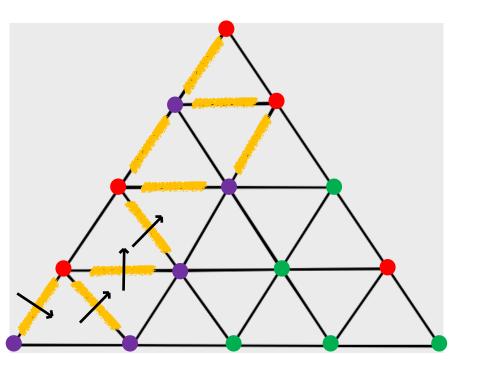


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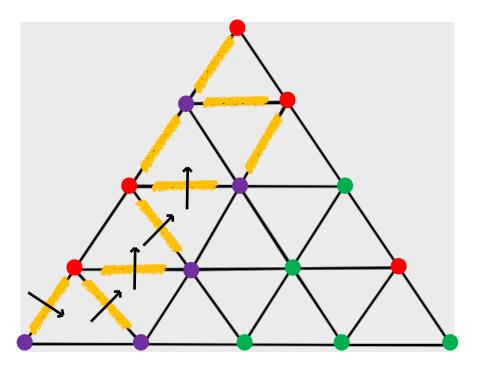
Enter the house through a door.

The room we entered can have either 1 or 2 doors

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Keep walking!

• reach a fully colored baby triangle



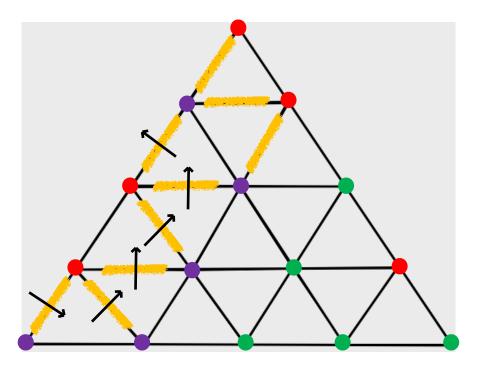
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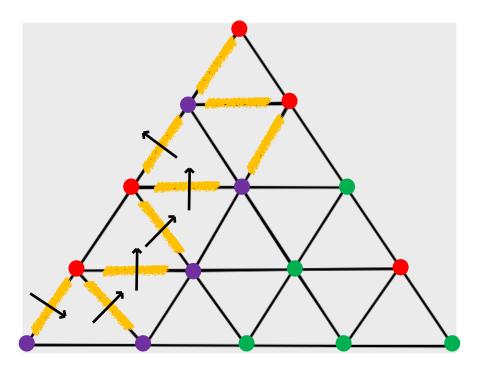
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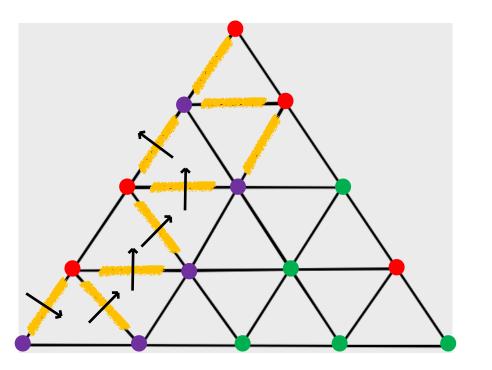
- reach a fully colored baby triangle
- thrown out of the house



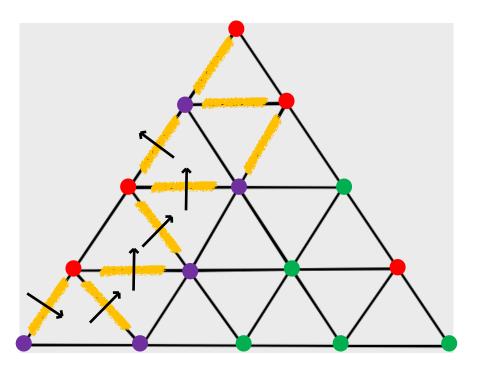
Thrown out?



Thrown out?
Cannot happen from (since no doors)

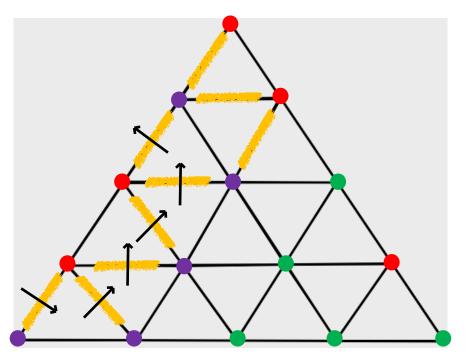


- Thrown out?
- Cannot happen from
 (since no doors)
- Entry and exit doors are paired up



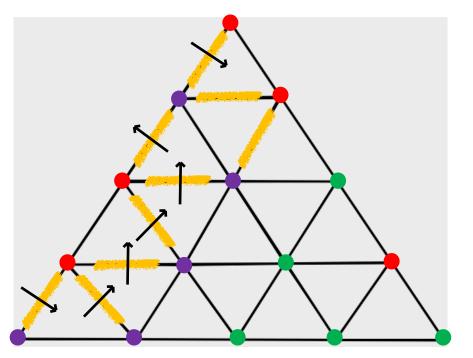
Thrown out?

- Cannot happen from
 (since no doors)
- Entry and exit doors are paired up
- There exists odd number of doors on the boundary.
 ⇒ we can enter again from another door!



Thrown out?

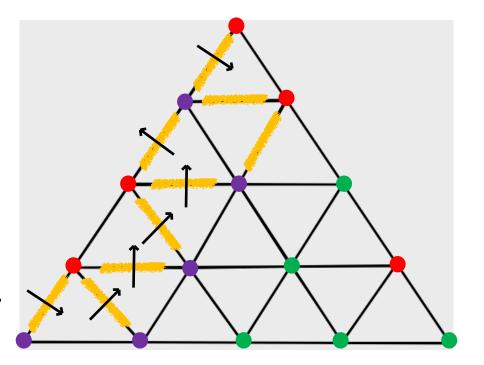
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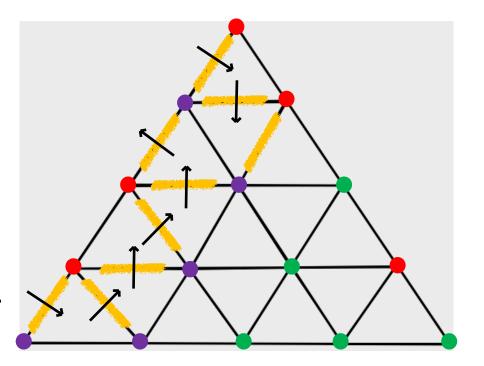
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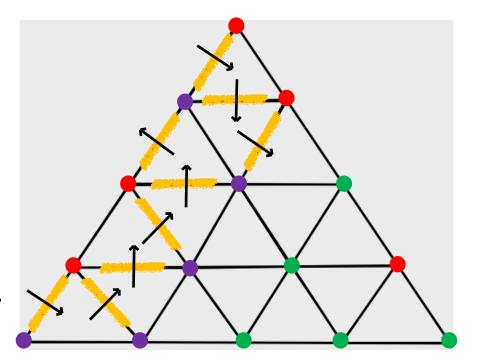
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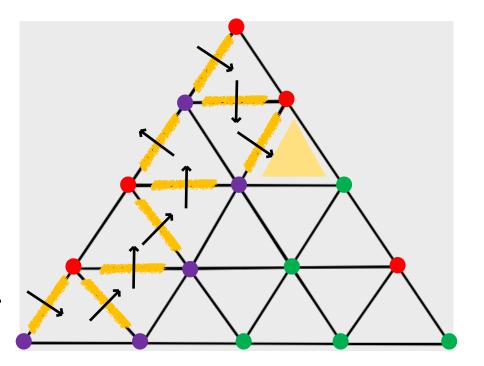
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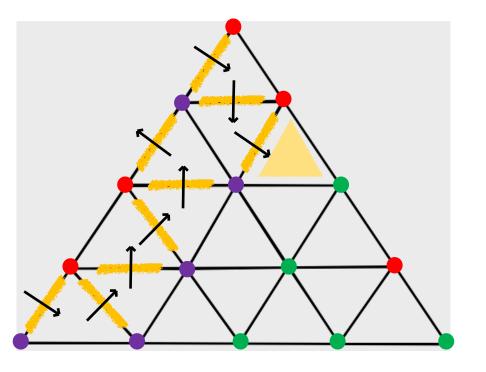


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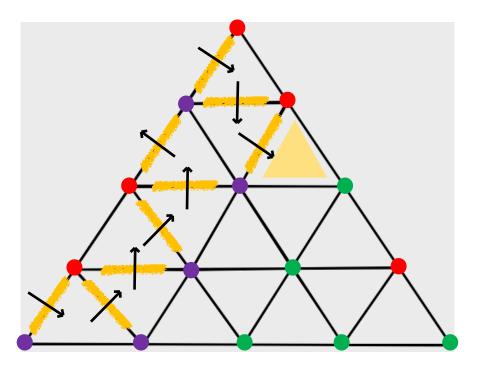




Think: Why cannot such walks cycle back on themselves?

- The number of rooms = finite
 - \implies the walk **terminates**

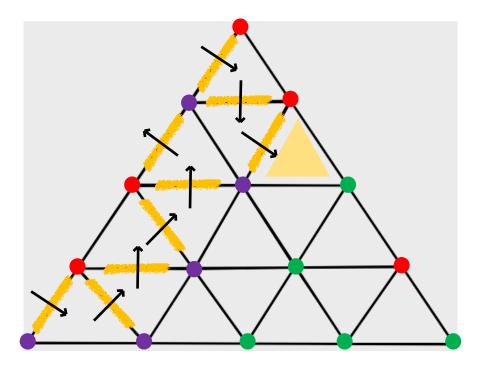
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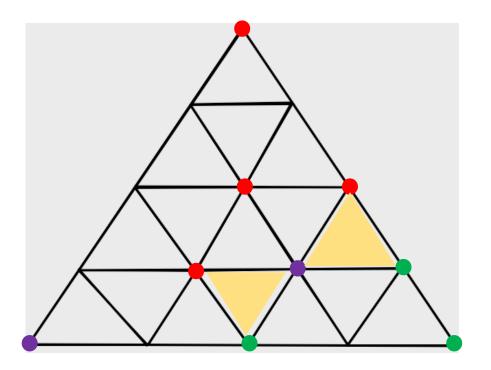
- The number of rooms = finite
 ⇒ the walk terminates
- ∃ at least one walk that will take us to a fully colored sperner solution

Think:

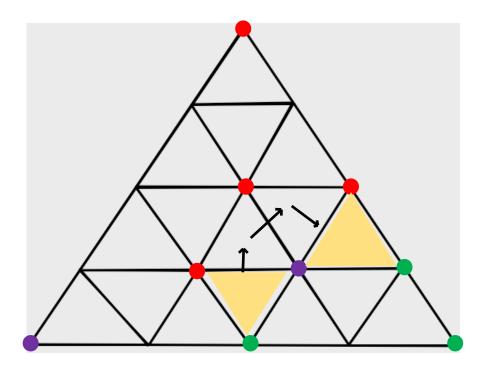
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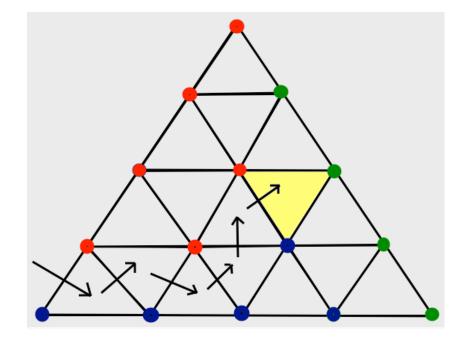
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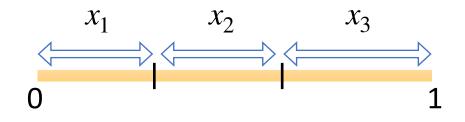


Holds true for any dimension

Forest Simmons, popularized by Francis Su [1999]

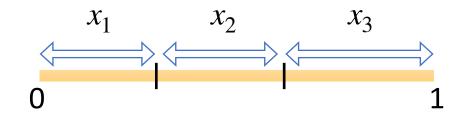
- The resource: cake [0,1] and n agents
- An allocation $(X_1, ..., X_n)$ is *envy-free* if $v_i(X_i) \ge v_i(X_j)$ for all i, j

- The resource: **cake** [0,1] and three **agents**
- An allocation (X_1, X_2, X_3) is *envy-free* if $v_i(X_i) \ge v_i(X_i)$ for all i, j



 (x_1, x_2, x_3) : a cut

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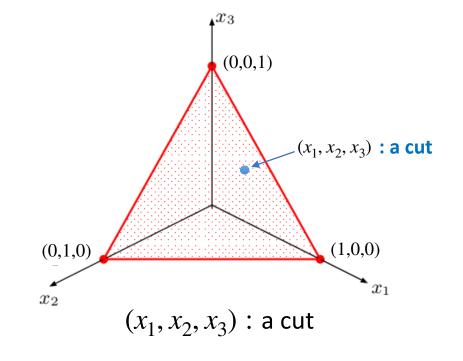


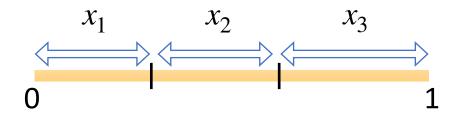
 $x_1 + x_2 + x_3 = 1$ and all $x_i \ge 0$

Space of all possible cuts

 (x_1, x_2, x_3) : a cut

- The resource: cake [0,1] and three agents
- An allocation (X_1, X_2, X_3) is *envy-free* if $v_i(X_i) \ge v_i(X_i)$ for all i, j

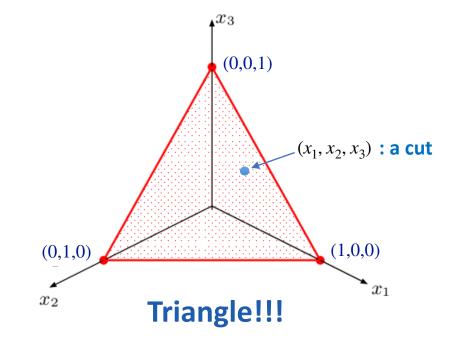


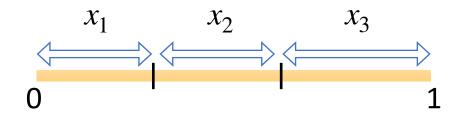


 $x_1 + x_2 + x_3 = 1 \text{ and all } x_i \ge 0$ (2-simplex)

Space of all possible cuts

- The resource: cake [0,1] and three agents
- An allocation (X_1, X_2, X_3) is *envy-free* if $v_i(X_i) \ge v_i(X_i)$ for all i, j

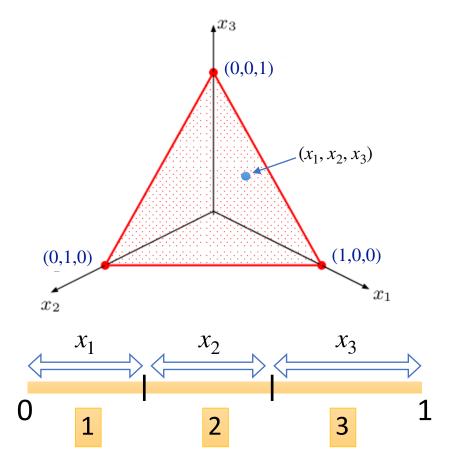




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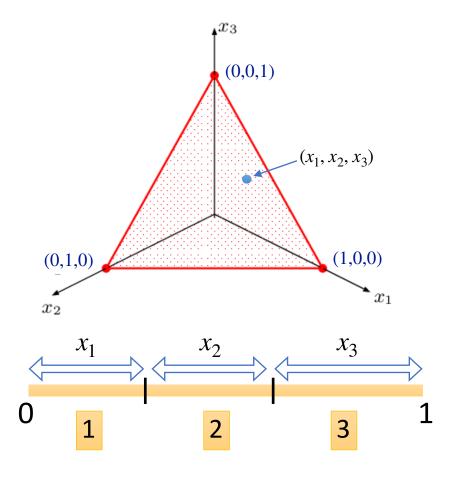
Space of all possible cuts

Assumptions on preferences/valuations



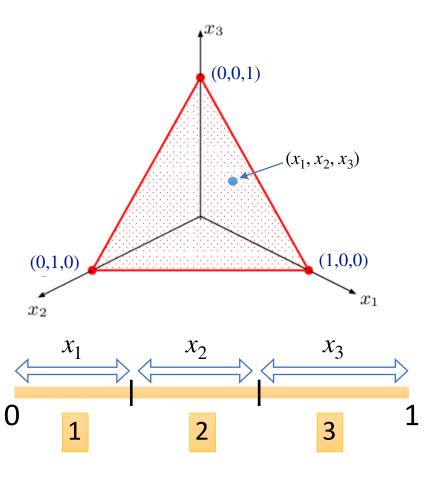
Assumptions on preferences/valuations

 Given any cut (x₁, x₂, x₃), each agent can point to its favorite piece



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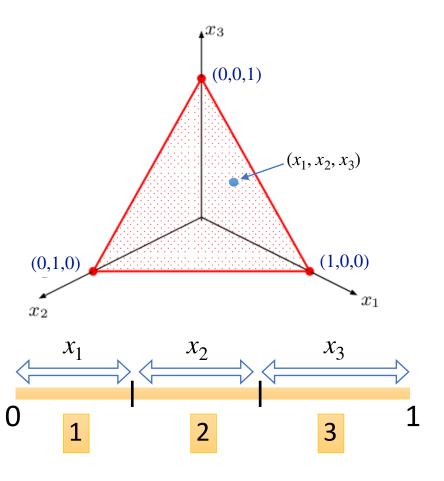
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Goal: to invoke Sperner's lemma somehow

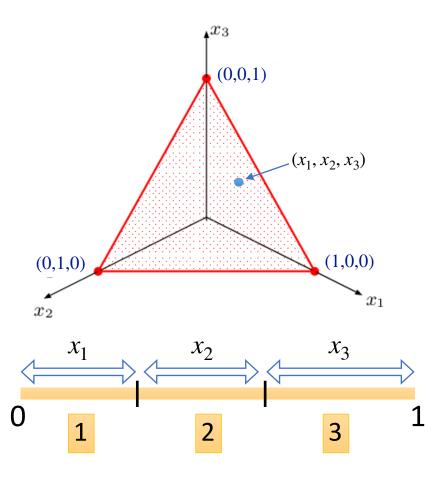


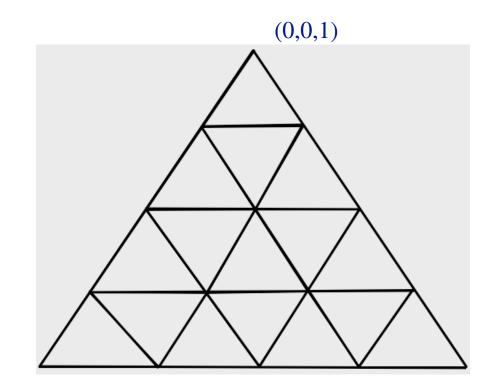
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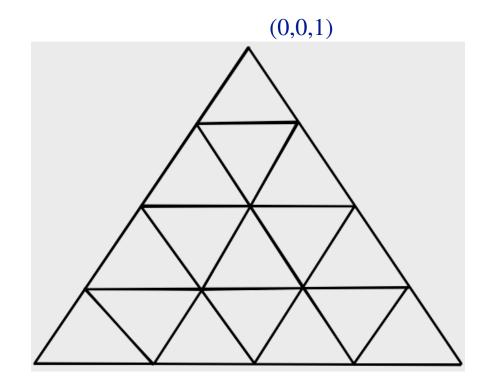
Set of agents: $\{A, B, C\}$





(0,1,0)

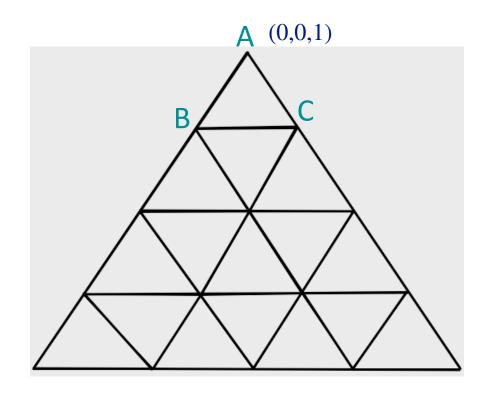
(1,0,0)



Assign ownerships to each vertex such that each baby triangle consists of all three owners {A, B, C}.

(0,1,0)

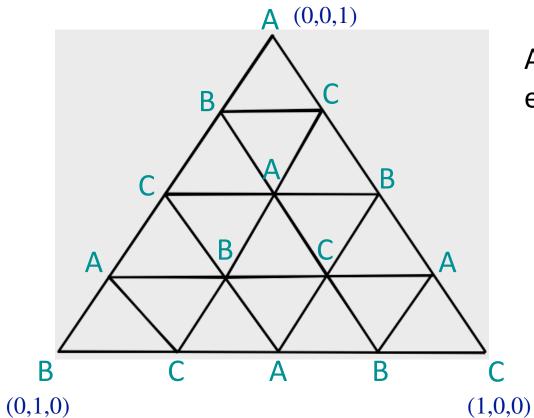
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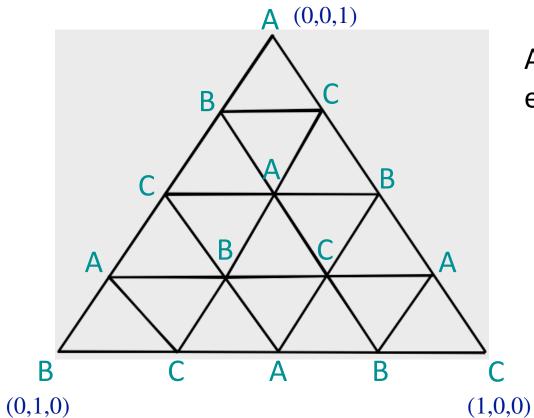
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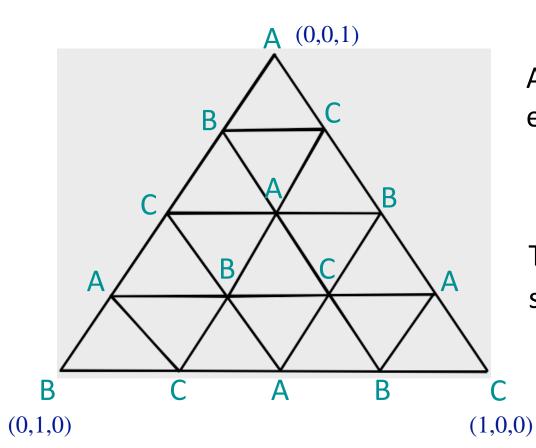
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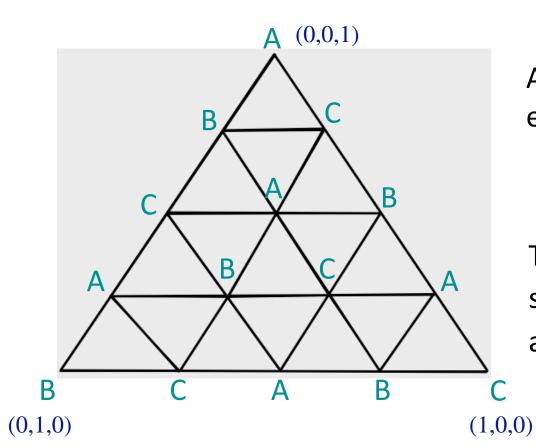


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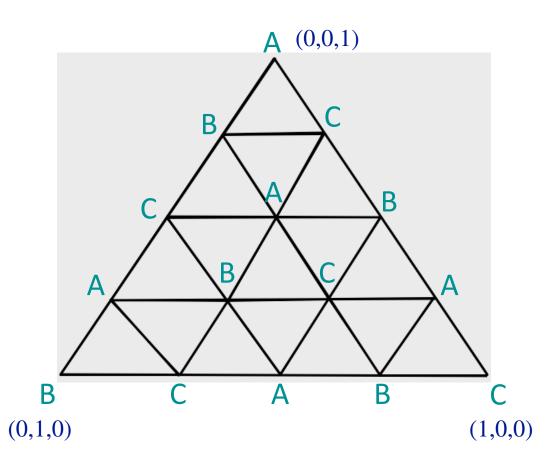
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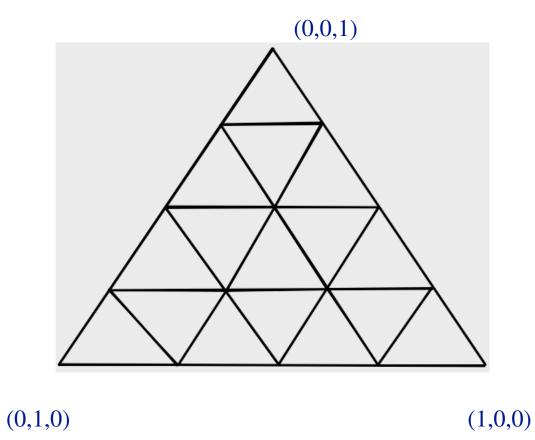
To generate a **Sperner coloring**, we go to a vertex, say some (x_1, x_2, x_3) , and

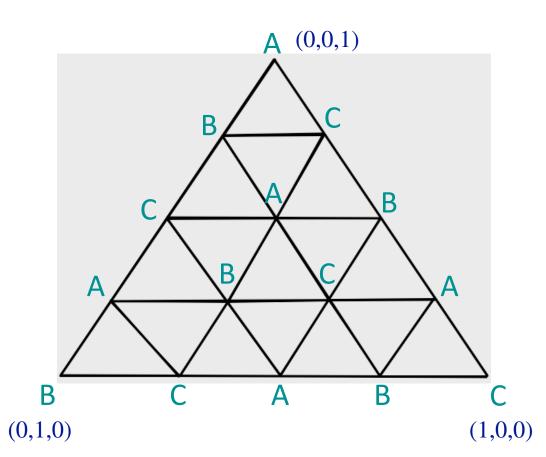


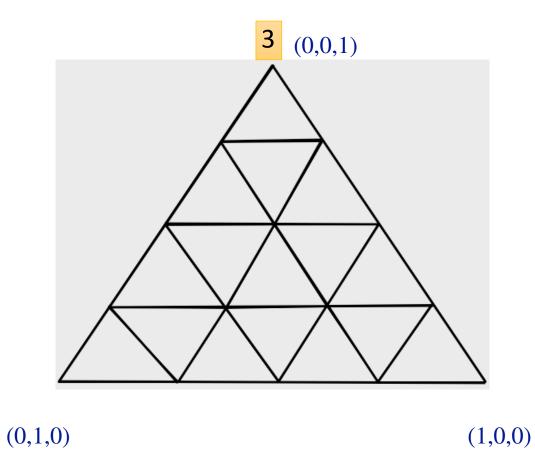
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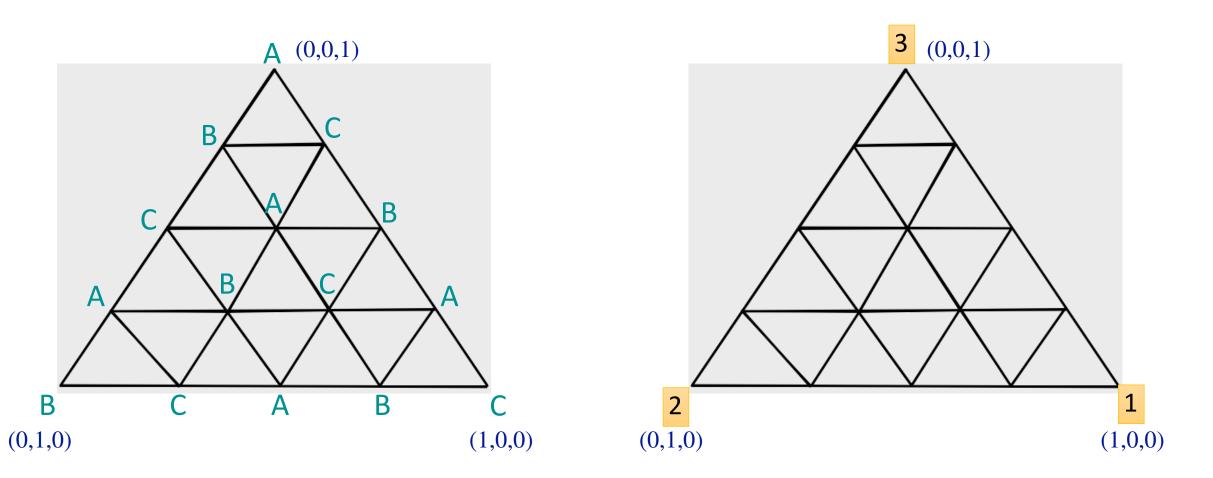
To generate a **Sperner coloring**, we go to a vertex, say some (x_1, x_2, x_3) , and ask its owner agent her most **favorite** piece in this cut

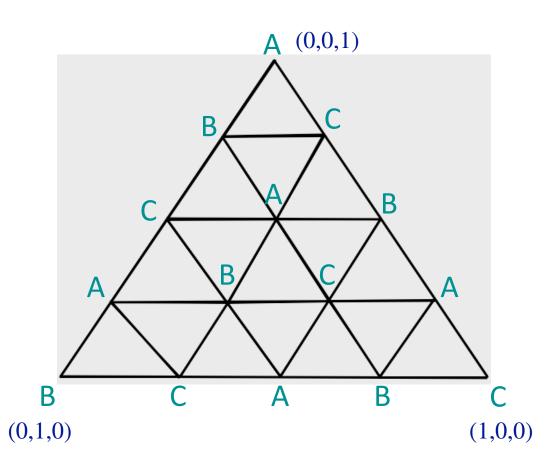


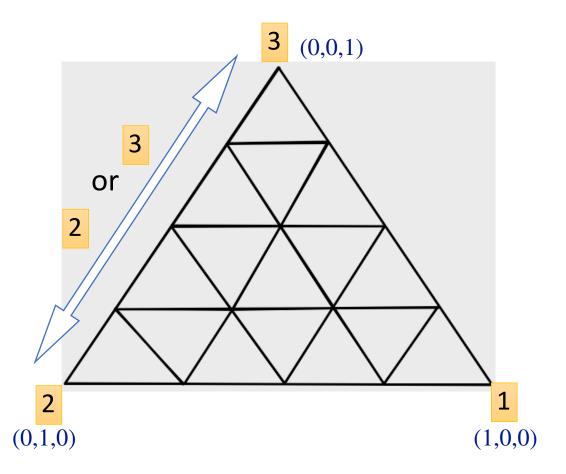


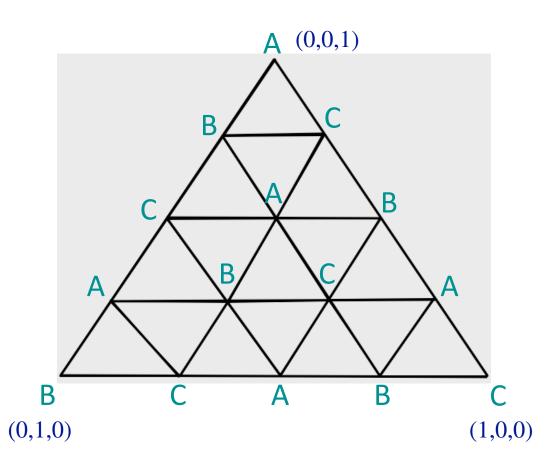


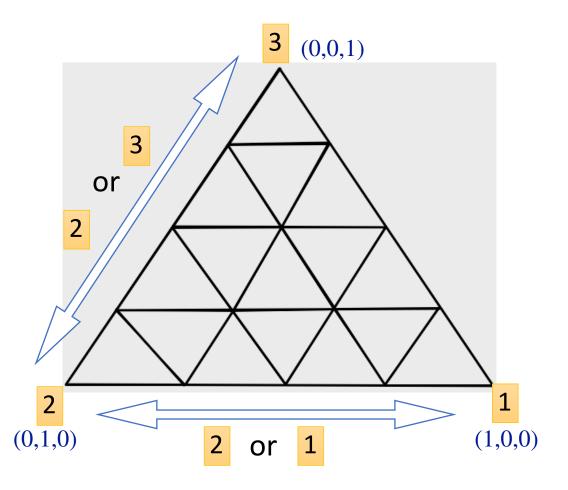


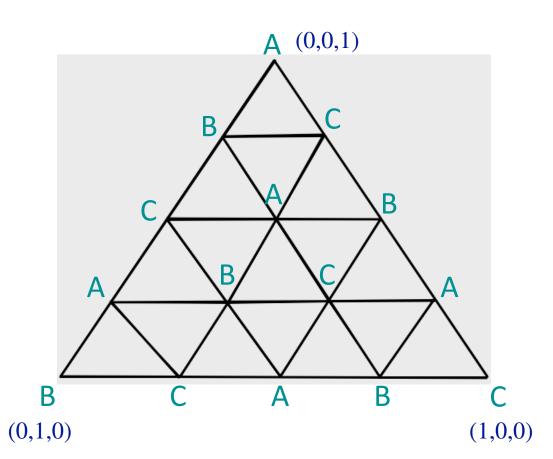


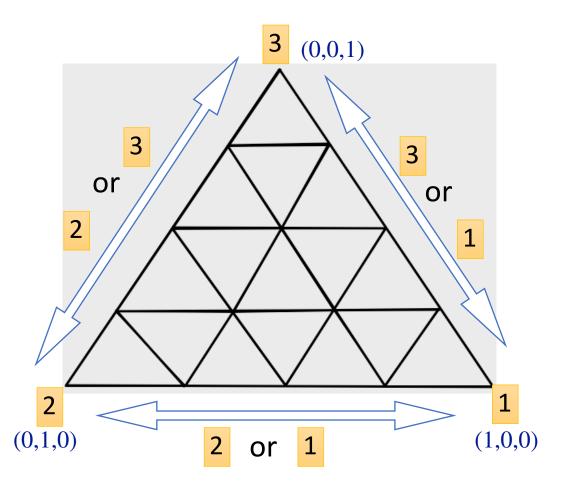


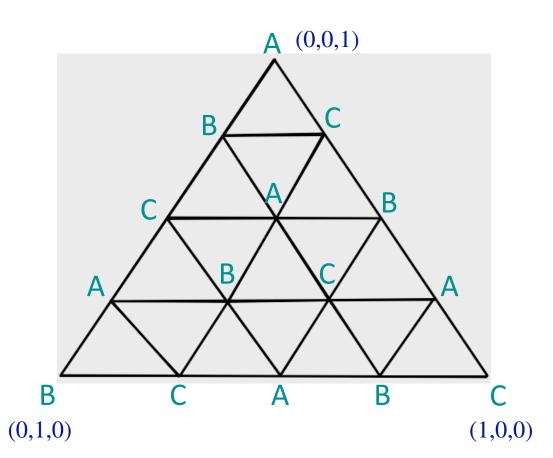


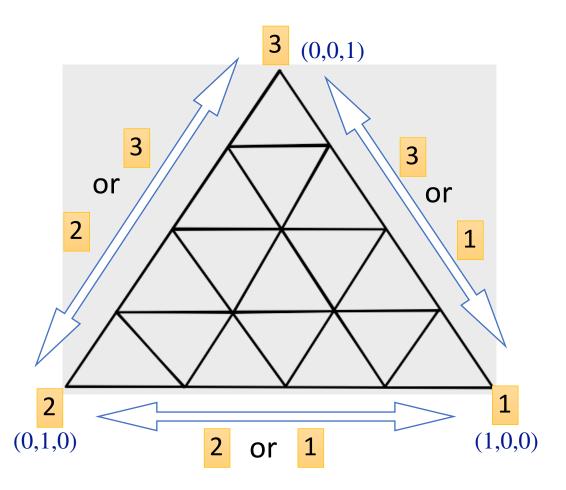








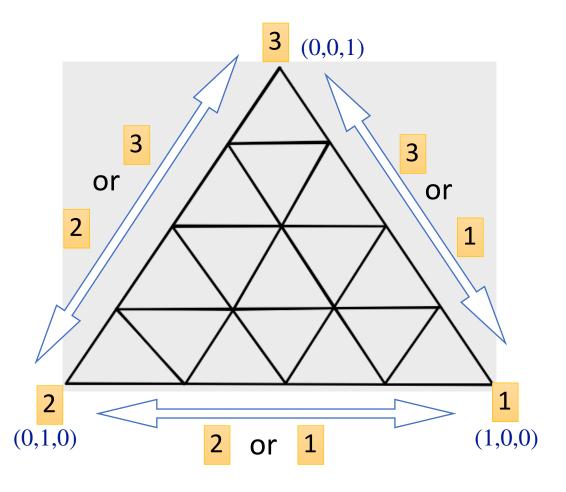




Ownership labeling

Sperner coloring

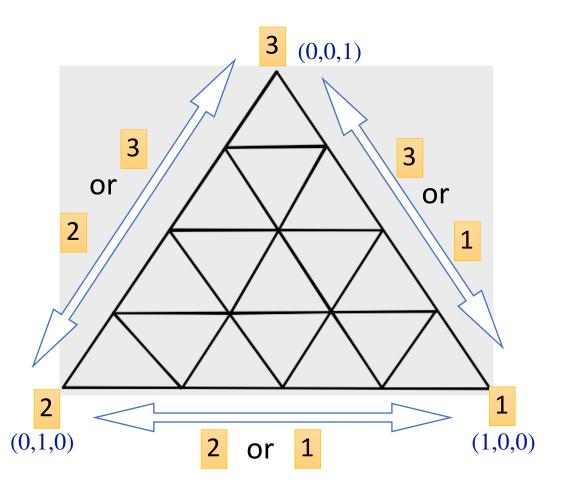
Sperner's lemma \implies

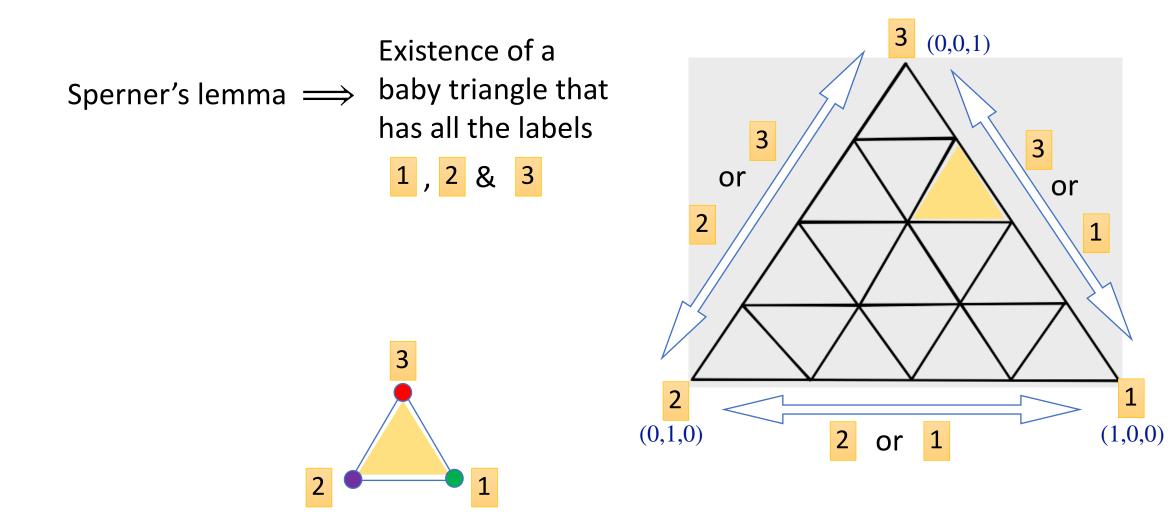


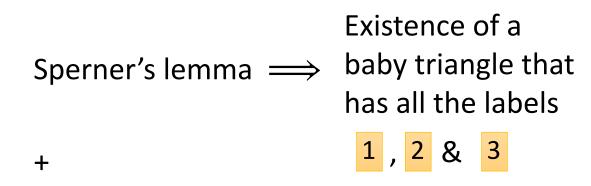
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Existence of a baby triangle that has all the labels

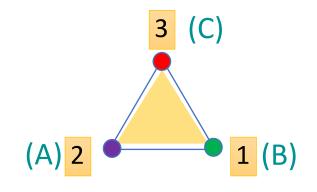
1,2&3

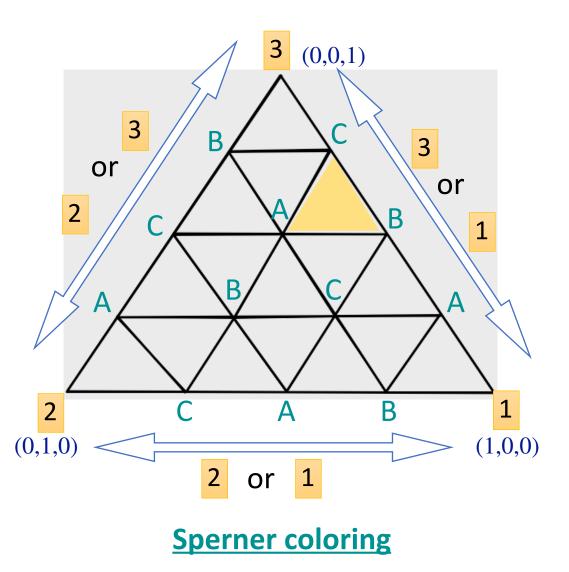


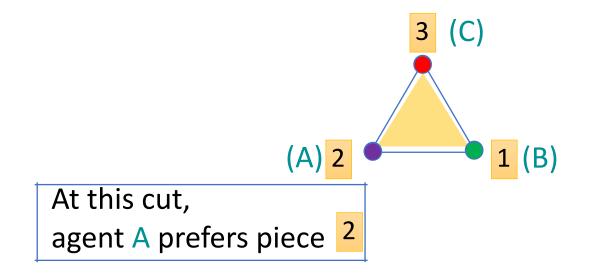


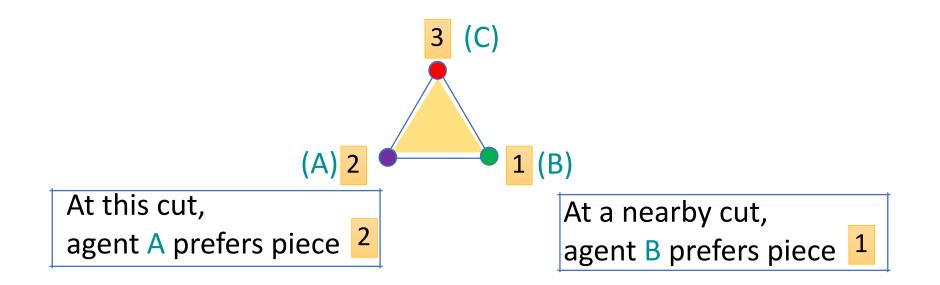


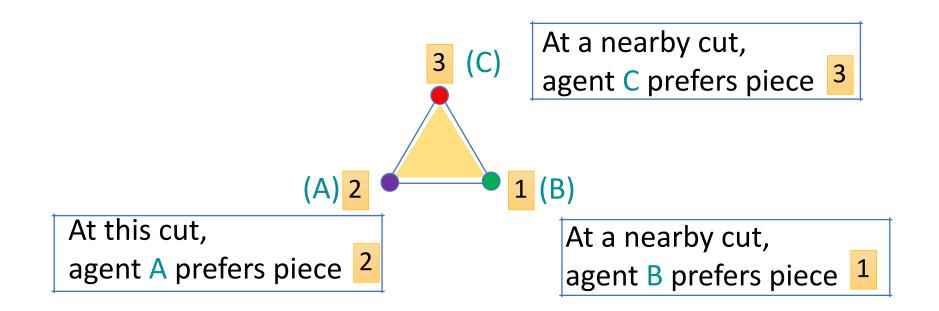
Ownership labeling \Longrightarrow

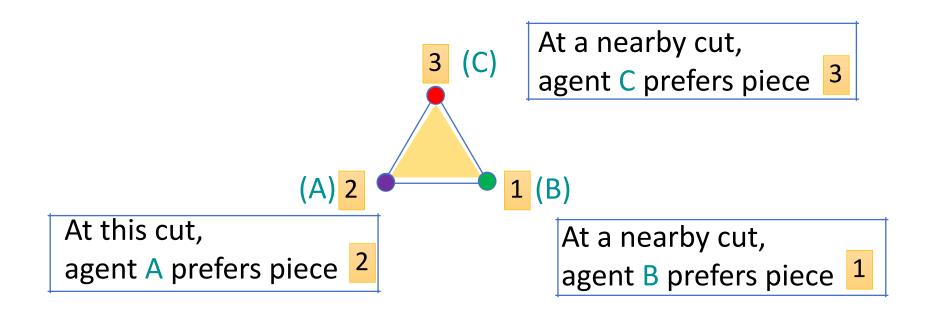




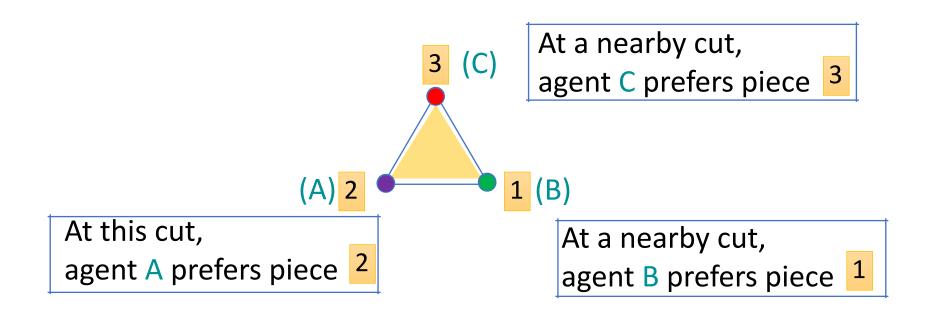






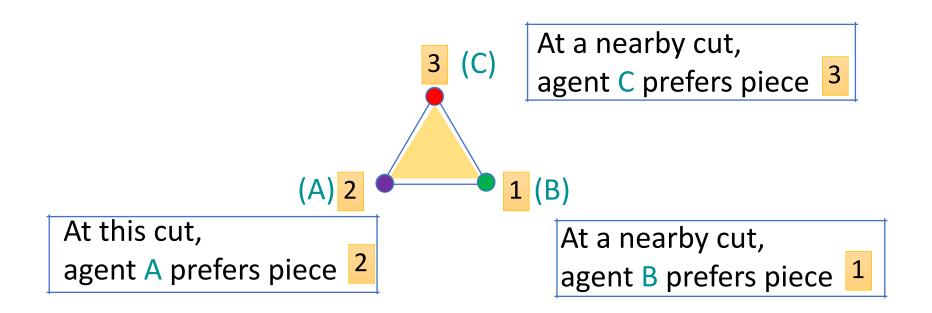


What we have is **not** a single cut (and hence not a single allocation), but three *nearby cuts*, where *envy-free-type* of thing is going on.



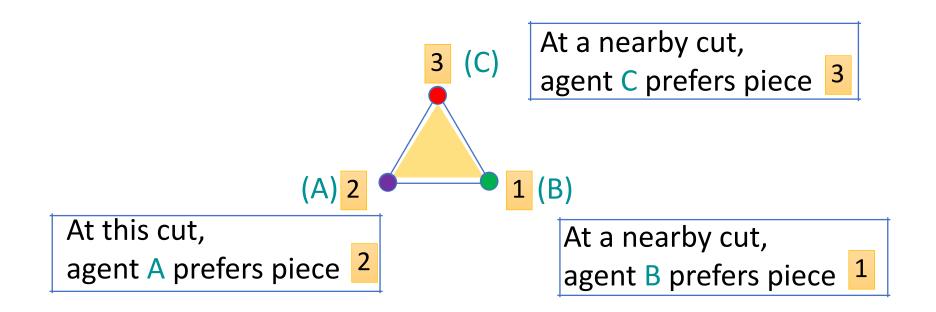
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A single cut where all three agents prefer different pieces \implies



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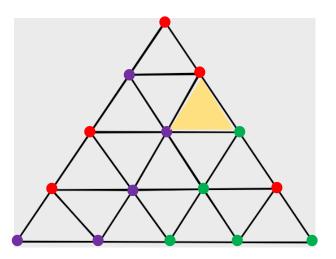
A single cut where all three agents prefer different pieces \implies EF cake division



Sperner's Lemma \implies A set of three '*nearby*' cuts where different agents prefer different pieces

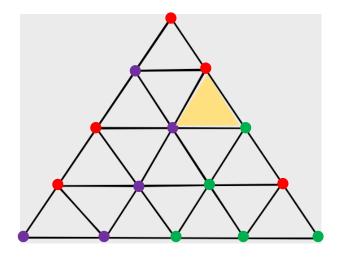
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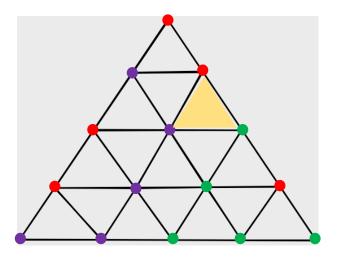
Imagine making this triangulation finer and finer



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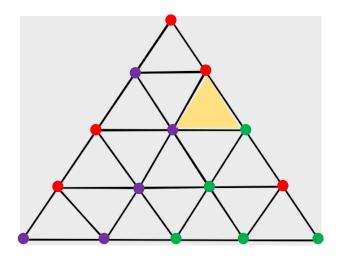
- we will have increasingly 'nearby' cuts
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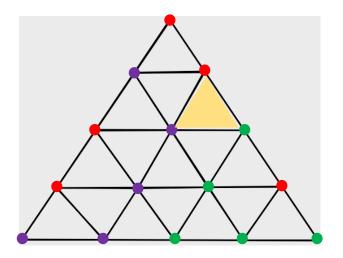


We can do something more: use convergence properties

Sperner's Lemma \implies A set of three '*nearby*' cuts where different agents prefer different pieces

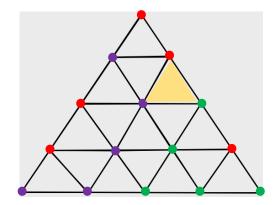
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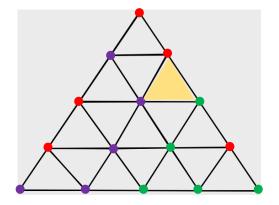


Valuations are (*topologically*) *closed* \implies the *limiting cut* has to be envy-free

Sperner's Lemma \implies A set of three '*nearby*' cuts where different agents prefer different pieces

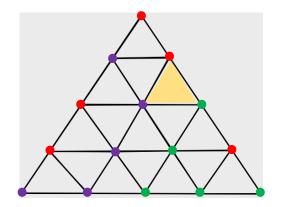


Third Assumption: valuations are closed



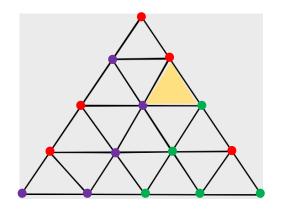
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Denote a cut $X = (x_1, x_2, x_3)$. Consider a sequence of cuts $X^{(1)}, X^{(2)}, X^{(3)}, ...$



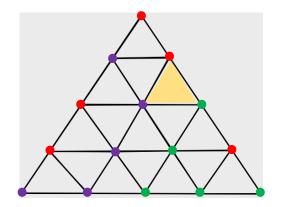
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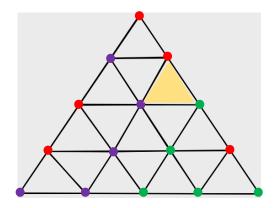
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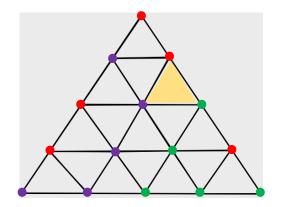
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(Using Bolzano-Weistrass convergence theorem)



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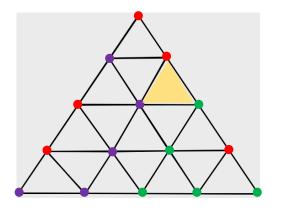
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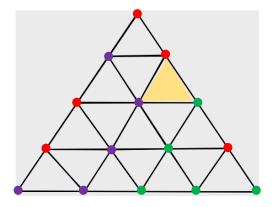
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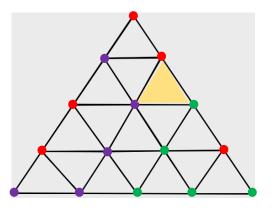
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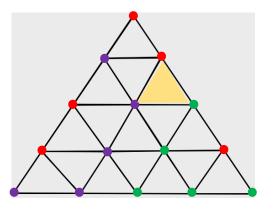
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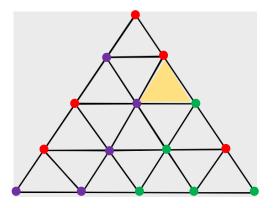


Third Assumption: valuations are closed **Closed under limit!**

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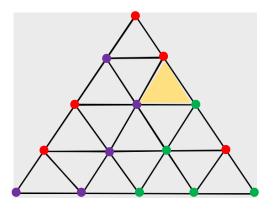


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Idea: There is a limiting cut where we can turn approximate EF into exact EF

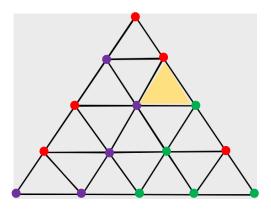


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We can take increasingly finer triangulations. They will all converge to a single cut-point, and at that cut, all three agents will prefer different pieces



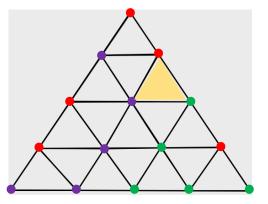
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Exact envy-free connected division



Sperner's Lemma

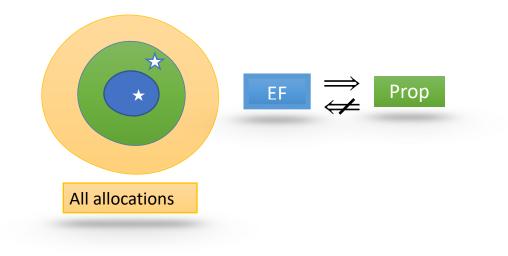
Convergence-based existential proof of envy-free cake divisions with connected pieces

Sperner's Lemma

Convergence-based existential proof of envy-free cake divisions with connected pieces (and hence, it does not lead to an efficient algorithm)

Sperner's Lemma

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Stromquist [1980], Su [1999]

Envy-free cake division with **connected** pieces exist for any number of agents

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Stromquist, J. of Combinatorics 2008

No finite-query protocols exists for connected EF cake division

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[ABKR] WINE'19

(Fair and Efficient Cake Division with Connected Pieces)(28 May)

An efficient algorithm: 1/2-EF + 1/3-NSW allocation for *connected* EF cake division

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[ABK**R**] *EC'20*

(Fair Cake Division under Monotone Likelihood Ratios)(25 June)

Efficient algorithms for connected EF cake division for a broad class of instances

Query Complexity of Envy-freeness

