## 

# Topics in Computational Social Choice Theory 

Lecture 4: Existence of envy-free cake divisions

Nidhi Rathi

## Last Lecture: Introduction to Cake Cutting

- The resource: Cake [0,1] (heterogeneous and divisible)
- Set of agents: $\{1,2, \ldots, n\}$
- Piece of a cake: finite union of subintervals of $[0,1]$
- Valuation function $v_{i}$ : Agent $i$ values piece $X$ at $v_{i}(X) \geq 0$



## Fairness Notions

## Allocation:

A partition $A=\left(A_{1}, A_{2}, \ldots, A_{n}\right)$ of the cake $[0,1]$ where piece $A_{i}$ belongs to agent $i$


- Proportionality: for each agent $i \in[n]$, we have $v_{i}\left(A_{i}\right) \geq 1 / n$ [Steinhaus, 1948]
- Envy-freeness: for every pair $\boldsymbol{i}, \boldsymbol{j} \in[n]$ of agents, we have $v_{i}\left(A_{i}\right) \geq v_{i}\left(A_{j}\right)$ [Foley 1967]


## Query Complexity of Proportionality

## Prop $\neq$ EF

for > two agents

Set of all Allocations


## Existence of Envy-free Cake Divisions

- Computing an envy-free cake division:
- Cut-and-choose: between two agents using 2 queries
- Selfridge-Conway: among three agents using 8 queries

What about $n \geq 4$ agents?

## Existence of Envy-free Cake Divisions



Stromquist [1980], Su [1999]
connected pieces
Envy-free cake division exist for any number of agents

Sperner's Lemma

## Sperner's Lemma

A beautiful lemma that, on the face of it, has nothing to do with cake division

## Sperner's Lemma



## Sperner's Lemma

## Ingredients:

1) A triangle that is subdivided into smaller triangles (Formal terms: simplex and its triangulation)


## Sperner's Lemma

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1) A triangle that is subdivided into baby triangles (Formal terms: simplex and its triangulation)


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1) A triangle that is subdivided into baby triangles (Formal terms: simplex and its triangulation)
2) Sperner coloring

- Main vertices have distinct colors



## Sperner's Lemma

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1) A triangle that is subdivided into baby triangles (Formal terms: simplex and its triangulation)
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- Main vertices have distinct colors
- Boundary vertices inherit colors of the adjacent main vertices



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(odd number)


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## Sperner's Lemma

- Entire main triangle: HOUSE
- Baby triangles: ROOMS
- $-\longrightarrow$ : DOOR



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## Observation 1:

Number of doors on the boundary is ODD


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Observation 2:
A room can have 0,1 , or 2 doors


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A room with 1 door is a fully colored baby triangle


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So, we enter the house through a door!


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Enter the house through a door.
The room we entered can have either 1 or 2 doors


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Keep walking!

- reach a fully colored baby triangle



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Keep walking!

- reach a fully colored baby triangle

- thrown out of the house


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Thrown out?


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- Cannot happen from (since no doors)



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- Entry and exit doors are paired up



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- There exists odd number of doors on the boundary. $\Longrightarrow$ we can enter again from another door!



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Think:
Why cannot such walks cycle back on themselves?


## Sperner's Lemma

- The number of rooms = finite $\Longrightarrow$ the walk terminates

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Why cannot such walks cycle back on themselves?


## Sperner's Lemma

- The number of rooms = finite $\Longrightarrow$ the walk terminates
- $\exists$ at least one walk that will take us to a fully colored sperner solution

Think:
Why cannot such walks cycle back on themselves?


## Sperner's Lemma

(odd number)

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(odd number)

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(odd number)

## Sperner's Lemma



Holds true for any dimension

## Cake division using Sperner's Lemma

Forest Simmons, popularized by Francis Su [1999]

## Cake division using Sperner's Lemma

- The resource: cake $[0,1]$ and n agents
- An allocation $\left(X_{1}, \ldots, X_{n}\right)$ is envy-free if $v_{i}\left(X_{i}\right) \geq v_{i}\left(X_{j}\right)$ for all $i, j$


## Cake division using Sperner's Lemma

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$$
\left(x_{1}, x_{2}, x_{3}\right): \text { a cut }
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## Cake division using Sperner's Lemma

Assumptions on preferences/valuations


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Goal: to invoke Sperner's lemma somehow


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- Given any cut $\left(x_{1}, x_{2}, x_{3}\right)$, each agent can point to its favorite piece
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Goal: to invoke Sperner's lemma somehow

$$
\text { Set of agents: }\{A, B, C\}
$$



Cake division using Sperner's Lemma

Cake division using Sperner's Lemma

(0,1,0)
(1,0,0)

Ownership labeling

## Cake division using Sperner's Lemma



Assign ownerships to each vertex such that each baby triangle consists of all three owners $\{A, B, C\}$.

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Assign ownerships to each vertex such that each baby triangle consists of all three owners $\{A, B, C\}$.
(There exists an efficient way to do this)

To generate a Sperner coloring, we go to a vertex, say some ( $x_{1}, x_{2}, x_{3}$ ), and ask its owner agent her most favorite piece in this cut

Ownership labeling

Cake division using Sperner's Lemma

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Sperner coloring

## Cake division using Sperner's Lemma

Sperner's lemma $\Longrightarrow$


## Cake division using Sperner's Lemma

## Existence of a

Sperner's lemma $\Longrightarrow$ baby triangle that has all the labels $1,2 \& 3$


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What we have is not a single cut (and hence not a single allocation), but three nearby cuts, where envy-free-type of thing is going on.

## Cake division using Sperner's Lemma



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A single cut where all three agents prefer different pieces $\Longrightarrow$

## Cake division using Sperner's Lemma



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A single cut where all three agents prefer different pieces $\Longrightarrow$ EF cake division

## Cake division using Sperner's Lemma



| Sperner's Lemma $\Longrightarrow$A set of three 'nearby' cuts where different agents <br> prefer different pieces |
| :---: |

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## Cake division using Sperner's Lemma

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| :---: |
|  |
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'Approximate' envy-free connected division

## Cake division using Sperner's Lemma



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## Cake division using Sperner's Lemma

Imagine making this triangulation finer and finer


> | Sperner's Lemma $\Longrightarrow$ |
| :---: |
|  |
|  |
| A set of three 'nearby' cuts where different agents |

'Approximate' envy-free connected division

## Cake division using Sperner's Lemma

Imagine making this triangulation finer and finer

- we will have increasingly 'nearby' cuts
- where we have envy-free like things happening


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## 'Approximate' envy-free connected division

## Cake division using Sperner's Lemma

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We can do something more: use convergence properties

$$
\begin{array}{|c|}
\hline \text { Sperner's Lemma } \Longrightarrow
\end{array} \begin{gathered}
\text { A set of three 'nearby' cuts where different agents } \\
\\
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\end{gathered}
$$

'Approximate' envy-free connected division

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Valuations are (topologically) closed $\Longrightarrow$ the limiting cut has to be envy-free

> | Sperner's Lemma $\Longrightarrow$ |
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|  |
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'Approximate' envy-free connected division

Cake division using Sperner's Lemma


## Cake division using Sperner's Lemma

Third Assumption: valuations are closed


## Cake division using Sperner's Lemma

Third Assumption: valuations are closed
Denote a cut $X=\left(x_{1}, x_{2}, x_{3}\right)$. Consider a sequence of cuts $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$


## Cake division using Sperner's Lemma

Third Assumption: valuations are closed
Denote a cut $X=\left(x_{1}, x_{2}, x_{3}\right)$. Consider a sequence of cuts $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$ Triangle is bounded


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Denote a cut $X=\left(x_{1}, x_{2}, x_{3}\right)$. Consider a sequence of cuts $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$ Triangle is bounded $\Longrightarrow$ the above sequence has a convergent subsequence


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(Using Bolzano-Weistrass convergence theorem)


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For any sequence of (converging) cuts $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$


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For any sequence of (converging) cuts $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$ if an agent $i$ prefers piece $k$ at each of $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$


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Triangle is bounded $\Longrightarrow$ the above sequence has a convergent subsequence.

For any sequence of (converging) cuts $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$ if an agent $i$ prefers piece $\boldsymbol{k}$ at each of $X^{(1)}, X^{(2)}, X^{(3)}, \ldots$ then she prefers the piece $\boldsymbol{k}$ in the limit as well!


## Cake division using Sperner's Lemma

## Third Assumption: valuations are closed Closed under limit!

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## Third Assumption: valuations are closed Closed under limit!

## Idea: There is a limiting cut where we can turn approximate EF into exact EF

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We can take increasingly finer triangulations. They will all converge to a single cut-point, and at that cut, all three agents will prefer different pieces


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Exact envy-free connected division


## Cake division using Sperner's Lemma

## Sperner's Lemma

Convergence-based existential proof of envy-free cake divisions with connected pieces

## Cake division using Sperner's Lemma

## Sperner's Lemma

Convergence-based existential proof of envy-free cake divisions with connected pieces (and hence, it does not lead to an efficient algorithm)

## Cake division using Sperner's Lemma

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Stromquist [1980], Su [1999]

## Envy-free cake divisions

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No finite-query protocols exists for connected EF cake division even for three agents!

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[ABKR] WINE'19
(Fair and Efficient Cake Division with Connected Pieces)(28 May)
An efficient algorithm: 1/2-EF +1/3-NSW allocation for connected EF cake division

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(Fair and Efficient Cake Division with Connected Pieces)(28 May)
An efficient algorithm: 1/2-EF +1/3-NSW allocation for connected EF cake division
[ABKR] EC'20
(Fair Cake Division under Monotone Likelihood Ratios)(25 June)
Efficient algorithms for connected EF cake division for a broad class of instances

## Query Complexity of Envy-freeness



