

Gibbard-Satterthwaite Theorem (and a bit about Revelation Principle)

Mechanism Design Without Money

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April 29, 2025





• voters $V = \{1, ..., n\}$, alternatives A with |A| = m

Social Choice Functions







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for every $a, b \in A$ with $a \neq b$

• a > c whenever a > b and b > c







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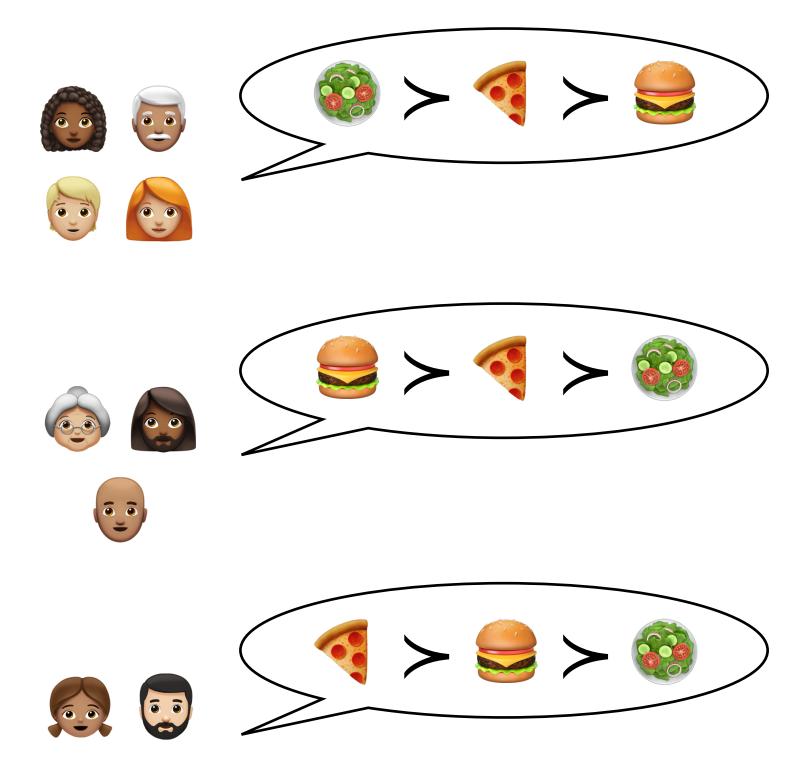






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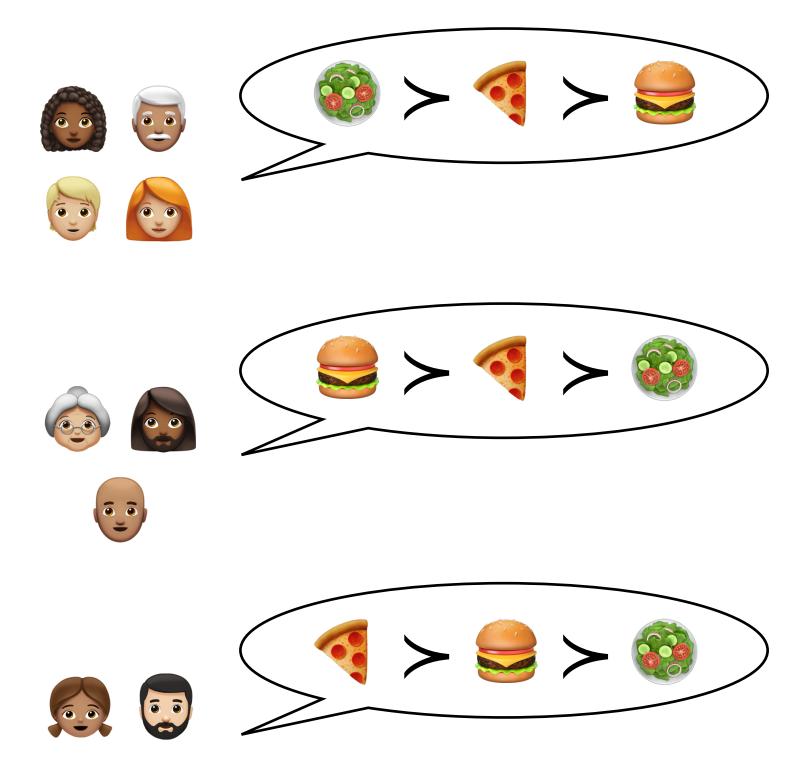






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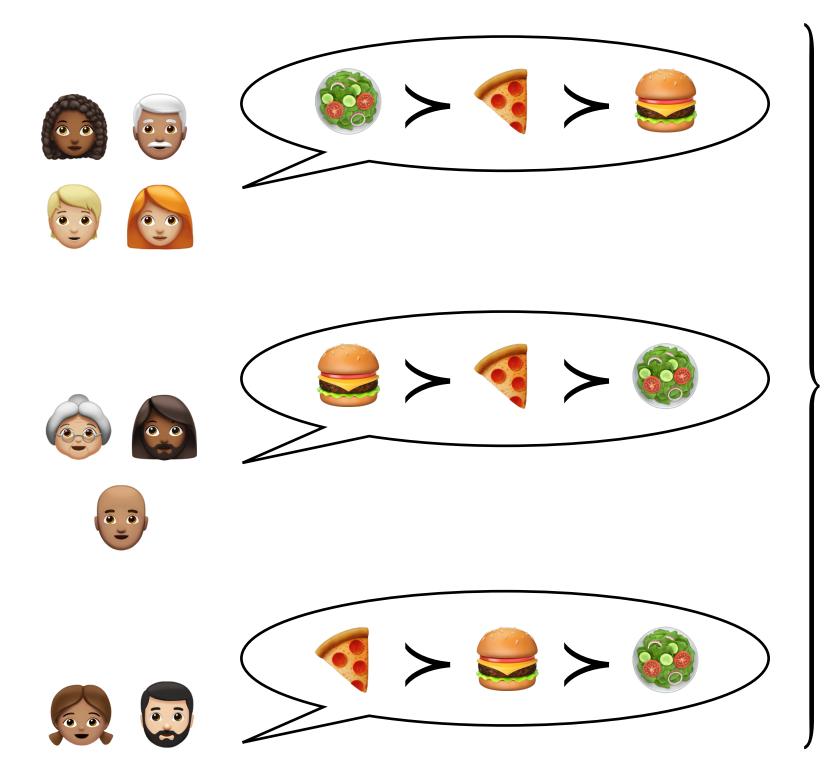






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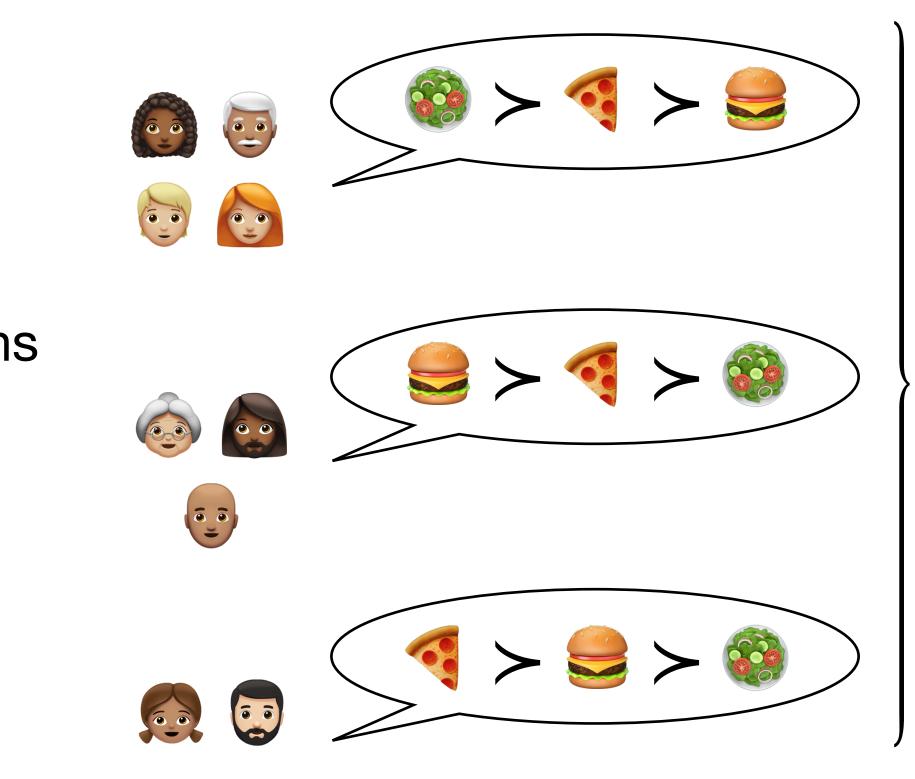






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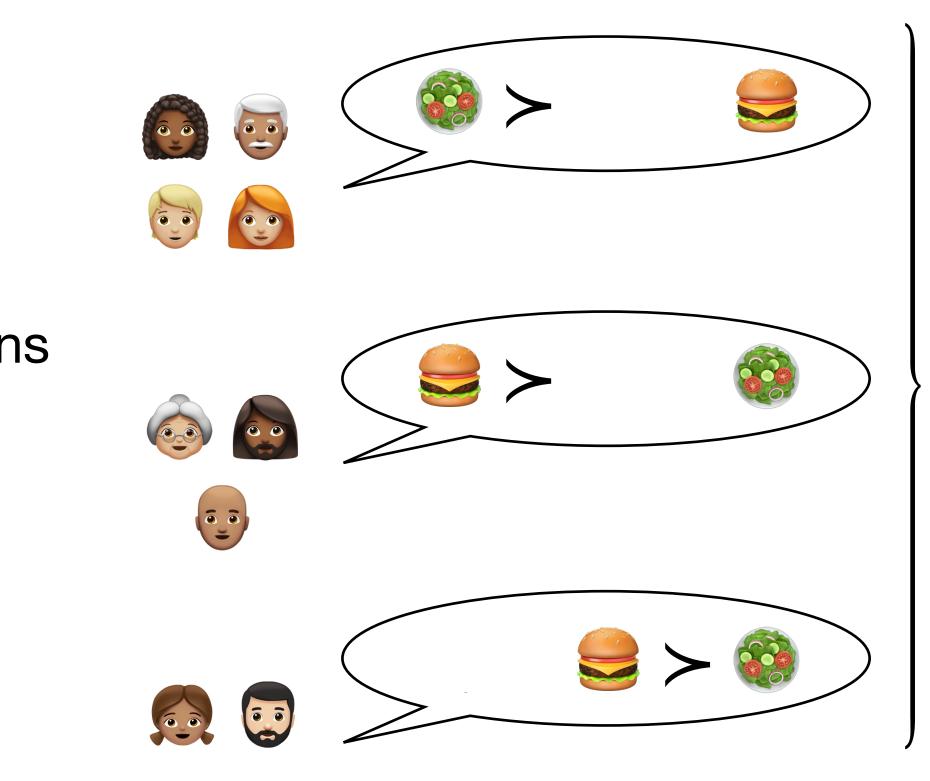






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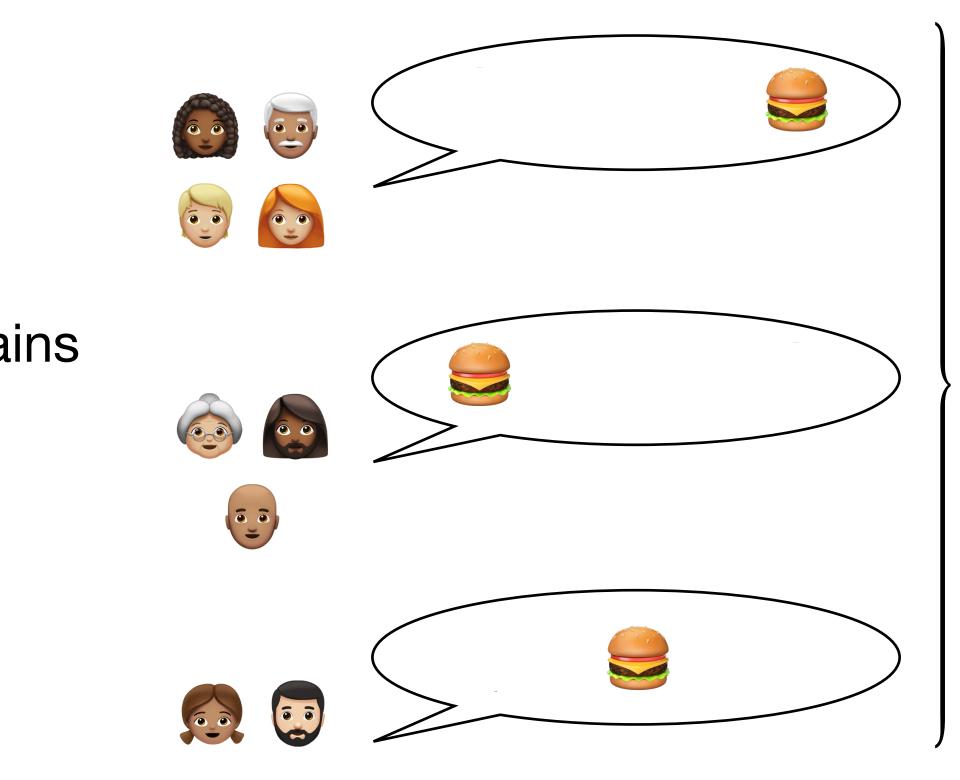






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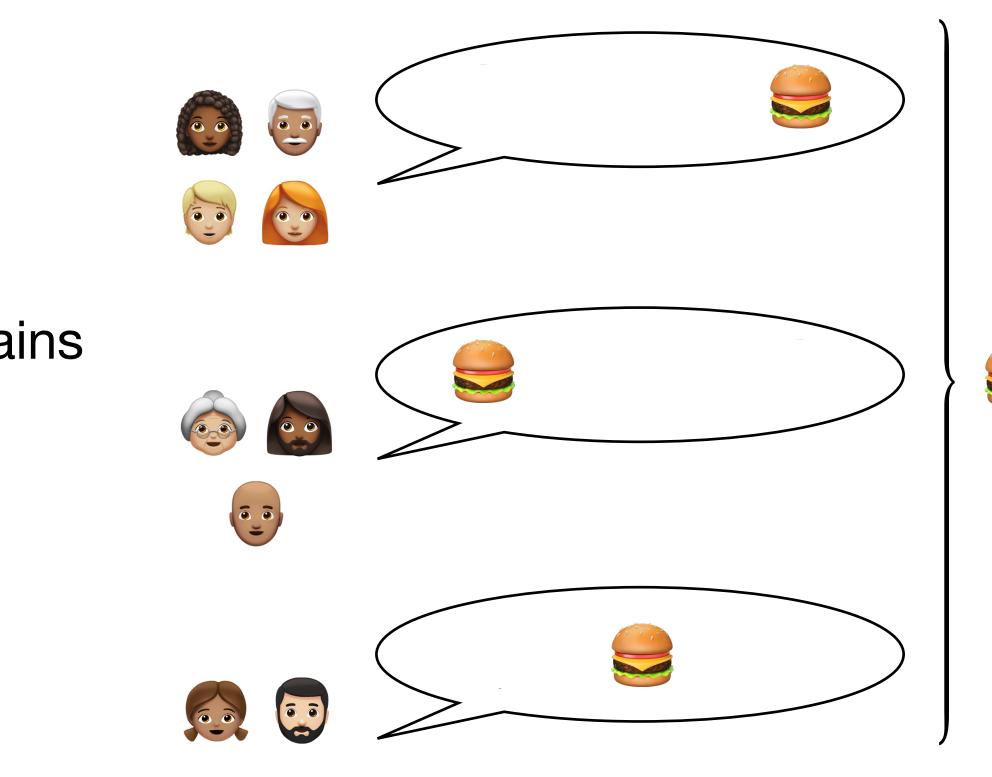






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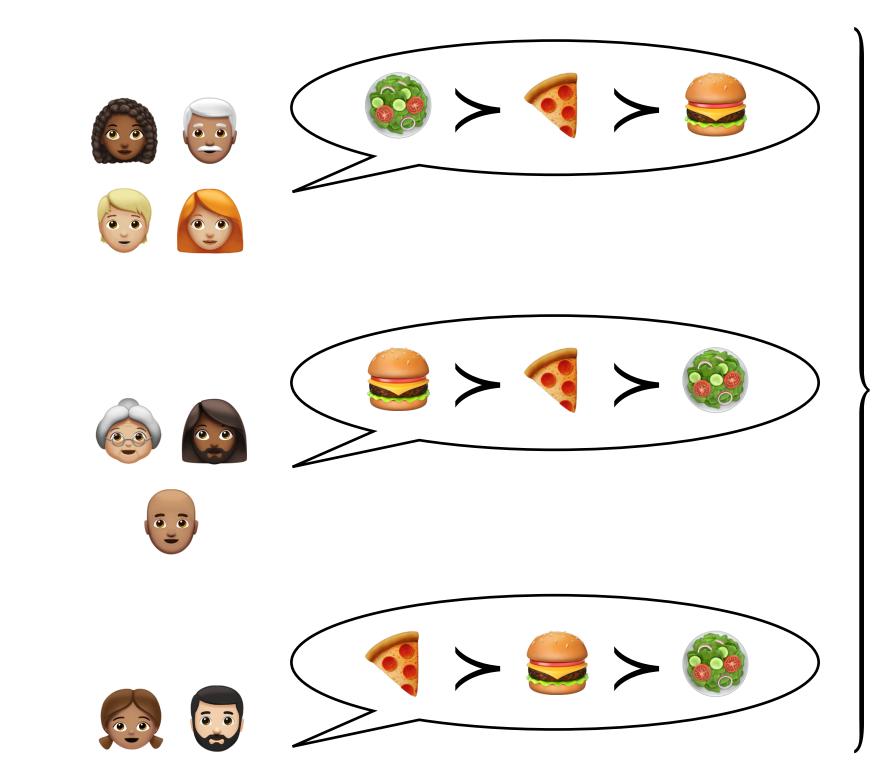






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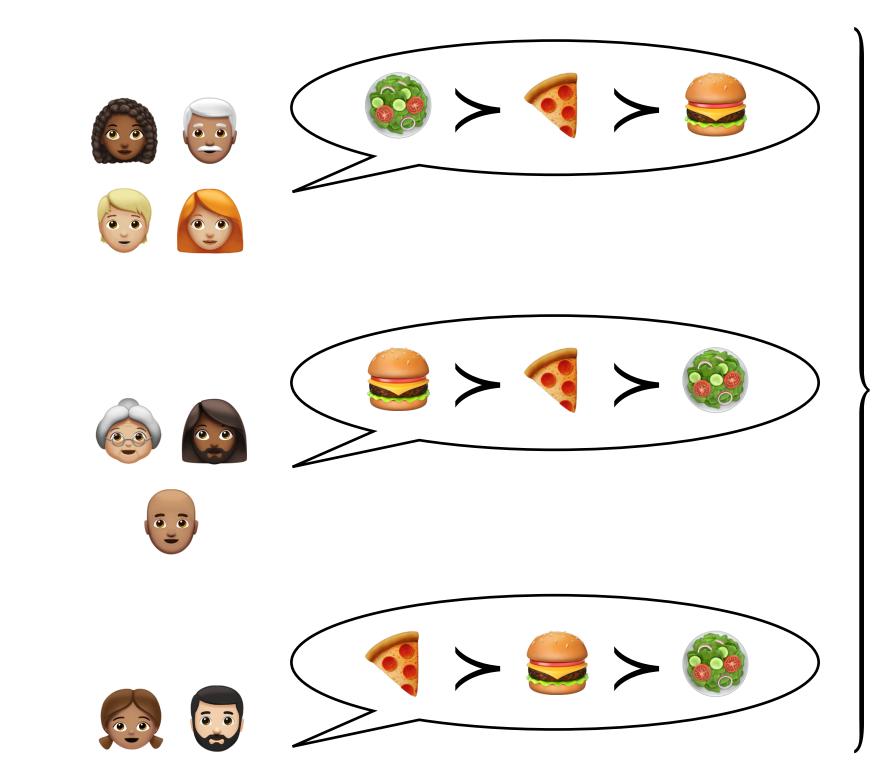
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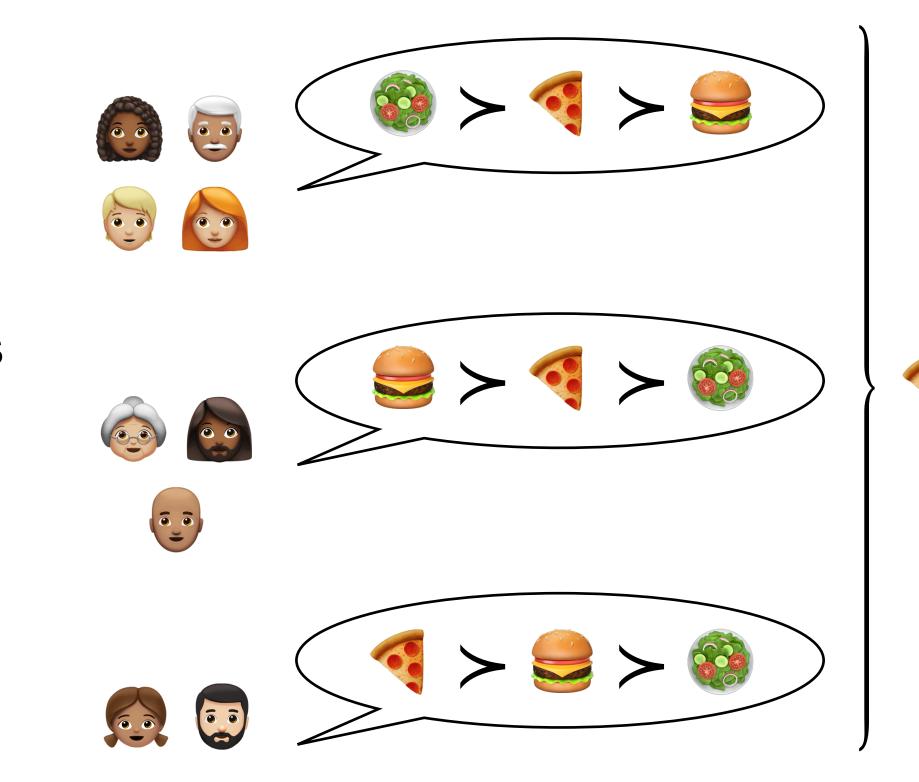
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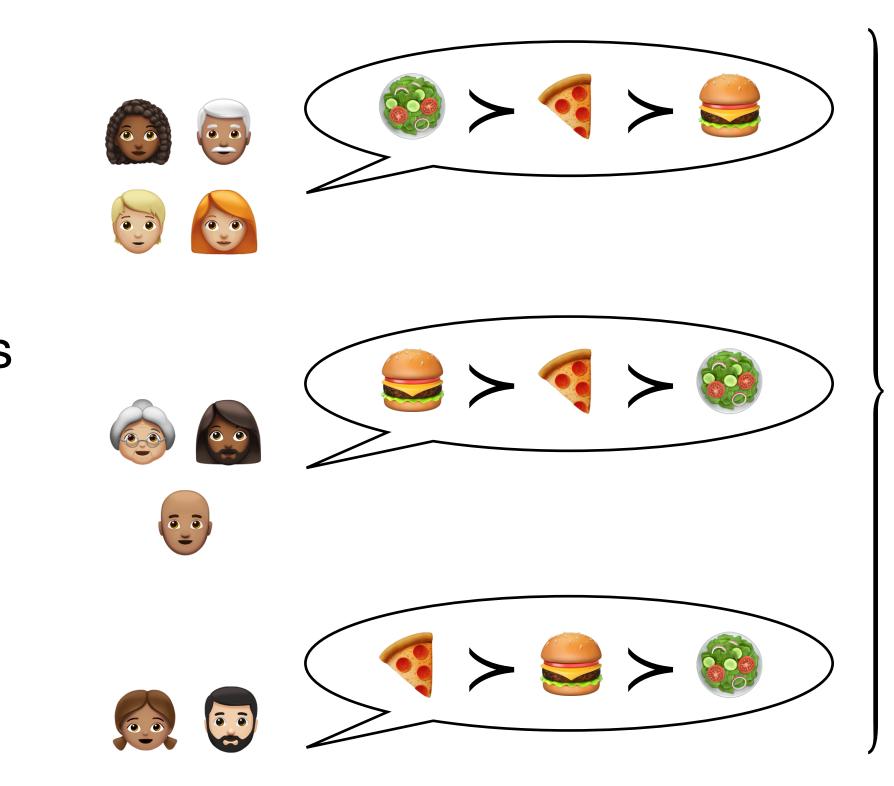






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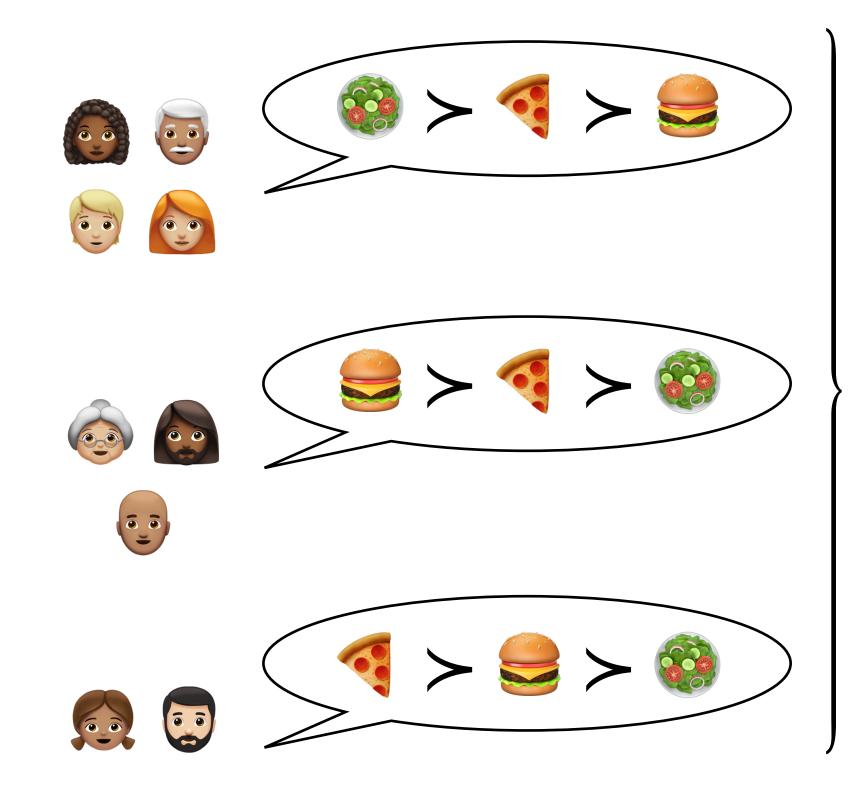


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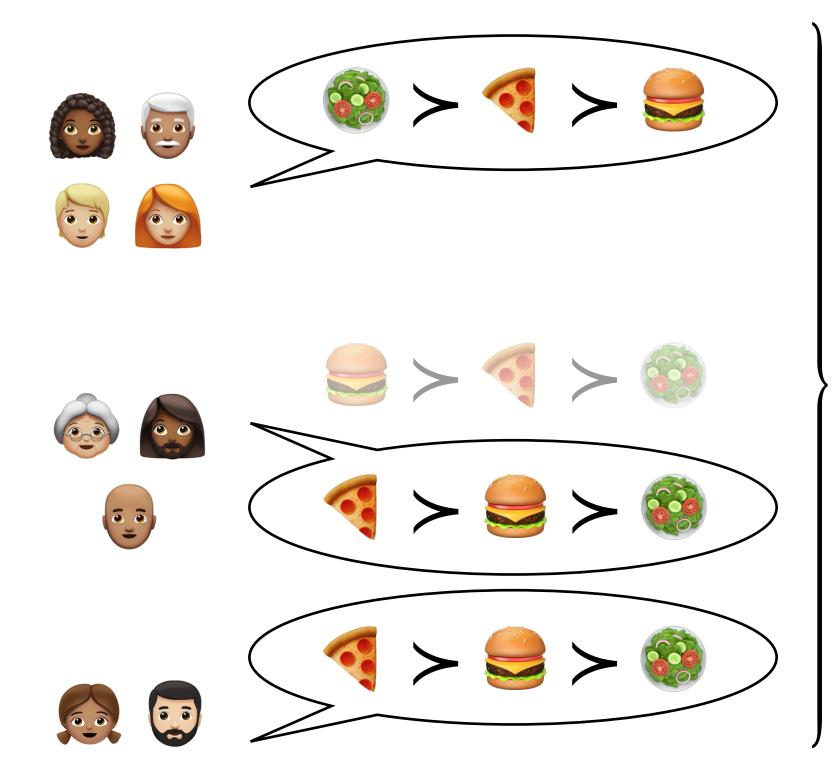


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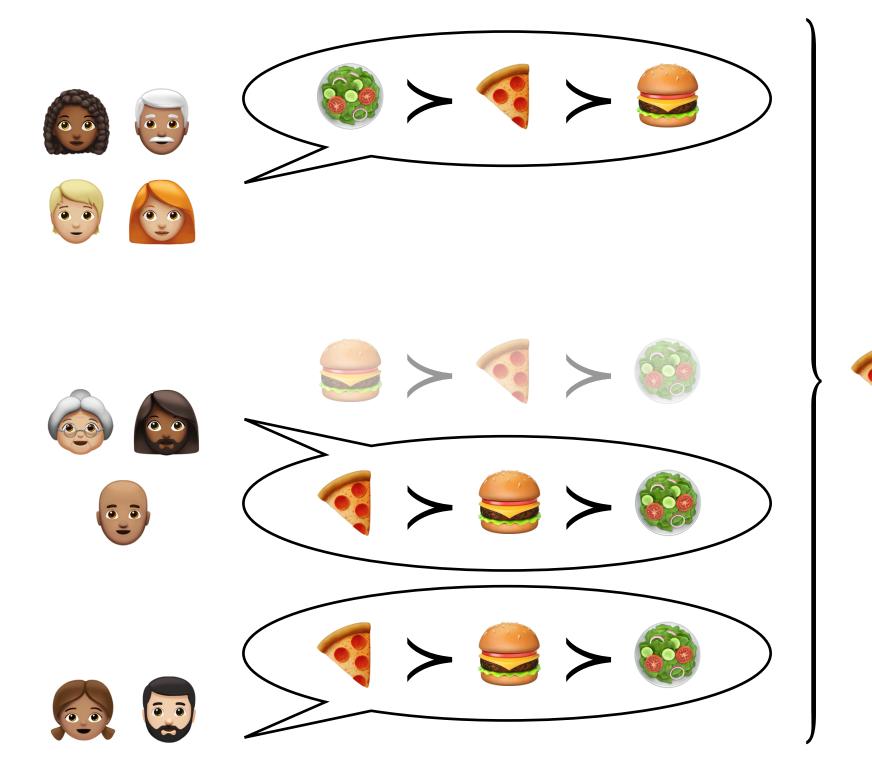


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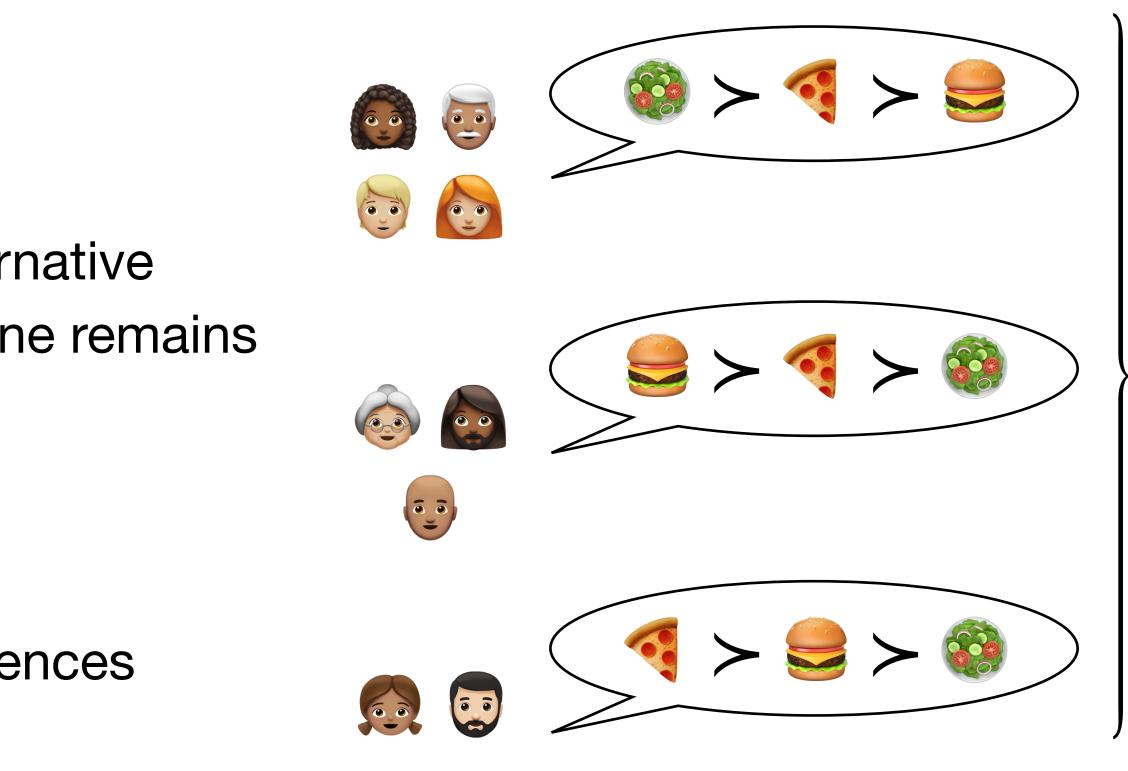


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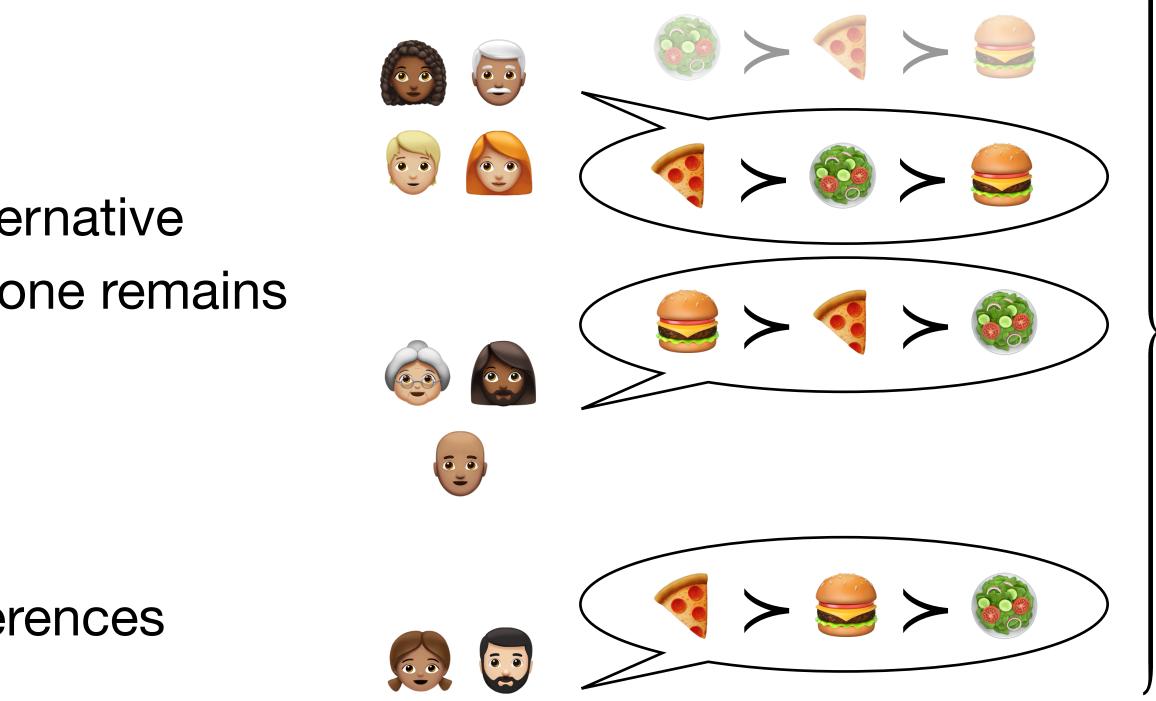


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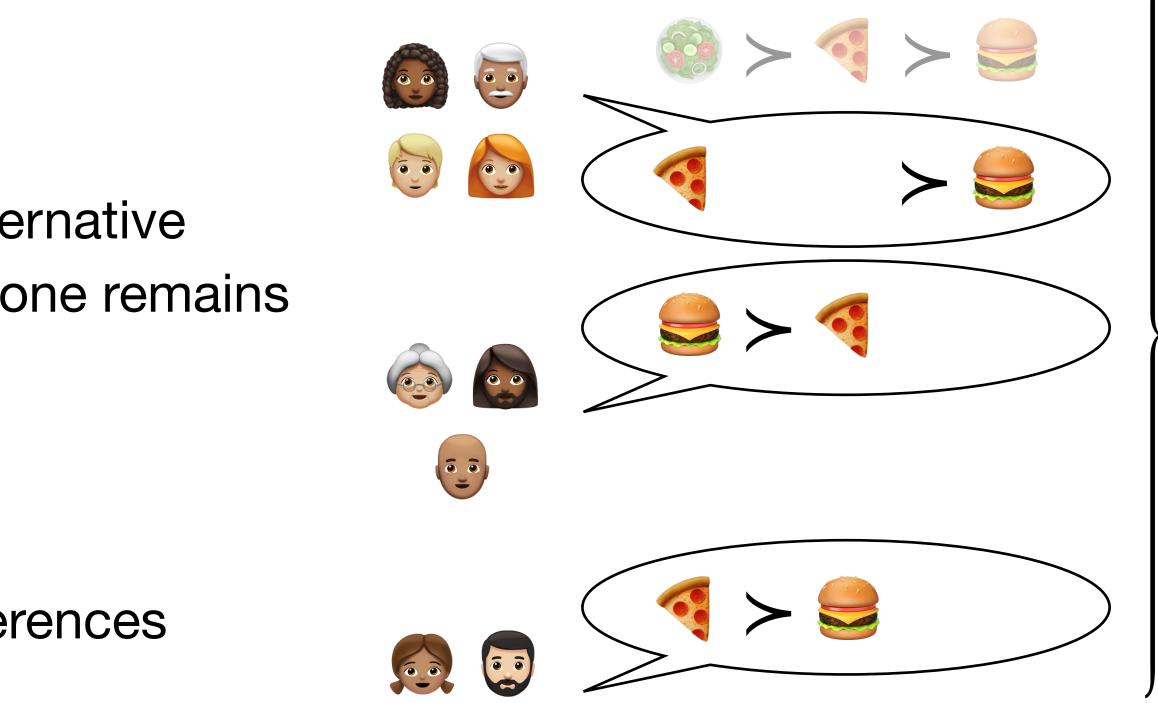


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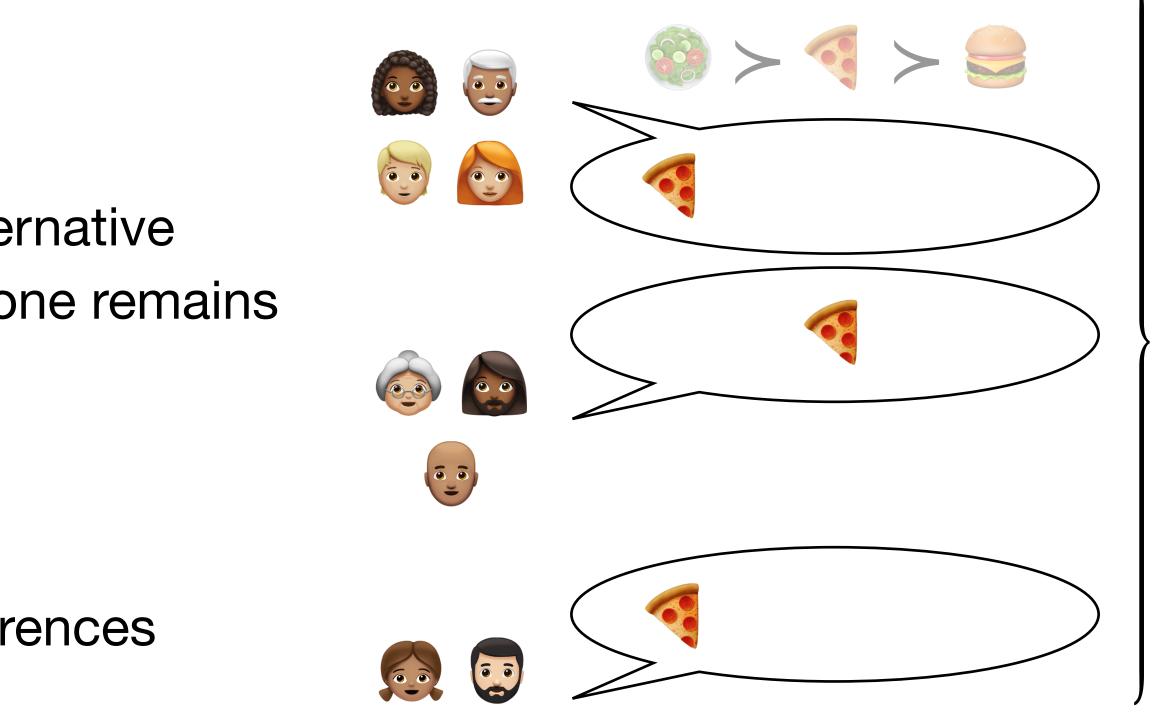


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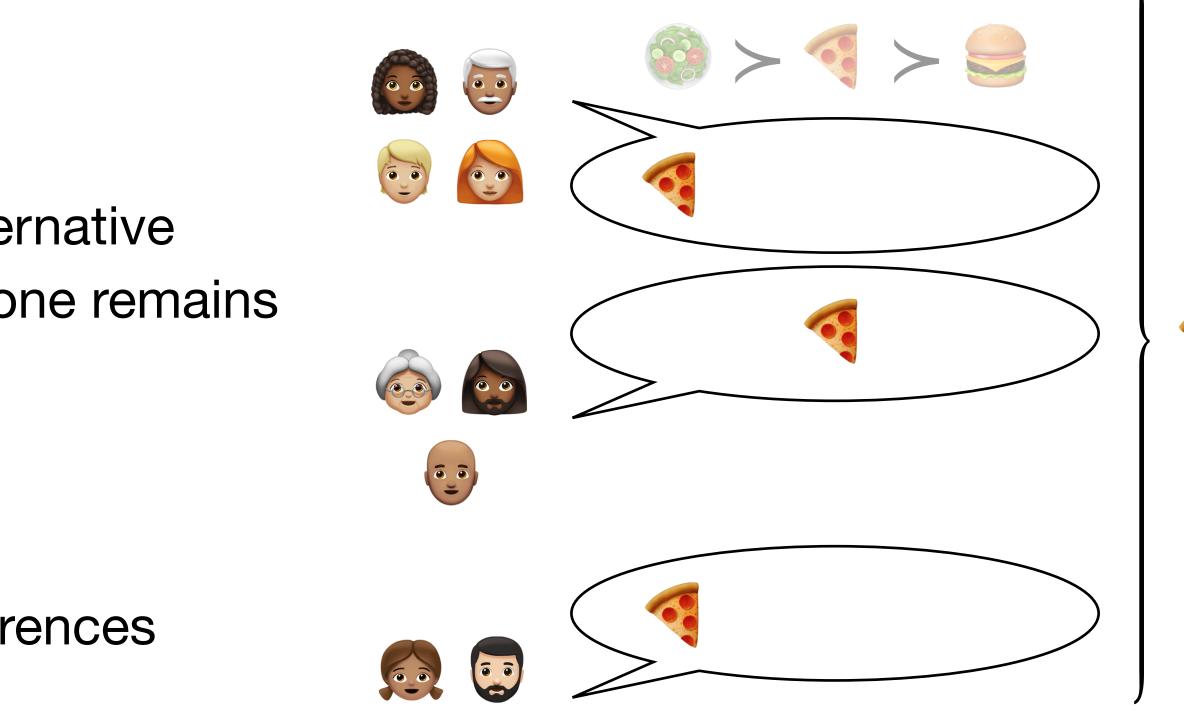


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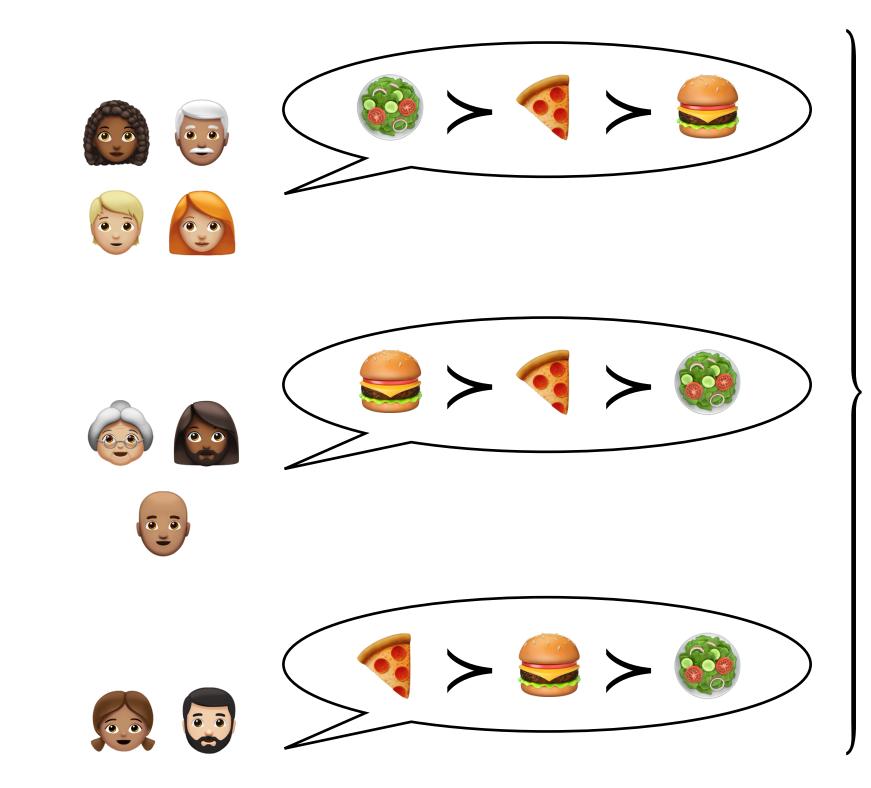




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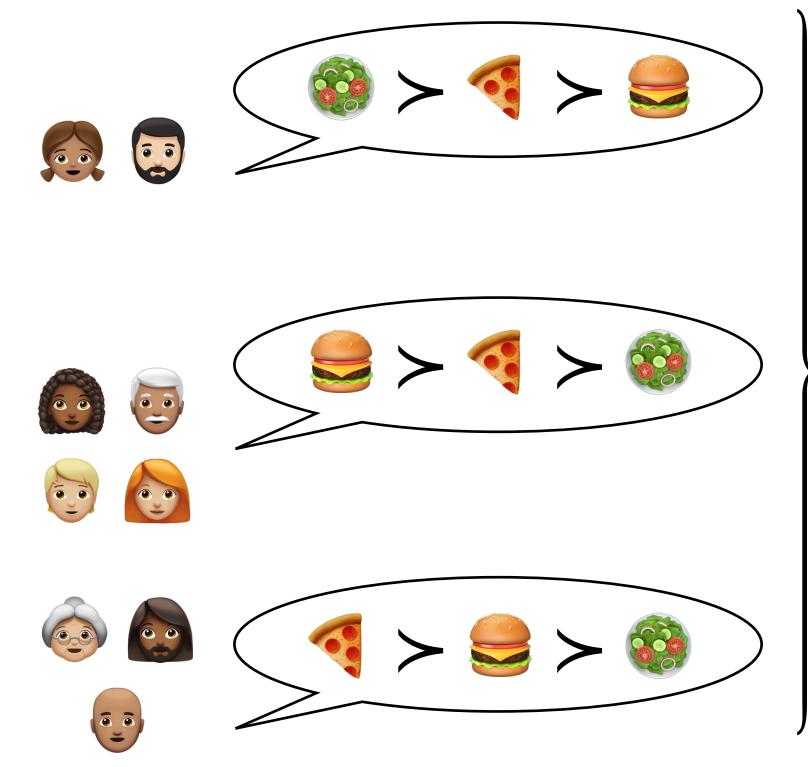




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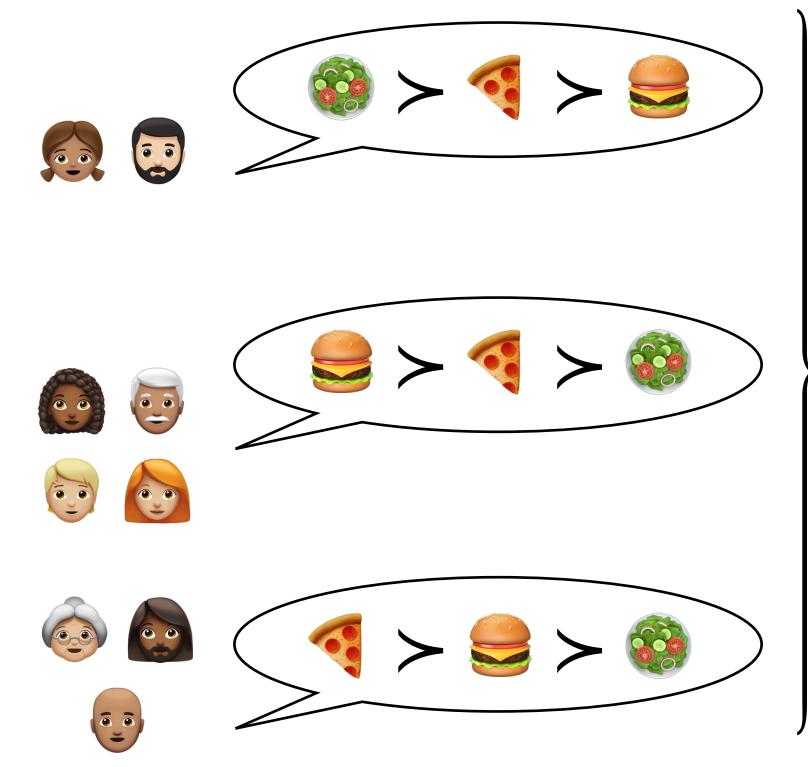
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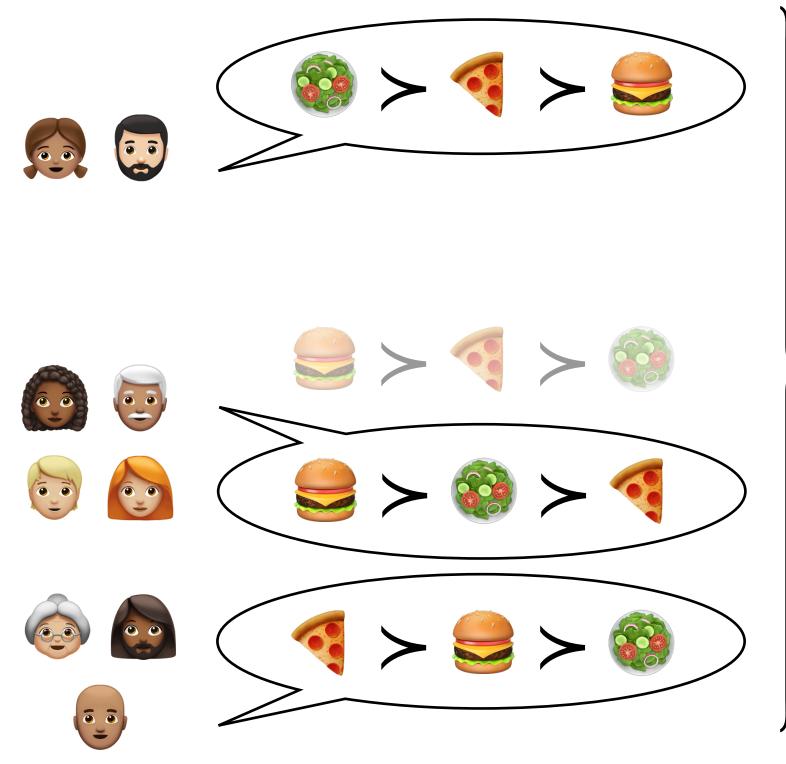
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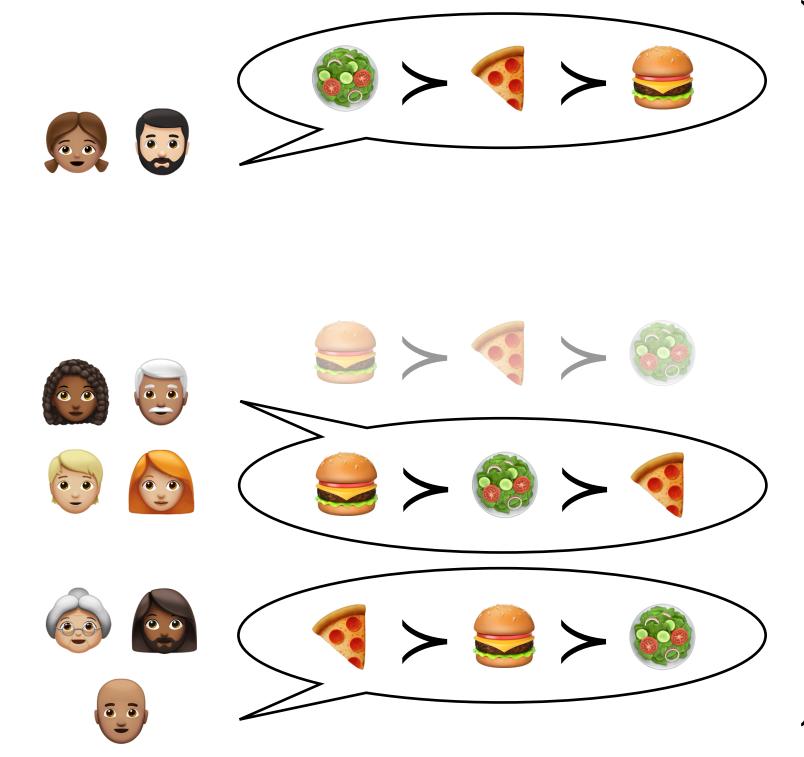


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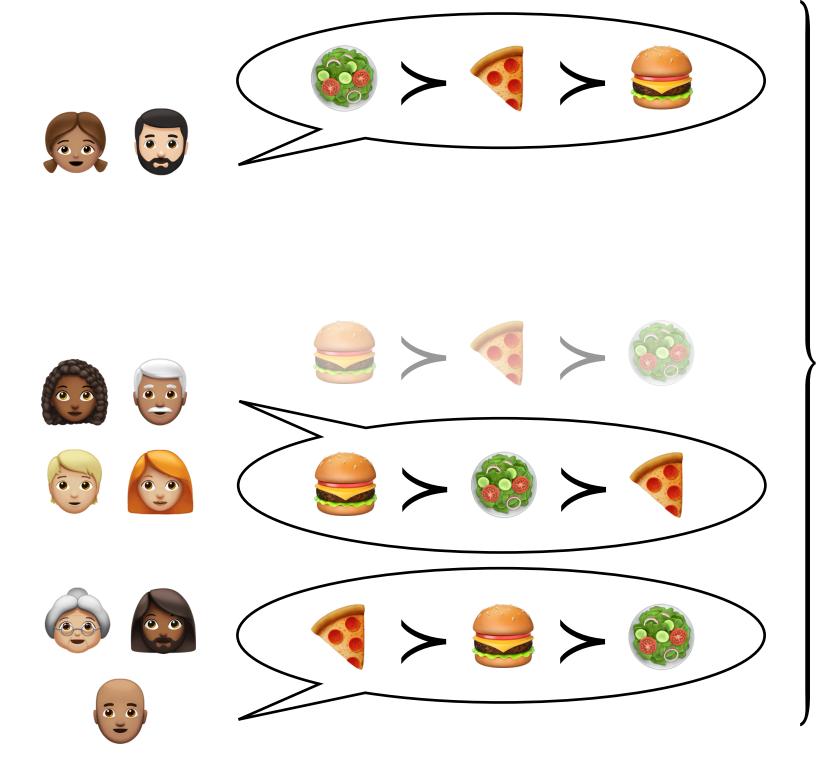


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\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_{n}
•	• • •	•	a	b	•••	•
	• • •	U	•	a	•••	b
•	• • •	•	•	•	•••	• •
a	• • •	•	b	•	• • •	•
b	• • •	a	•	•	•••	•







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\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_{κ}
•	• • •	•	a	b	•••	•
•	•••	b	•	a	•••	b
• • •	• • •	• •	• •	• •	• • •	• • •
a	• • •	•	b	•	•••	•
b	• • •	a	•	•	•••	•
			a			







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\succ_1	•••	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
b	• • •	b	a	b	• • •	b
•	• • •	•	b	•	• • •	•
•	• • •	•	• •	•	• • •	•
•	•••	•	•	•	• • •	•
a	•••	a	•	a	• • •	a
			a			



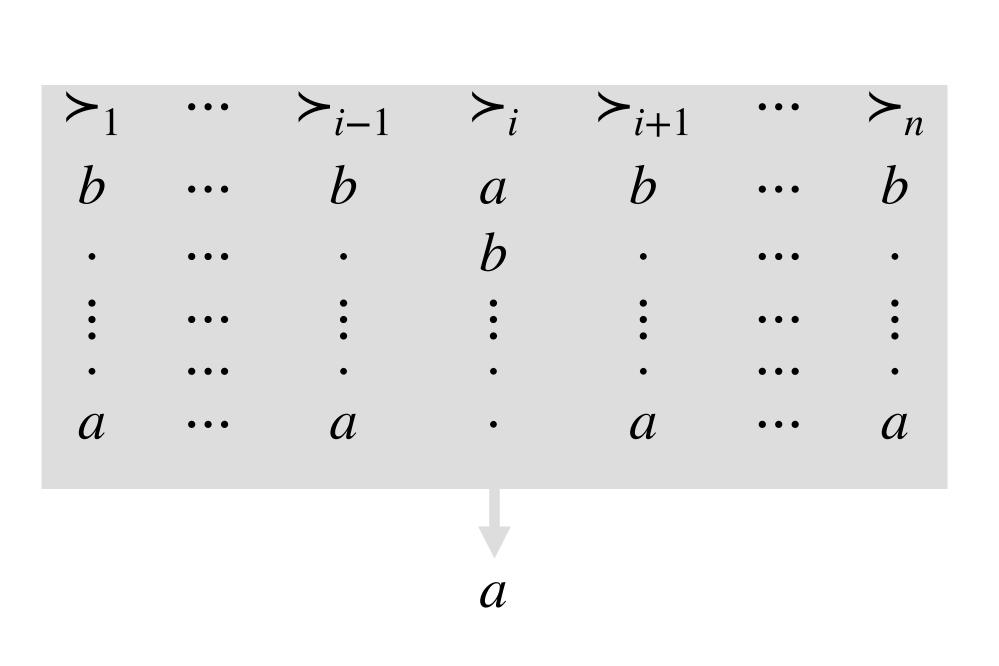








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- ▶ f is dictatorial if there is $i \in V$ such that $f(\succ) \succ_i a$ for every $\succ \in \mathscr{L}(A)^n$ and $a \in A \setminus f(\succ)$ top choice of *i* is always selected
- any dictatorship is surjective and strategyproof

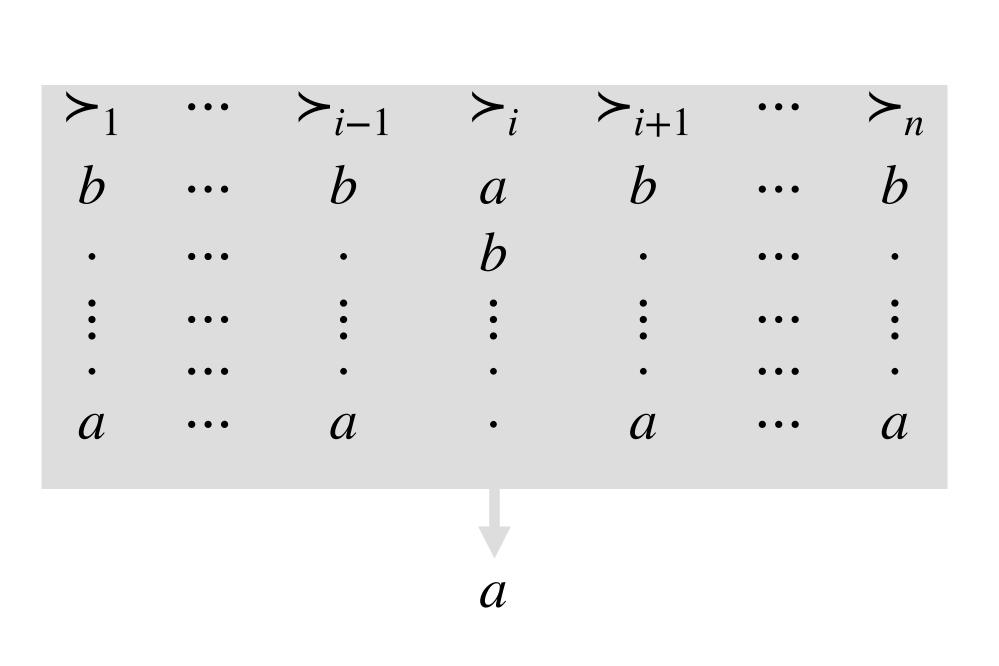








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- the converse is also true!





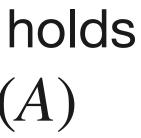


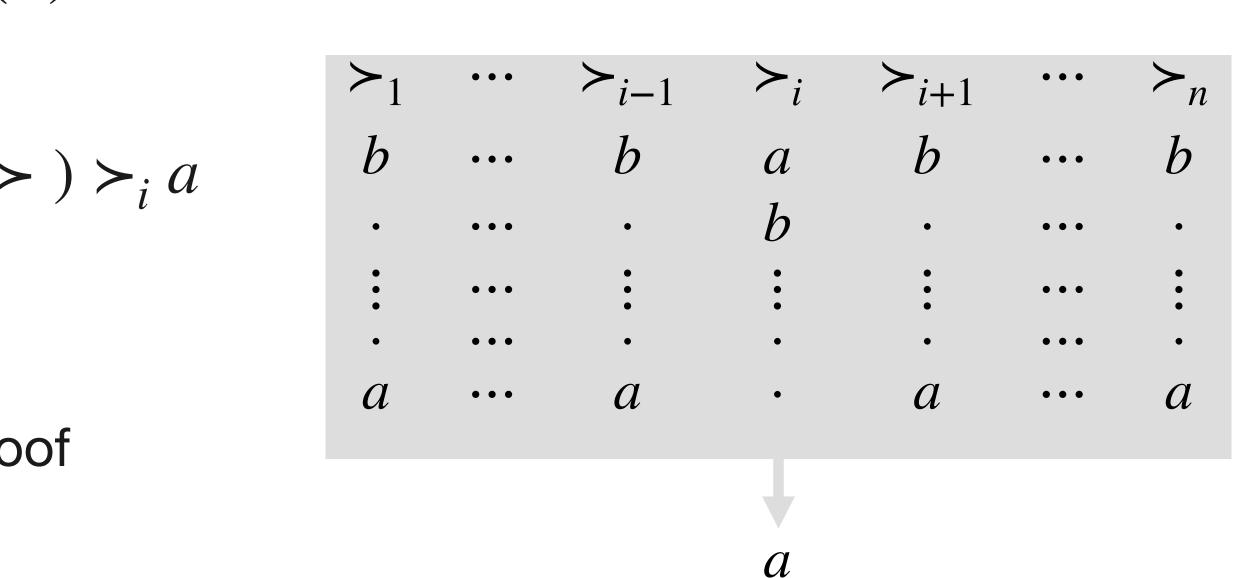


The Gibbard-Satterthwaite Theorem

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Theorem [Gibbard '73, Satterthwaite '75] Let $f: \mathscr{L}(A)^n \to A$ be a surjective and strategyproof social choice function, where $|A| \geq 3$. Then, f is dictatorial.











• f is unanimous if f(>) = a whenever $a \succ_i b$ for every $i \in V$ and $b \in A \setminus \{a\}$

when all voters have the same top choice, it is selected

The Muller-Satterthwaite Theorem





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The Muller-Satterthwaite Theorem

\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
a	• • •	a	a	a	• • •	a
•	• • •	•	•	•	•••	• •
•	• • •	•	•	•	• • •	•

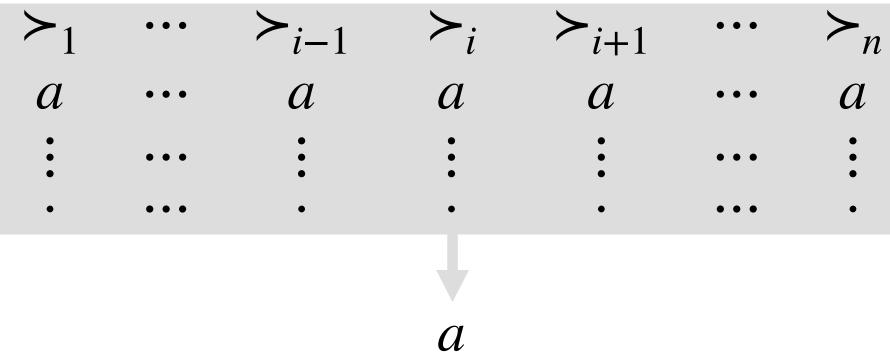




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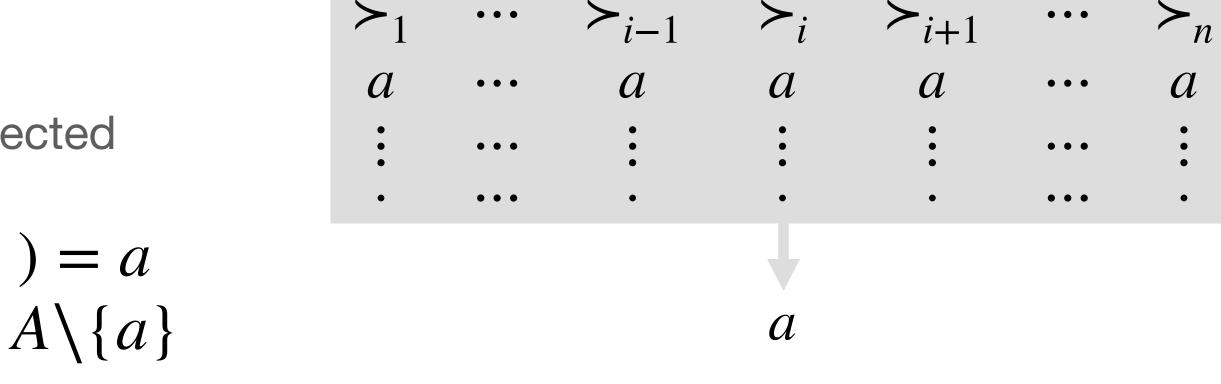
The Muller-Satterthwaite Theorem

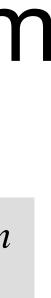






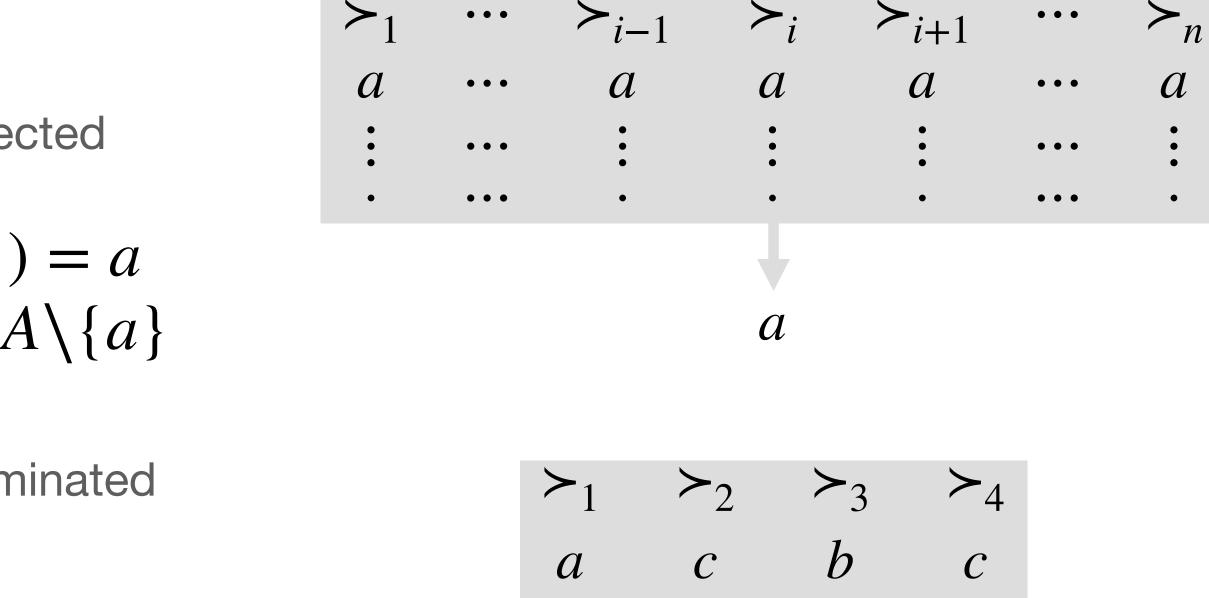
- f is unanimous if $f(\succ) = a$ whenever $a \succ_i b$ for every $i \in V$ and $b \in A \setminus \{a\}$ when all voters have the same top choice, it is selected
- f is monotone if $f(\succ') = a$ whenever $f(\succ) = a$ and $a \succ_i b \Rightarrow a \succ'_i b$ for all $i \in V$ and $b \in A \setminus \{a\}$ a selected alternative remains selected if dominated alternatives in all rankings remain dominated







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b

С

d

d

a

b

 \mathcal{A}

d

С

 \mathcal{A}

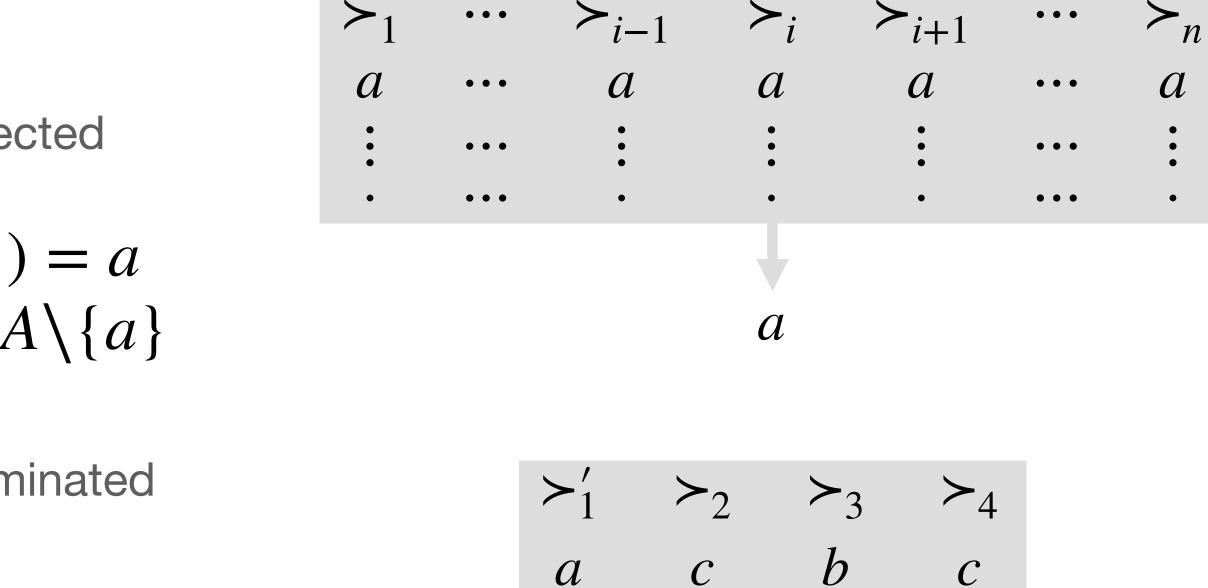
d

b





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a

d

b

C

d

b

a

C

 \mathcal{A}

d

b

a

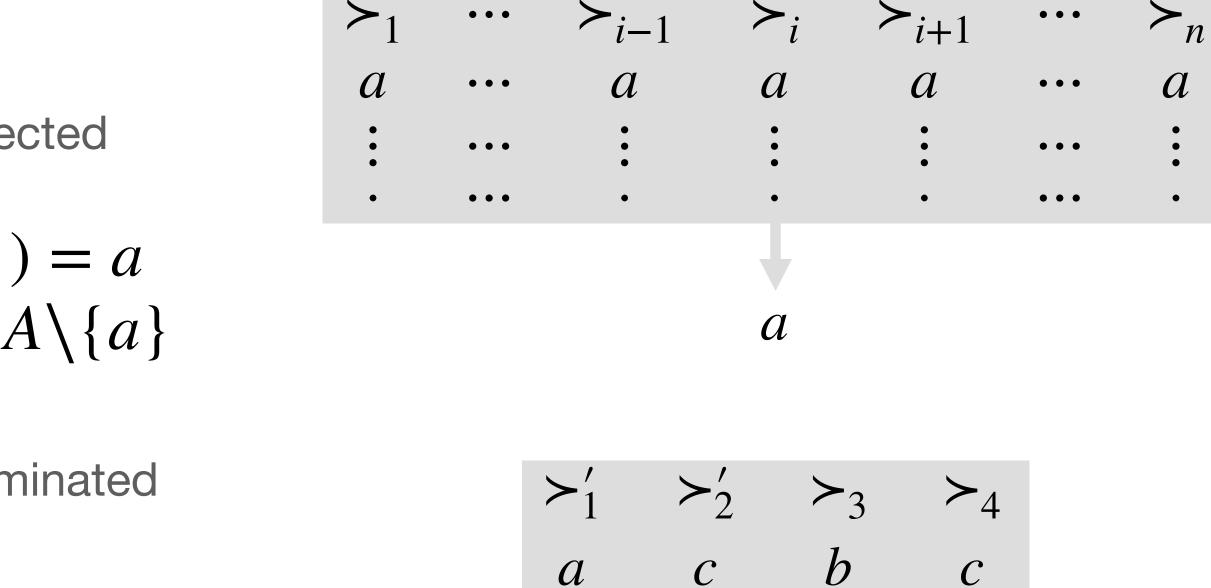
d

С





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d

b

C

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a

d

b

 \boldsymbol{a}

d

С

 \mathcal{A}

d

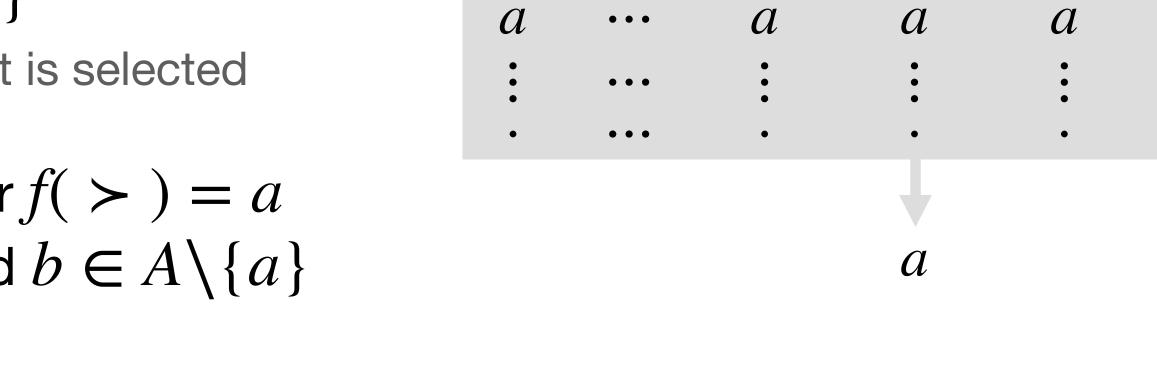
b





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 $>_{i-1}$



 \succ_1

\succ'_1	\succ_2'	≻'3	\succ_4
a	С	b	С
d	a	a	a
b	b	С	d
С	d	d	b
	C	l	

 \succ_i

 \succ_{i+1}

a

• • •

• • •

• • •

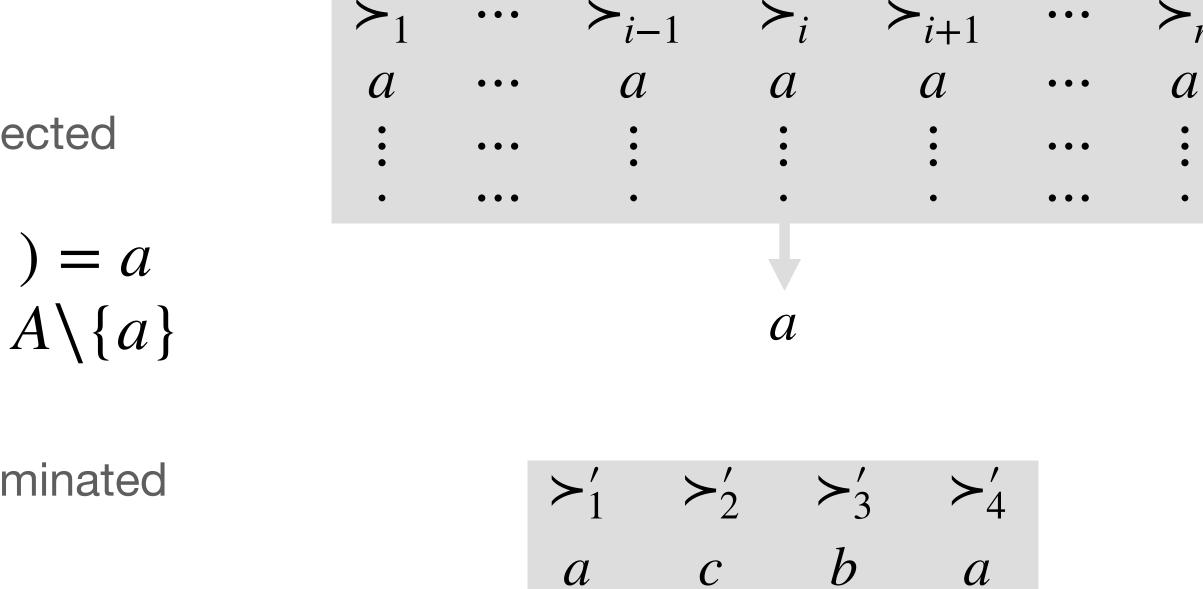
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 \mathcal{A}



- f is unanimous if $f(\succ) = a$ whenever $a \succ_i b$ for every $i \in V$ and $b \in A \setminus \{a\}$ when all voters have the same top choice, it is selected
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d

b

С

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a

b

d

a

С

d

b

d

С

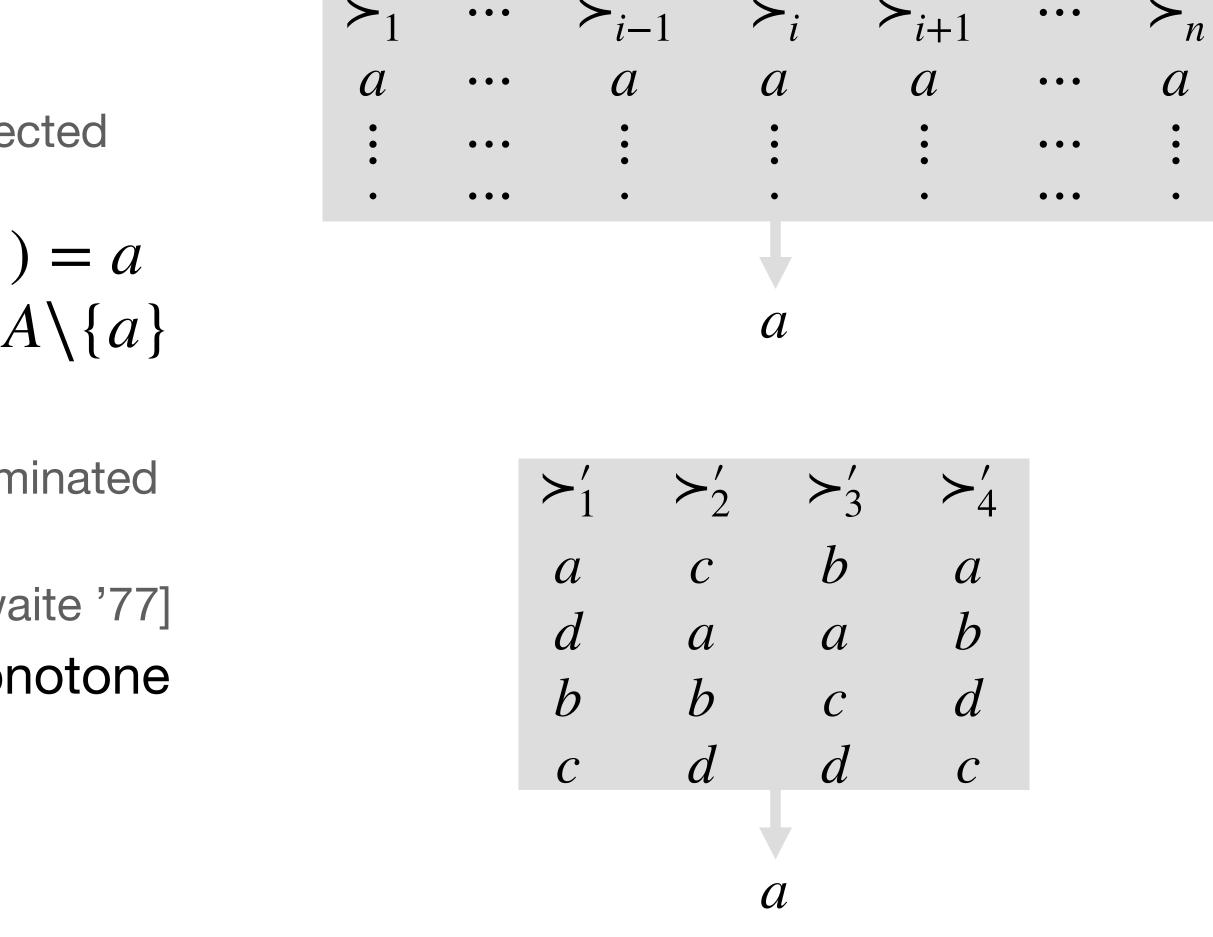




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Theorem [Muller, Satterthwaite '77] Let $f: \mathscr{L}(A)^n \to A$ be a unanimous and monotone social choice function, where $|A| \geq 3$. Then, f is dictatorial.

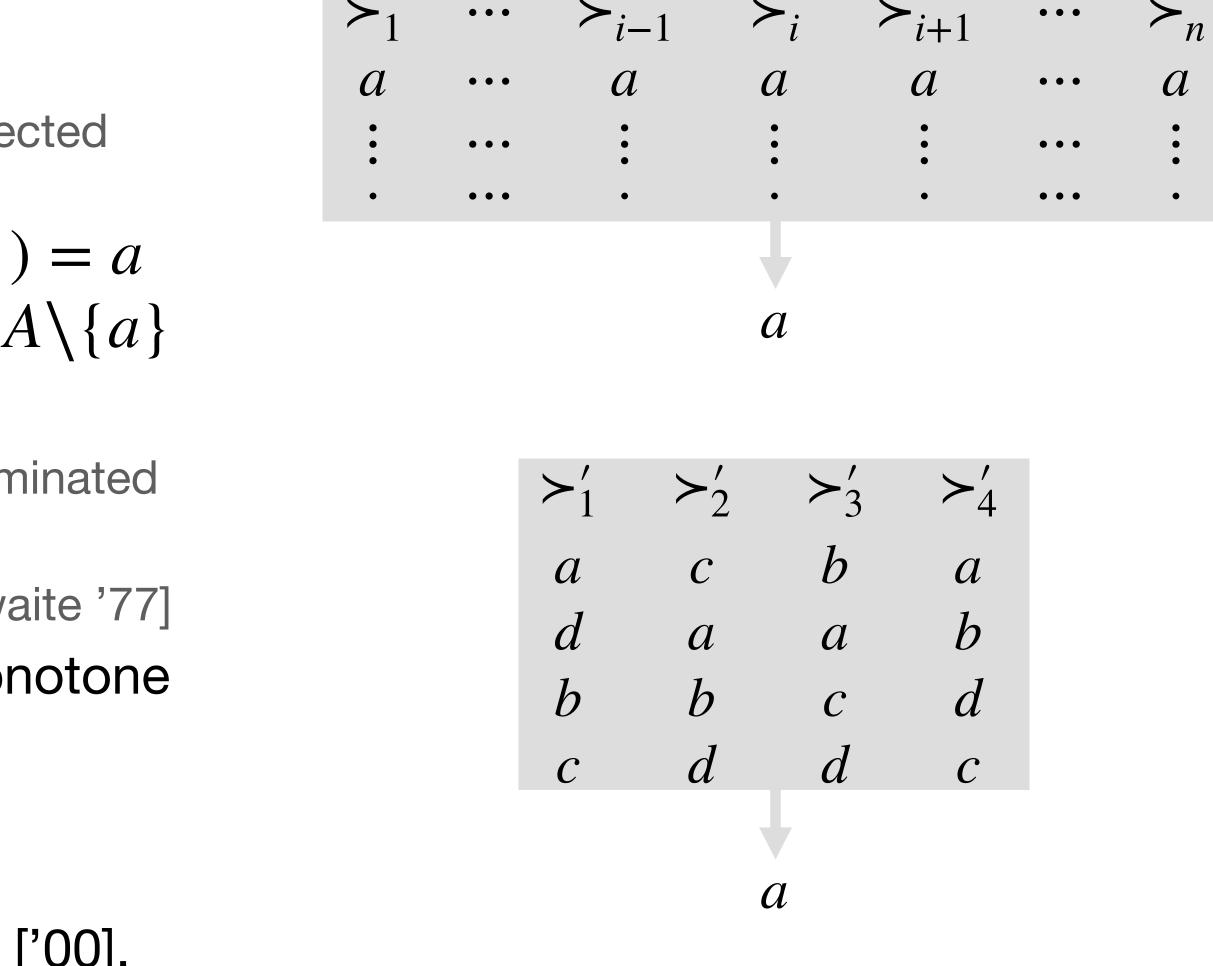
The Muller-Satterthwaite Theorem







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 - Theorem [Muller, Satterthwaite '77] Let $f: \mathscr{L}(A)^n \to A$ be a unanimous and monotone social choice function, where $|A| \geq 3$. Then, f is dictatorial.
- we give a proof of this theorem due to Reny ['00], so let f be as in the statement







• we consider two fixed alternatives $a, b \in A$

Pivotal Voter







- we consider two fixed alternatives $a, b \in A$
- we move from a profile where *a* is ranked highest by all voters to a profile where b is moving *b* one position at a time and one voter at a time







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\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
a	• • •	a	a	a	• • •	a
•	• • •	•	•	•	• • •	•
•	• • •	• • •	• • •	•	• • •	• • •
•	• • •	•	•	•	• • •	•
b	• • •	b	b	b	• • •	b
			a			







- we consider two fixed alternatives $a, b \in A$
- we move from a profile where *a* is ranked highest by all voters to a profile where b is moving *b* one position at a time and one voter at a time

\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
a	• • •	a	a	a	•••	a
•	• • •	•	•	•	• • •	•
•	• • •	• •	• •	• •	• • •	• • •
b	• • •	•	•	•	•••	•
•	• • •	b	b	b	• • •	b
			a			







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- we move from a profile where *a* is ranked highest by all voters to a profile where b is moving *b* one position at a time and one voter at a time

\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	•••	\succ_n
a	• • •	a	a	a	• • •	a
b	• • •	•	•	•	• • •	•
•	• • •	•	• •	• •	• • •	•
•	• • •	•	•	٠	• • •	•
•	• • •	b	b	b	•••	b
			a			







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\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
b	•••	a	a	a	• • •	a
a	•••	•	•	•	• • •	•
• •	•••	• • •	• • •	• • •	•••	•
•	• • •	•	•	•	•••	•
•	•••	b	b	b	• • •	b
			a			







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\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
b	• • •	a	a	a	• • •	a
a	•••	b	•	•	•••	•
•	• • •	•	• •	• •	•••	• •
•	• • •	•	•	•	• • •	•
•	• • •	•	b	b	• • •	b
			a			







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\succ_1	•••	\succ_{i-1}	\succ_i	\succ_{i+1}	•••	\succ_n
b	• • •	b	a	a	• • •	a
a	• • •	a	•	•	• • •	٠
•	• • •	• •	• •	• •	•••	• •
•	• • •	•	•	•	• • •	•
•	• • •	•	b	b	• • •	b
			a			







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\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
b	•••	b	a	a	•••	a
a	•••	a	b	•	•••	•
•	• • •	• •	• •	• •	• • •	• •
•	• • •	•	•	•	• • •	•
•	•••	•	•	b	•••	b
			a			







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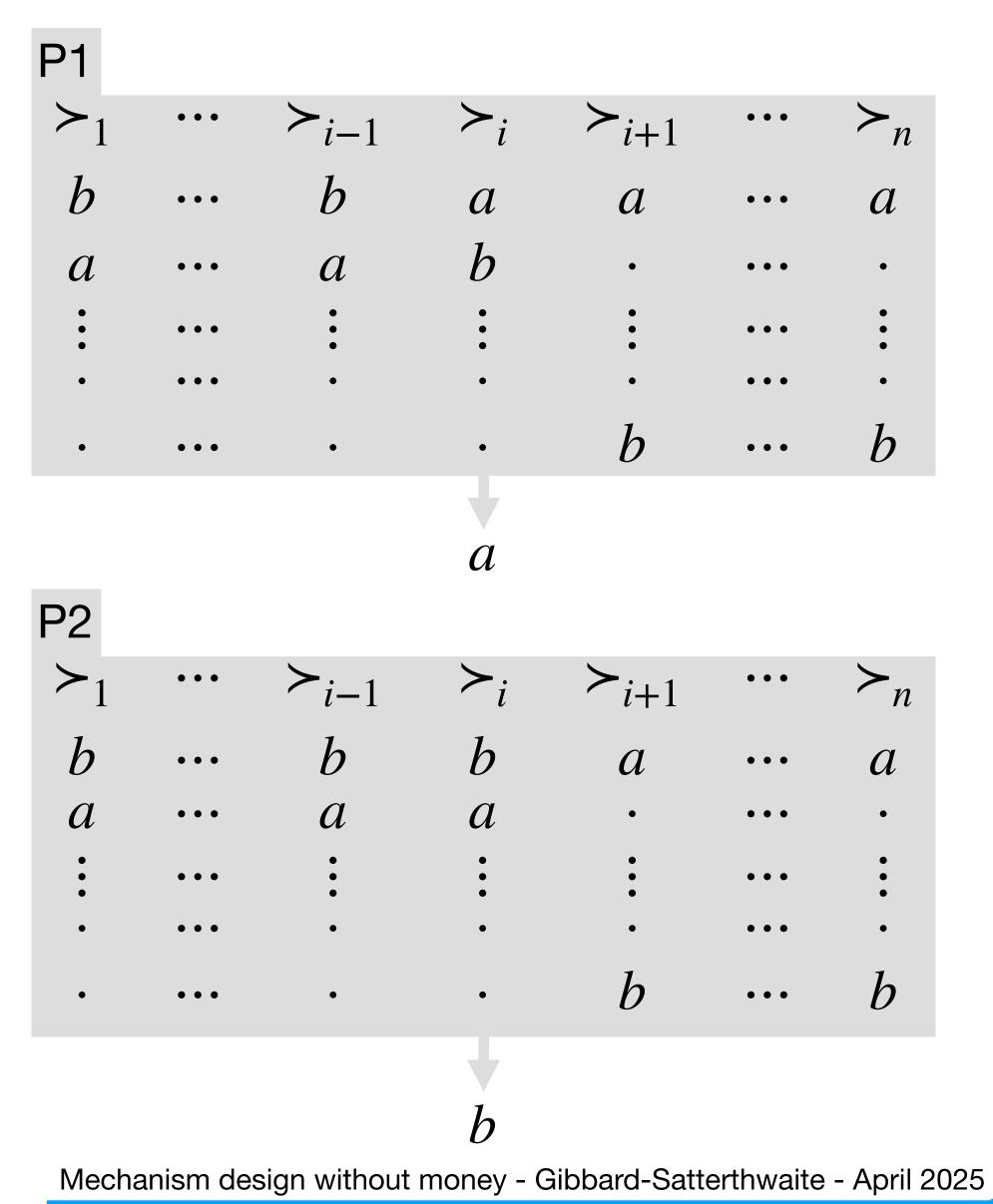
\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
b	•••	b	a	a	•••	a
a	•••	a	b	•	•••	٠
•	•••	• •	•	• •	•••	• •
•	•••	•	•	•	•••	•
•	•••	•	•	b	•••	b
			a			
•						
\succ_1	•••	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
\succ_1 b	•••		\succ_i b	\succ_{i+1}	•••	\succ_n
1	•••	<i>i</i> -1		$\iota + 1$	•••	11
b	•••	b	b	$\iota + 1$	•••	11
b	•••	b	b	$\iota + 1$	•••	π
b	•••	b	b	$\iota + 1$	•••	π
b	•••	b	b	$\iota + 1$	•••	11
b	•••	b	b	$\iota + 1$	•••	11







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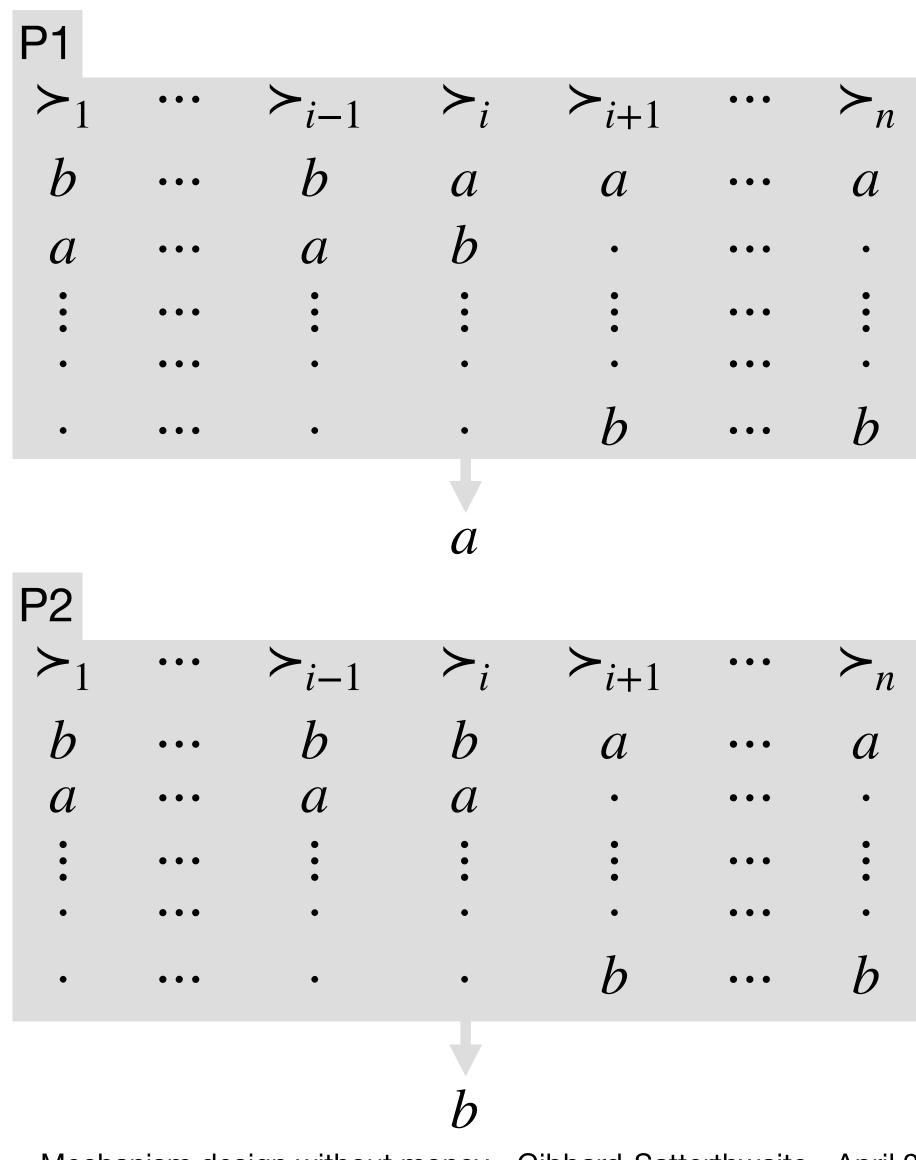








- we consider two fixed alternatives $a, b \in A$
- we move from a profile where *a* is ranked highest by all voters to a profile where b is moving b one position at a time and one voter at a time
- by unanimity, the outcome must change from a to b voter *i* changes profile P1 to P2

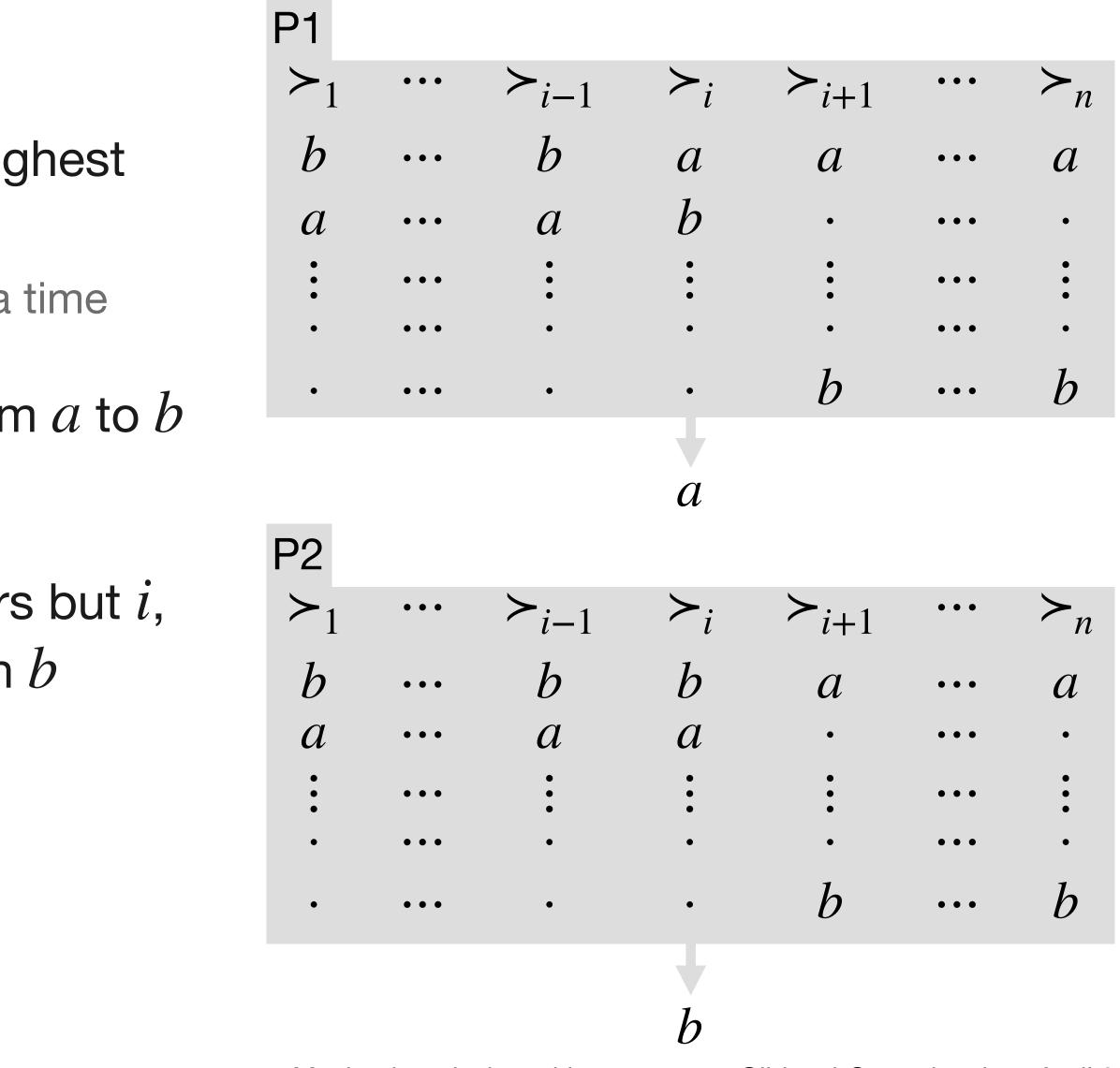








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- by unanimity, the outcome must change from a to b voter *i* changes profile P1 to P2
- start from P2 and move a below for all voters but i, without changing pairwise relationships with b

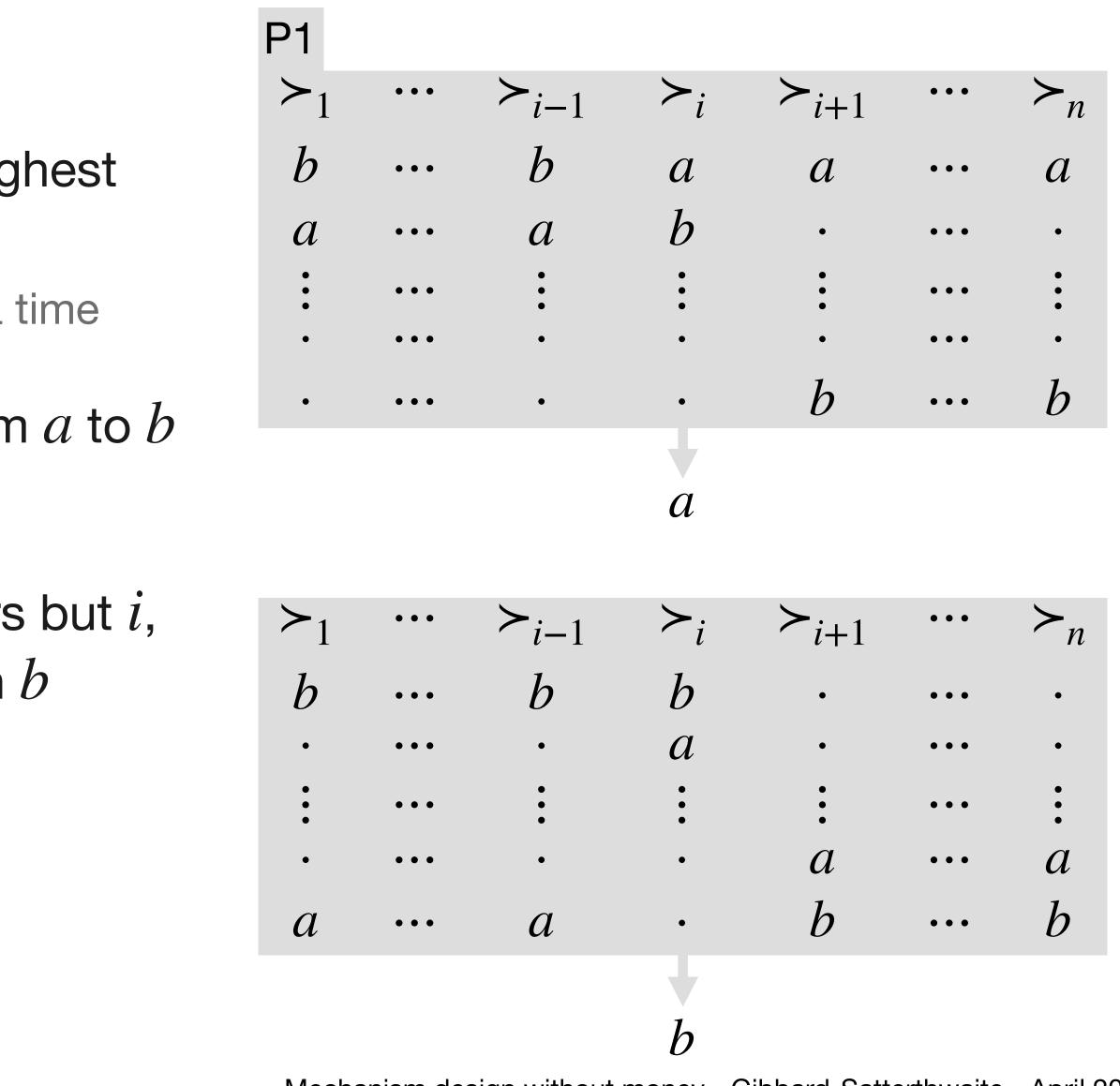








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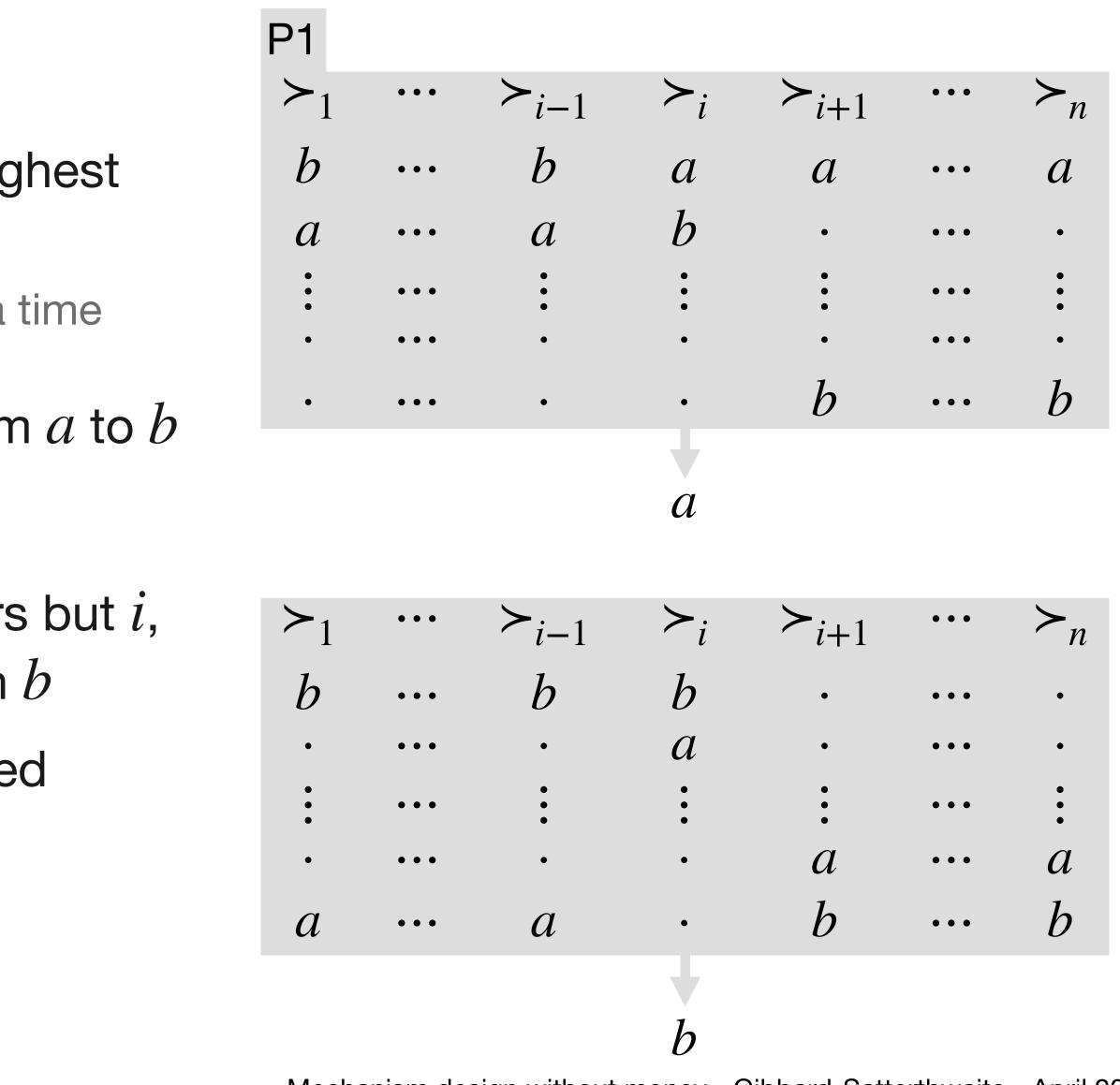








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- by unanimity, the outcome must change from a to bvoter *i* changes profile P1 to P2
- start from P2 and move a below for all voters but i, without changing pairwise relationships with b
 - by monotonicity, b must remain selected

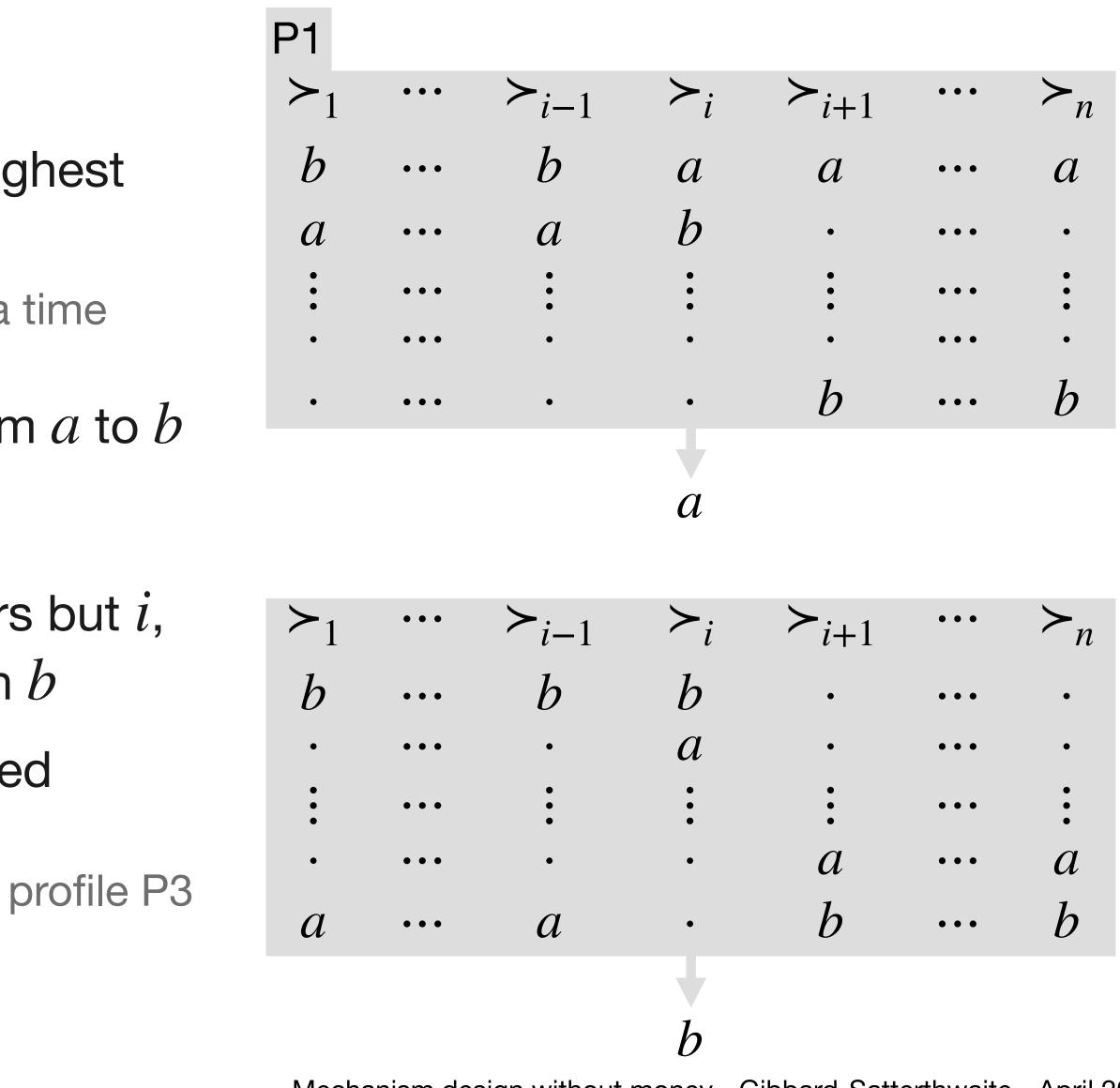








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- start from P2 and move a below for all voters but i, without changing pairwise relationships with b
 - by monotonicity, b must remain selected
- flip a and b in i's ranking and call the resulting profile P3

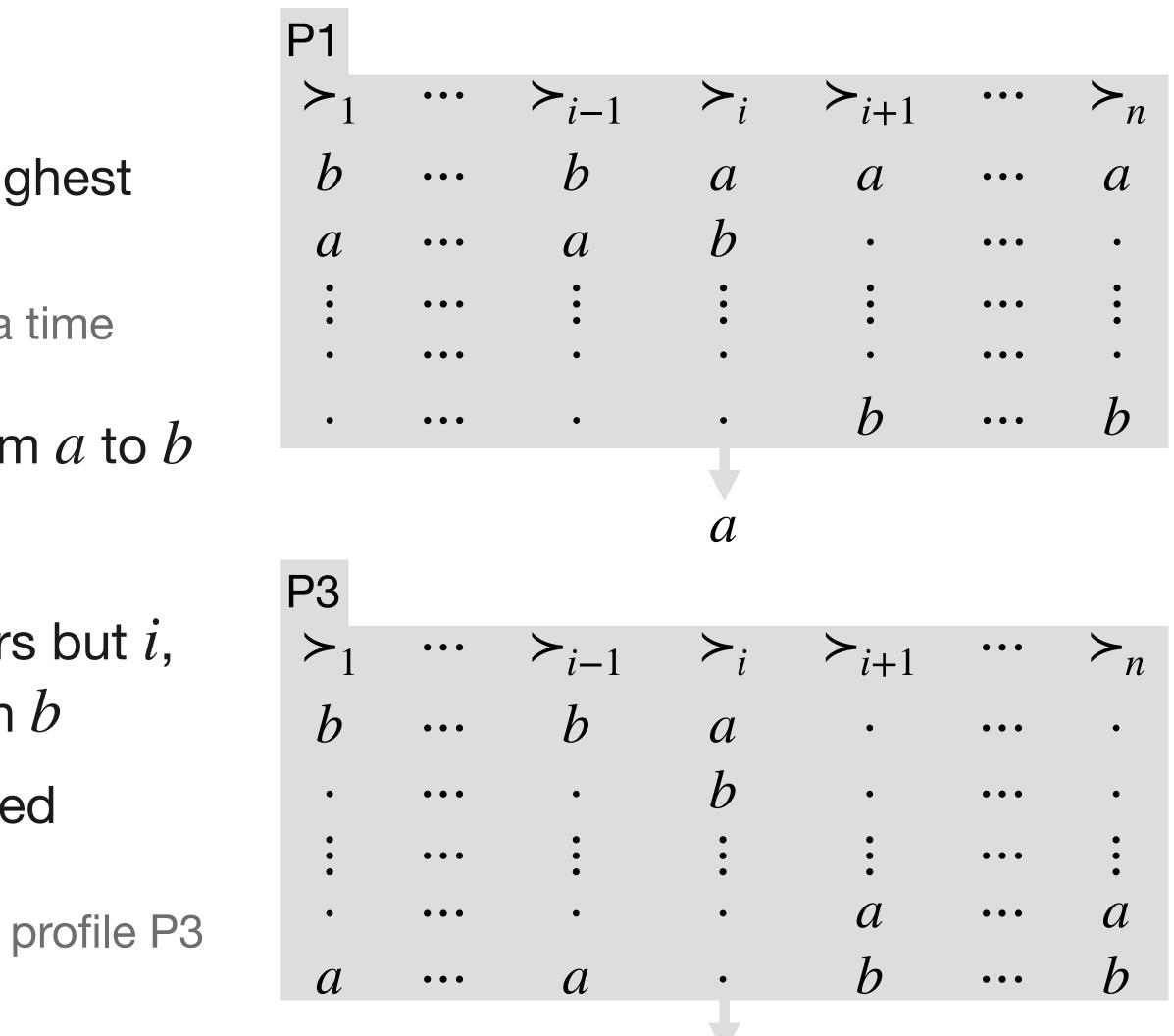








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 - by monotonicity, b must remain selected
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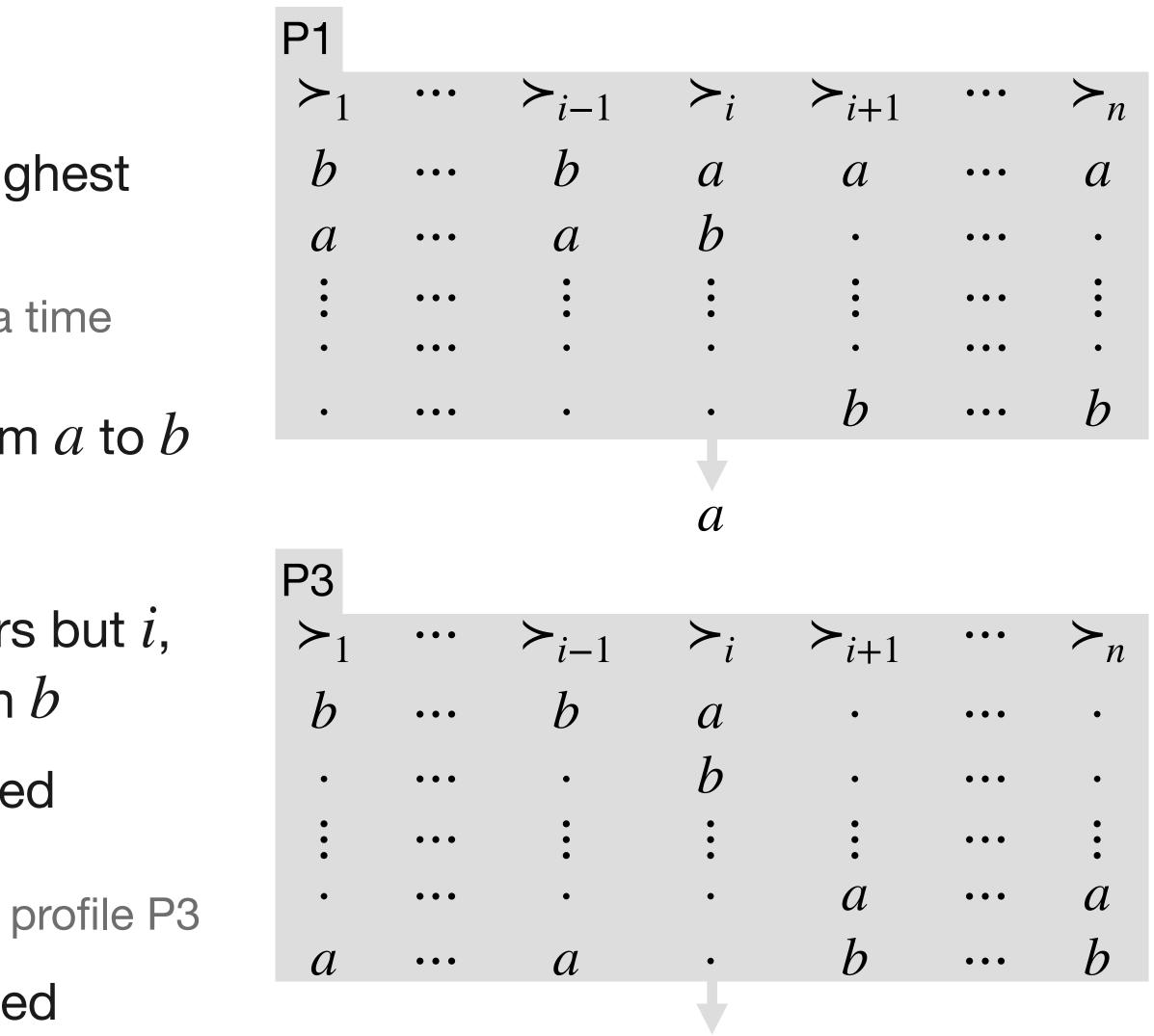








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- start from P2 and move a below for all voters but i, without changing pairwise relationships with b
 - by monotonicity, b must remain selected
- flip a and b in i's ranking and call the resulting profile P3 by monotonicity, either a or b is selected

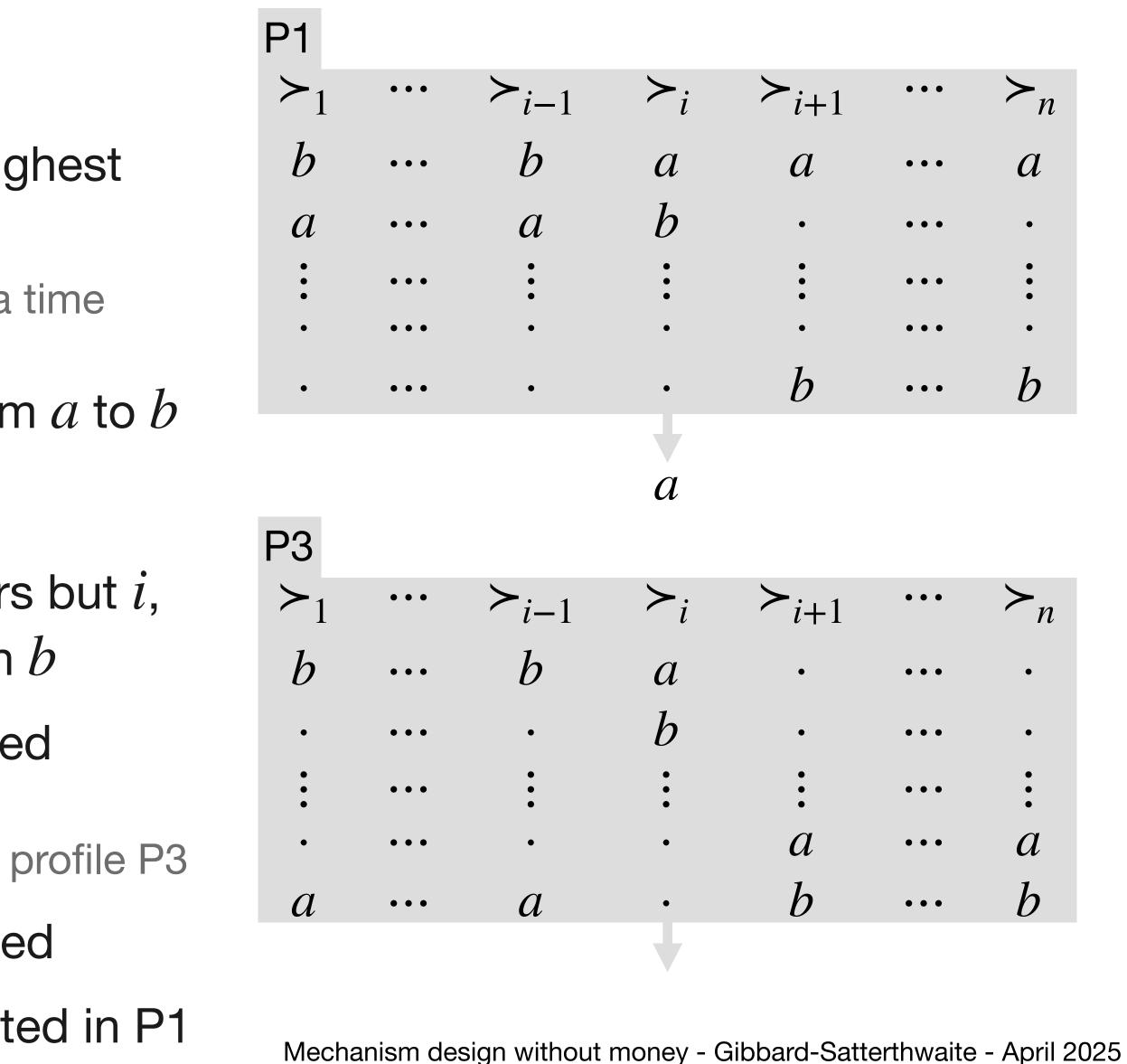








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- we move from a profile where *a* is ranked highest by all voters to a profile where b is moving b one position at a time and one voter at a time
- by unanimity, the outcome must change from a to bvoter *i* changes profile P1 to P2
- start from P2 and move a below for all voters but i, without changing pairwise relationships with b
 - by monotonicity, b must remain selected
- flip a and b in i's ranking and call the resulting profile P3
 - by monotonicity, either a or b is selected
 - selecting b would imply that b is selected in P1

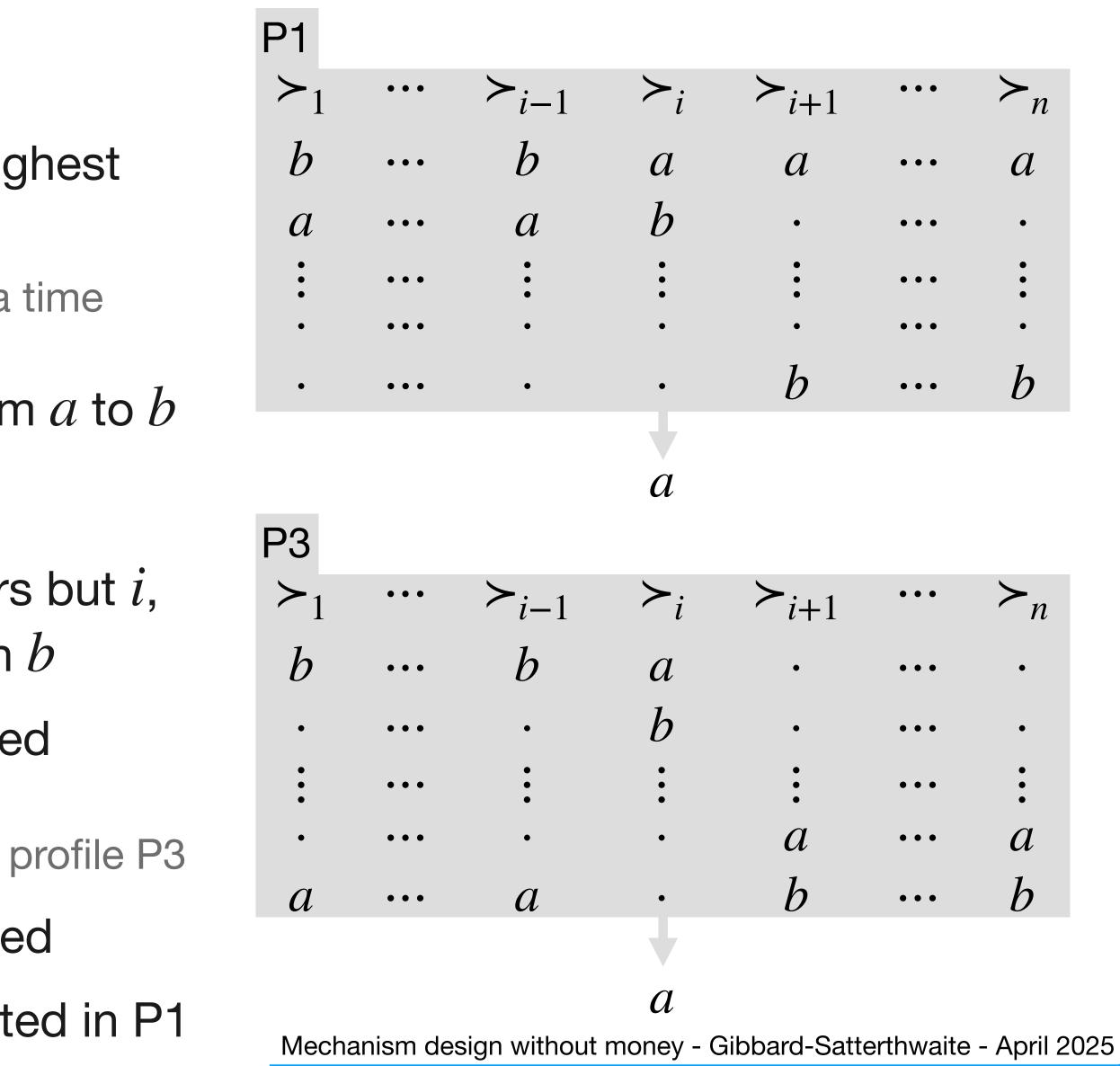








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 - by monotonicity, b must remain selected
- flip a and b in i's ranking and call the resulting profile P3
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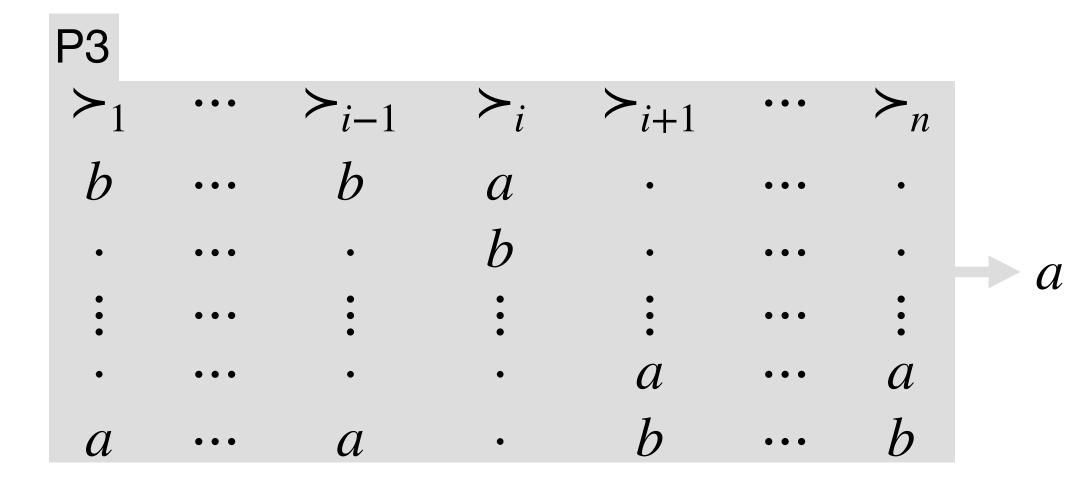


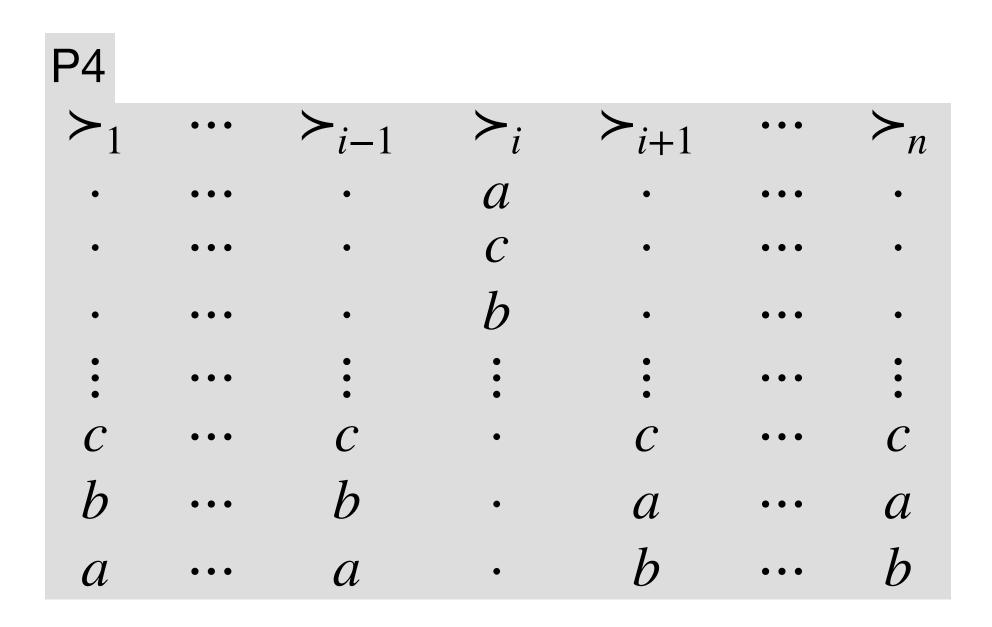






From Pivotal to Dictator







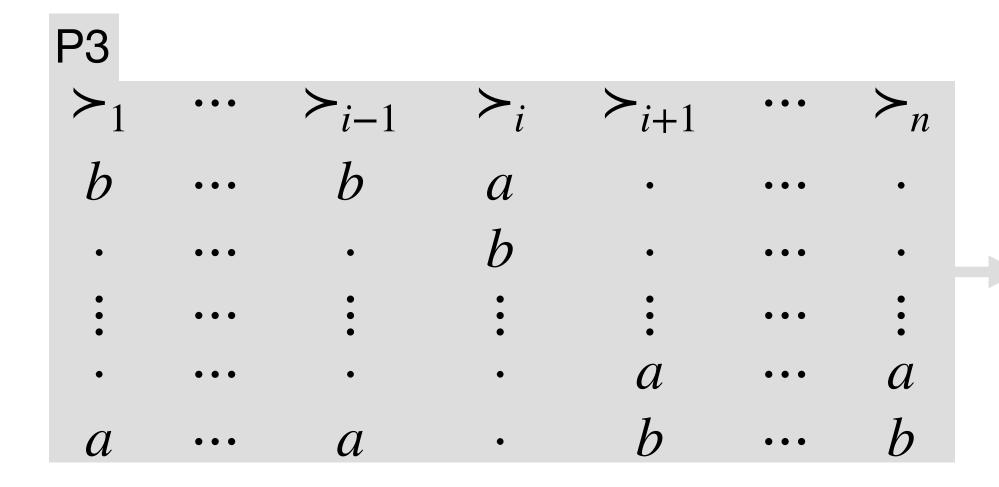


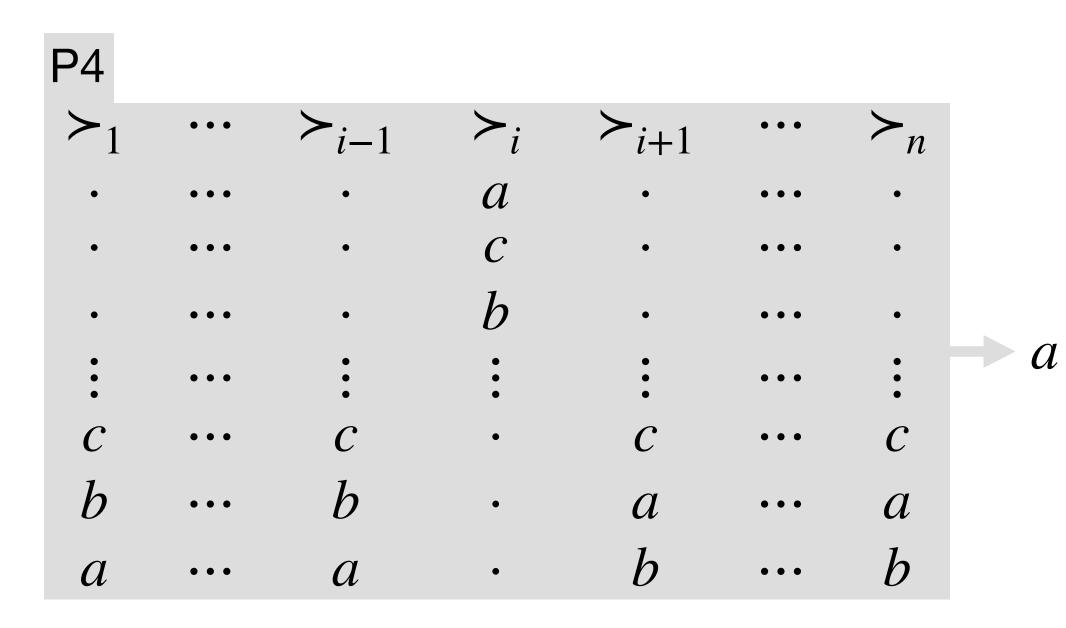




• by monotonicity, *a* remains selected

From Pivotal to Dictator









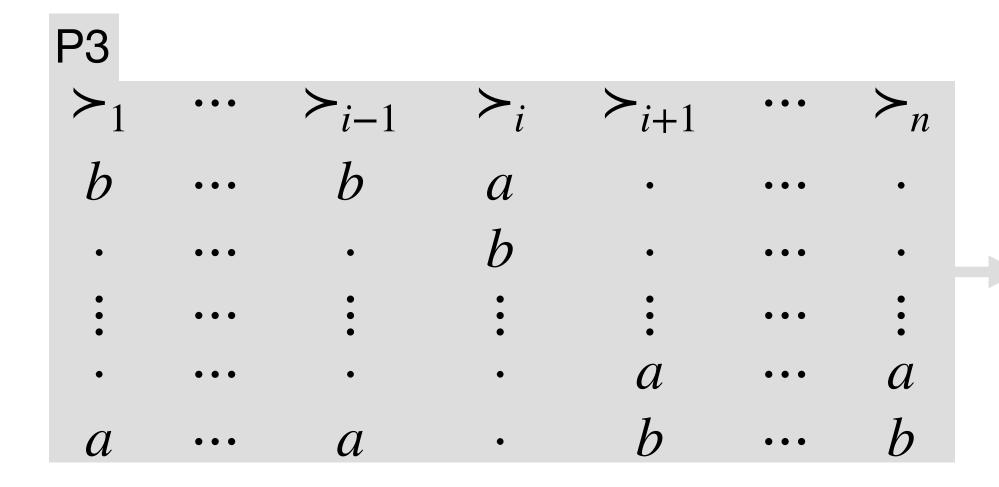


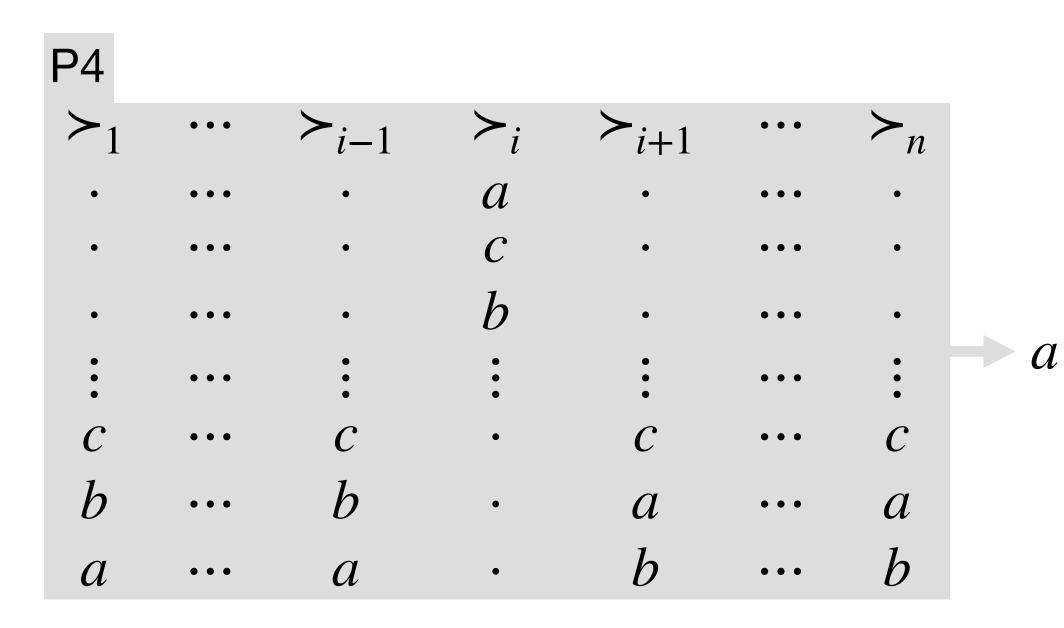




- by monotonicity, *a* remains selected
- Fip a and b in the ranking of all voters j > i

From Pivotal to Dictator









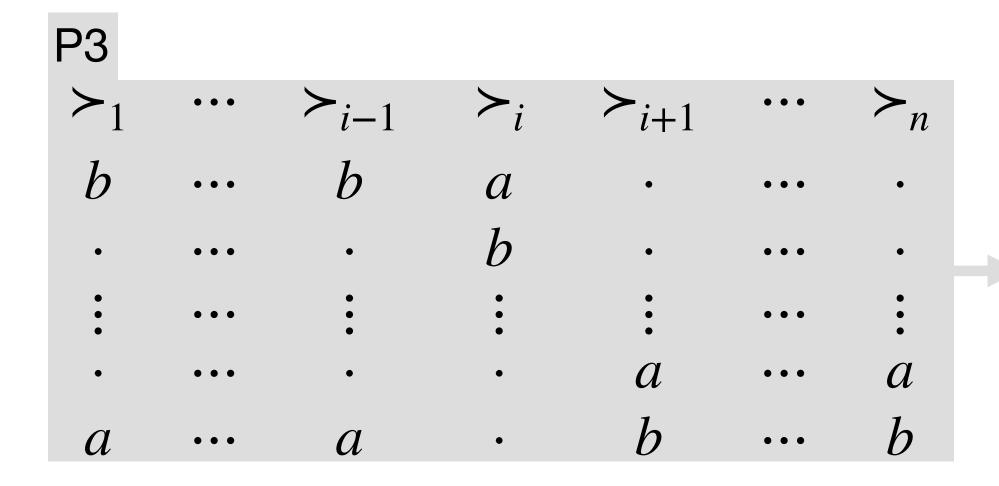






- by monotonicity, *a* remains selected
- Fip a and b in the ranking of all voters j > i

From Pivotal to Dictator



\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	•••	\succ_n
•	• • •	•	a	•	• • •	•
•	• • •	•	С	•	• • •	•
•	• • •	•	b	•	• • •	•
•	• • •	• •	•	• •	•••	•
С	• • •	С	•	С	•••	С
b	• • •	b	•	b	• • •	b
a	• • •	a	•	a	•••	a



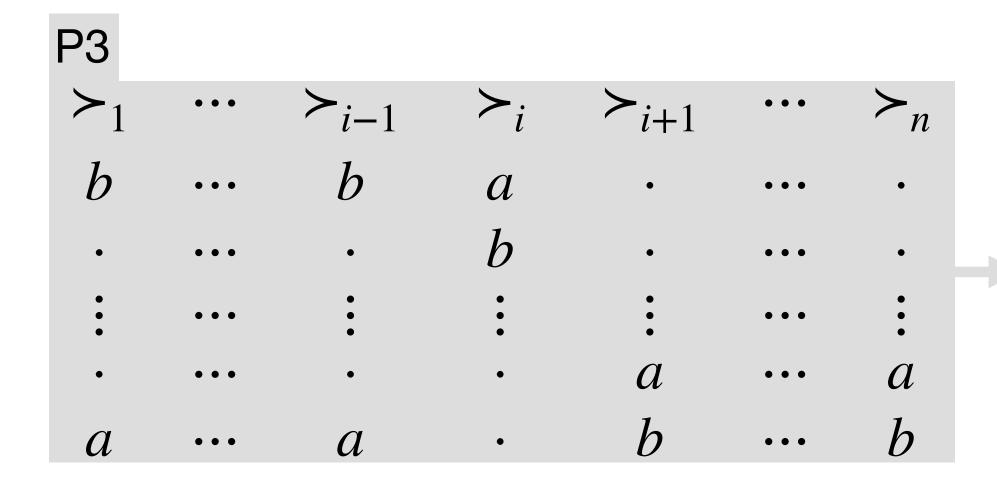






- by monotonicity, *a* remains selected
- Fip a and b in the ranking of all voters j > i
 - by monotonicity, either *a* or *b* is selected

From Pivotal to Dictator



\succ_1	• • •	\succ_{i-1}	\succ_i	\succ_{i+1}	• • •	\succ_n
•	• • •	•	a	٠	• • •	•
•	• • •	•	С	•	• • •	•
•	• • •	•	b	•	• • •	•
•	• • •	• •	• •	• •	• • •	•
С	• • •	С	•	С	•••	С
b	• • •	b	•	b	• • •	b
a	• • •	a	•	a	•••	a



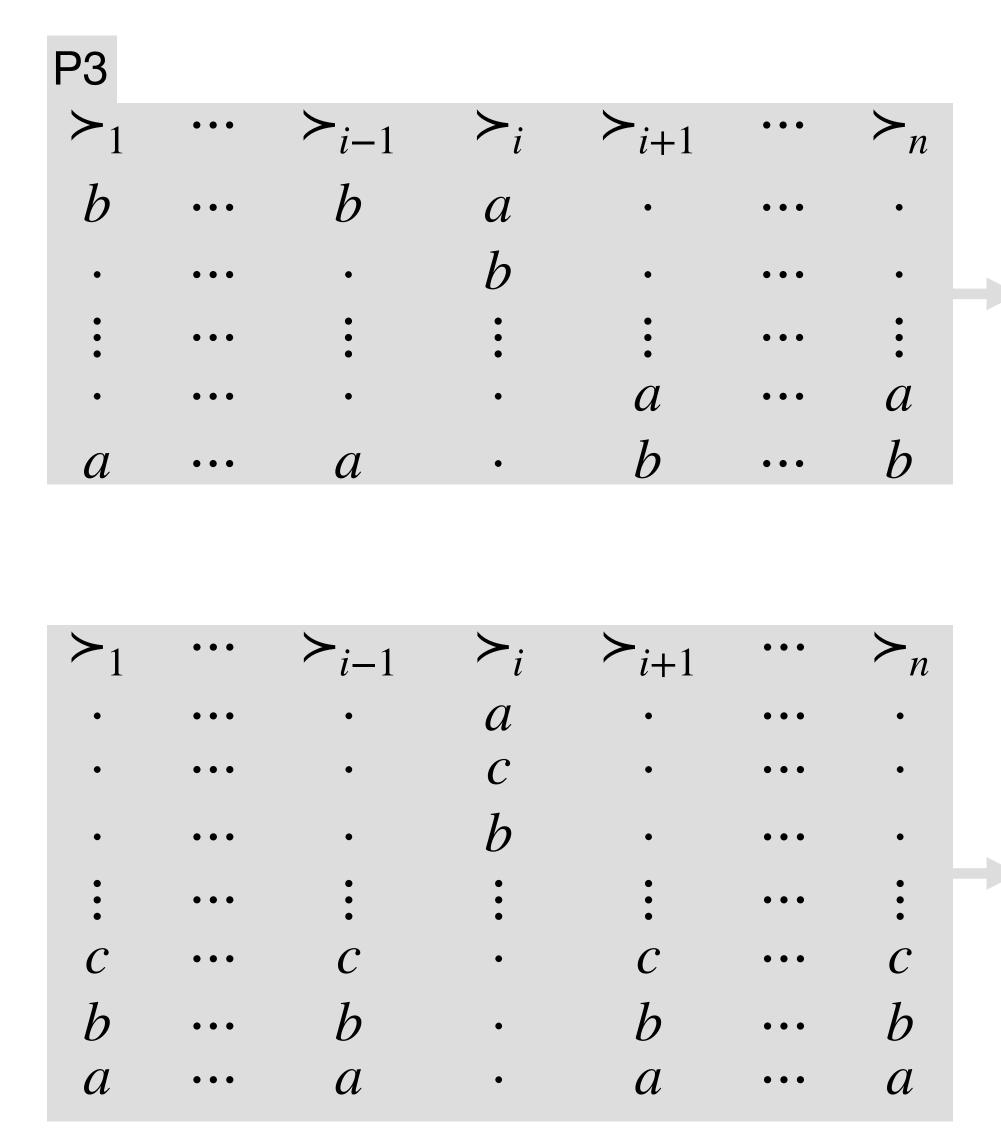






- by monotonicity, *a* remains selected
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From Pivotal to Dictator





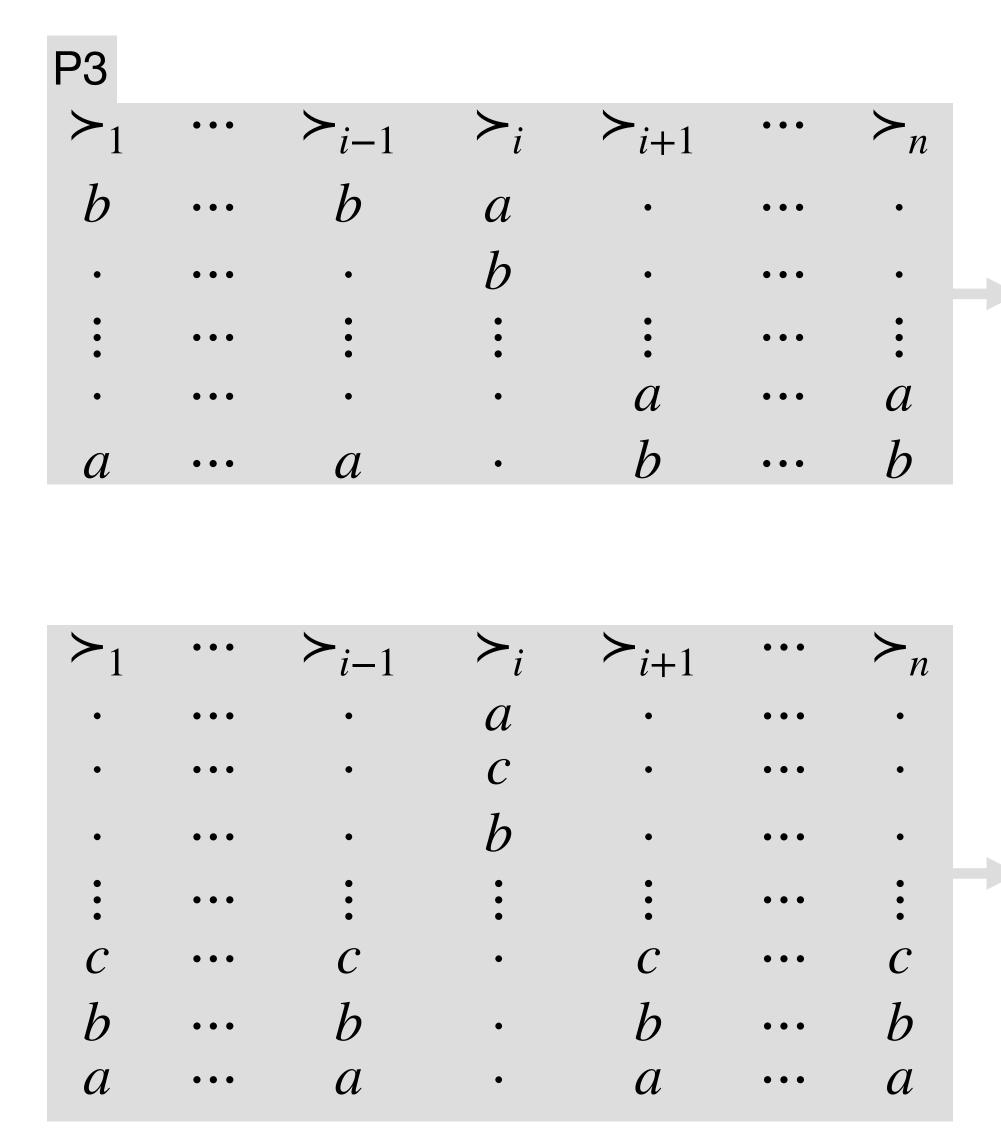






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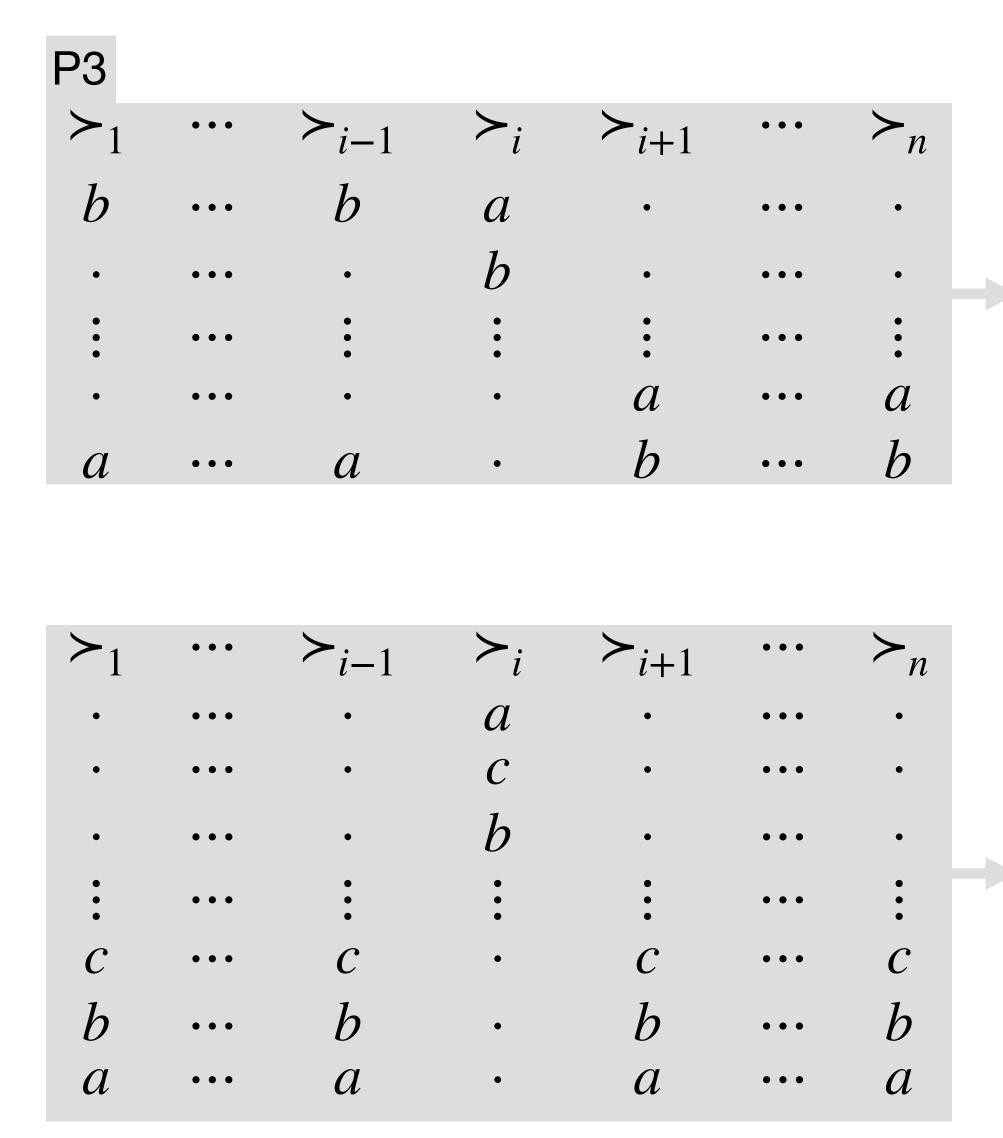






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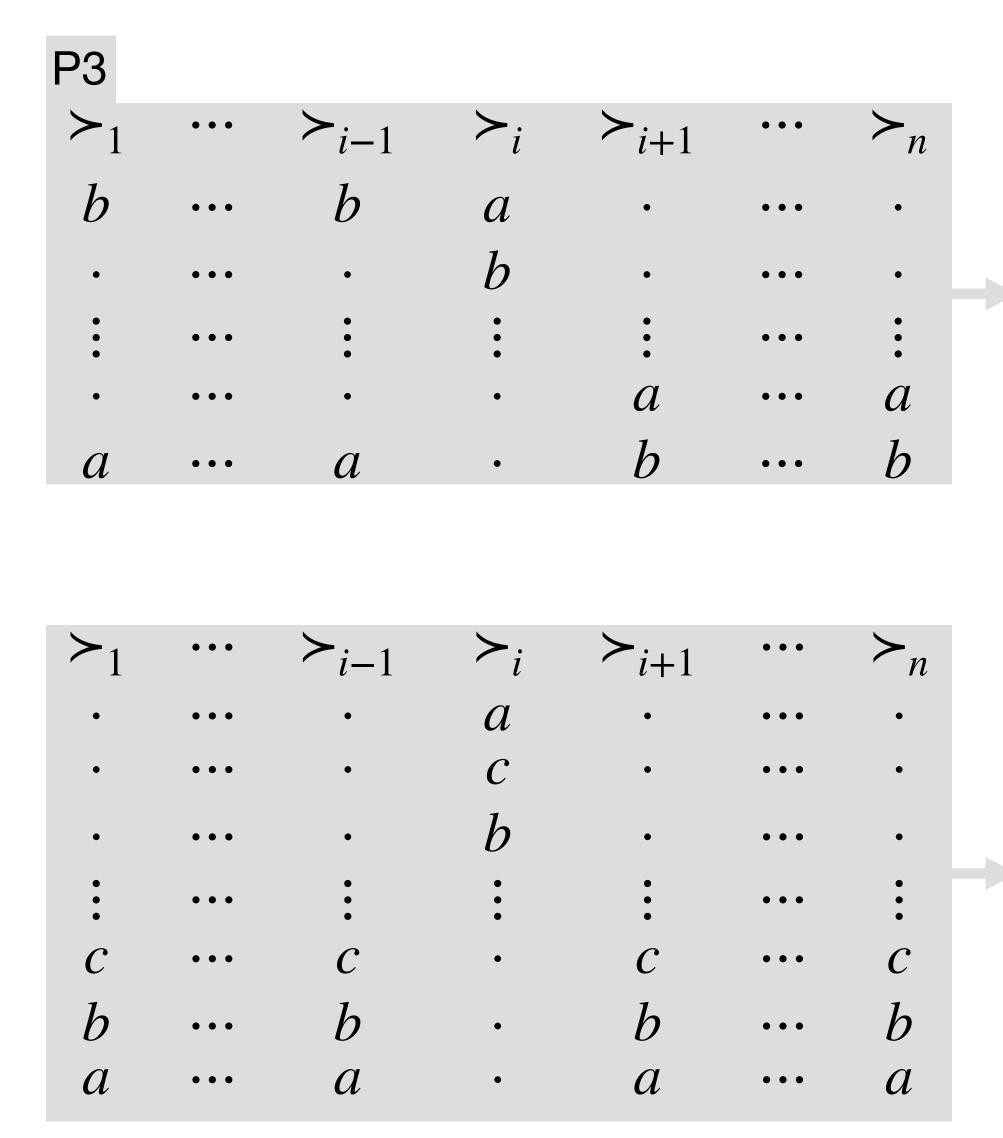






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 - there is a unique dictator for all alternatives

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Lemma If $f: \mathscr{L}(A)^n \to A$ is surjective and strategyproof, then it is unanimous and monotone.

From Muller-Satterthwaite to Gibbard-Satterthwaite

[Muller, Satterthwaite '77]







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by monotonicity, this holds when we move a to the top of all rankings and shuffle the rest







▶ player $i \in V$ has type $\theta_i \in \Theta_i$ and utility $u_i : A \times \Theta_i \to \mathbb{R}$

General Mechanisms







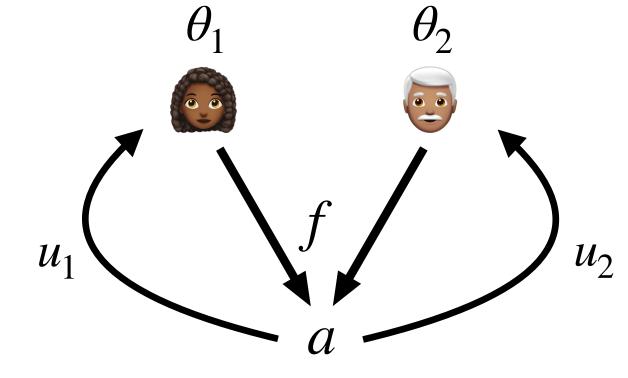
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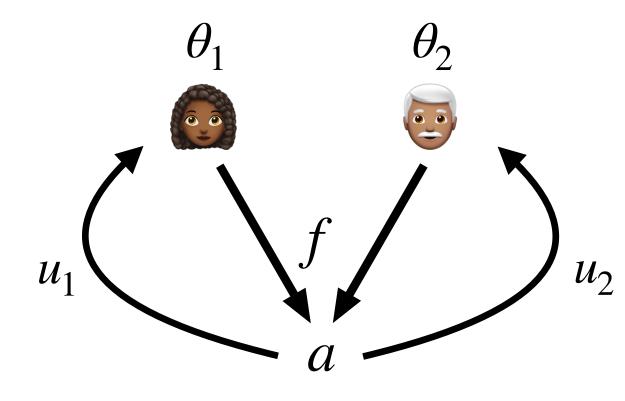






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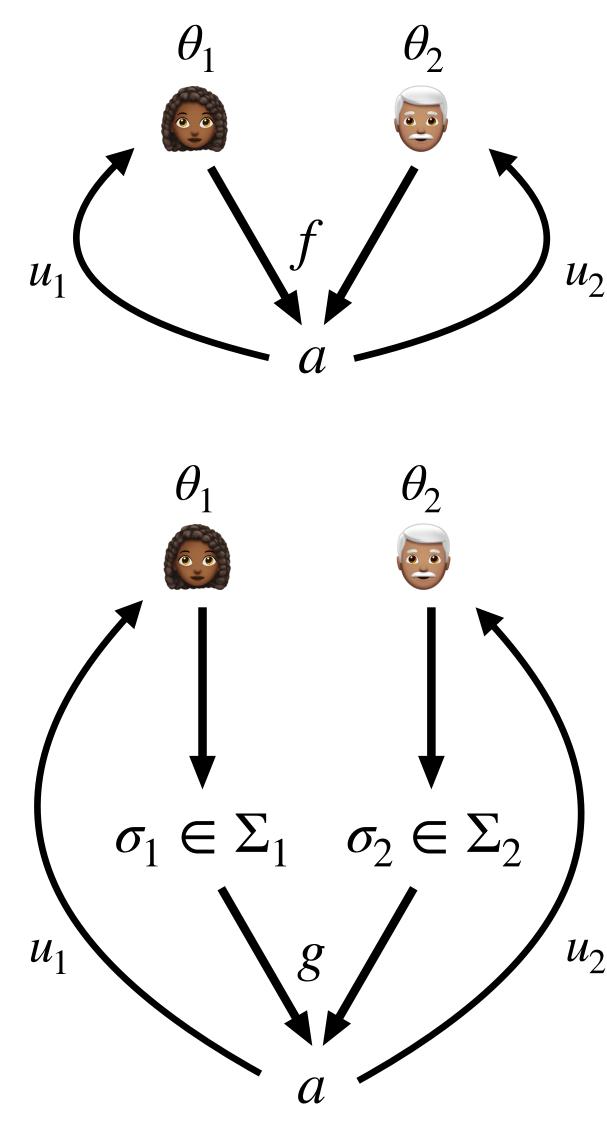




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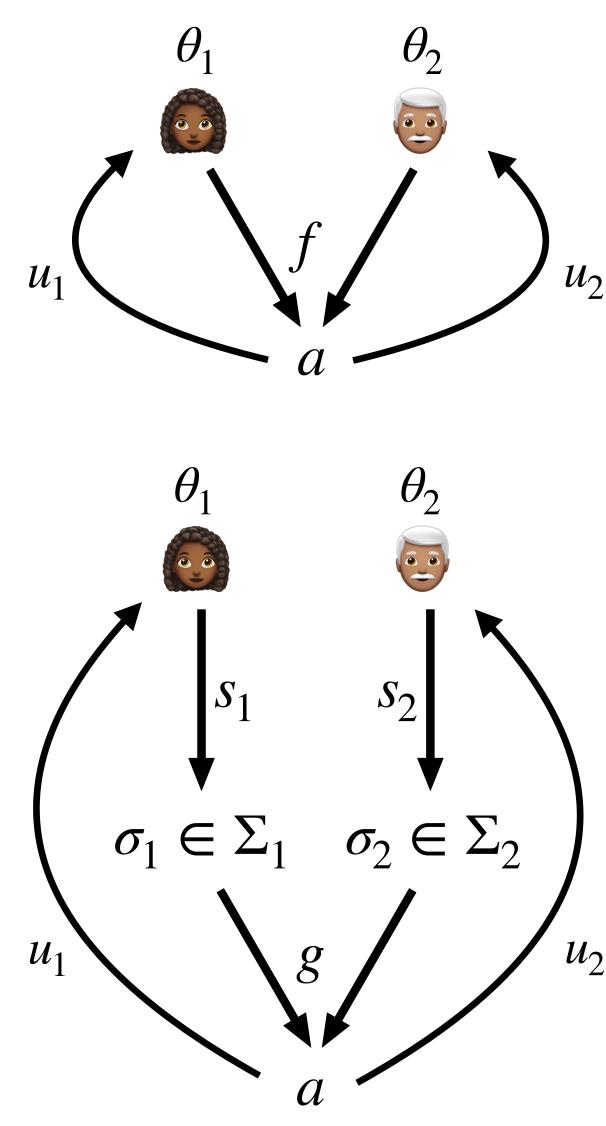




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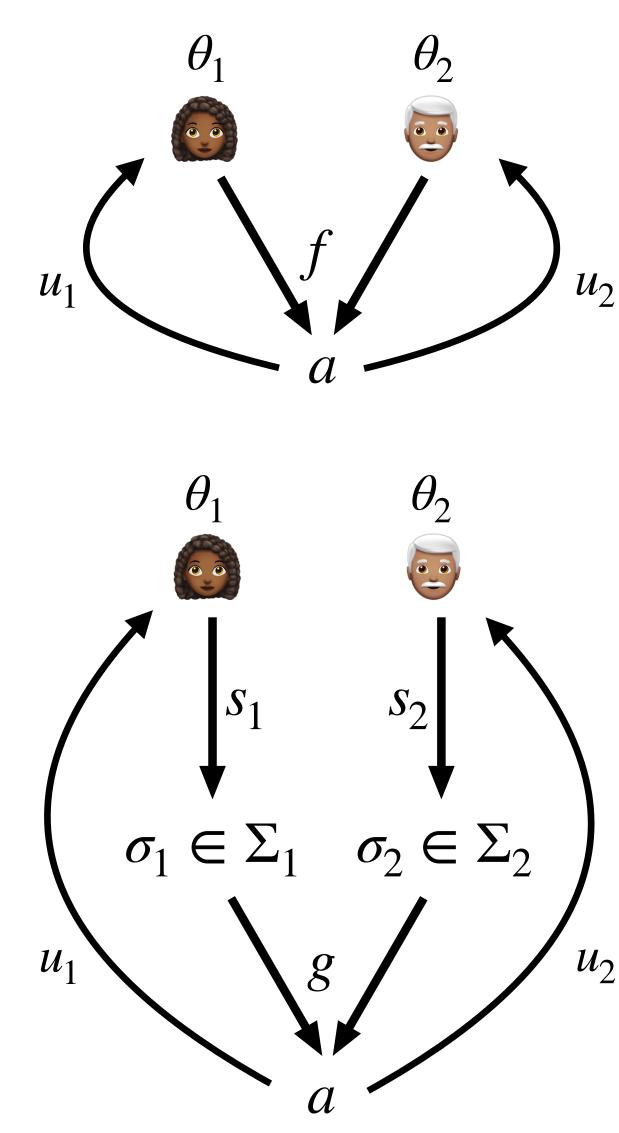








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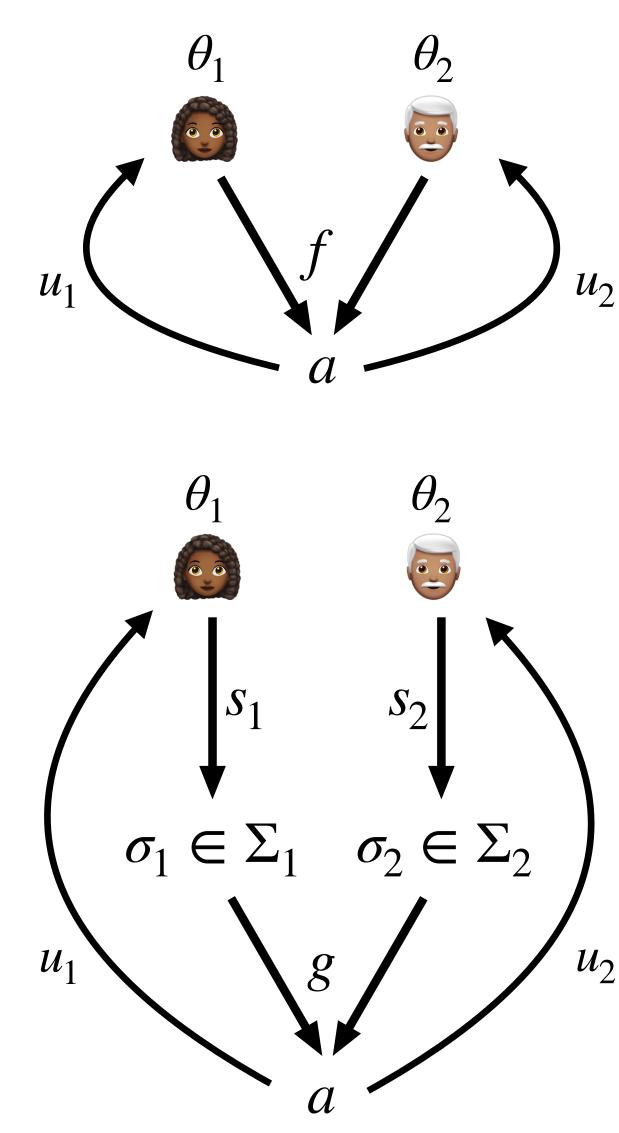






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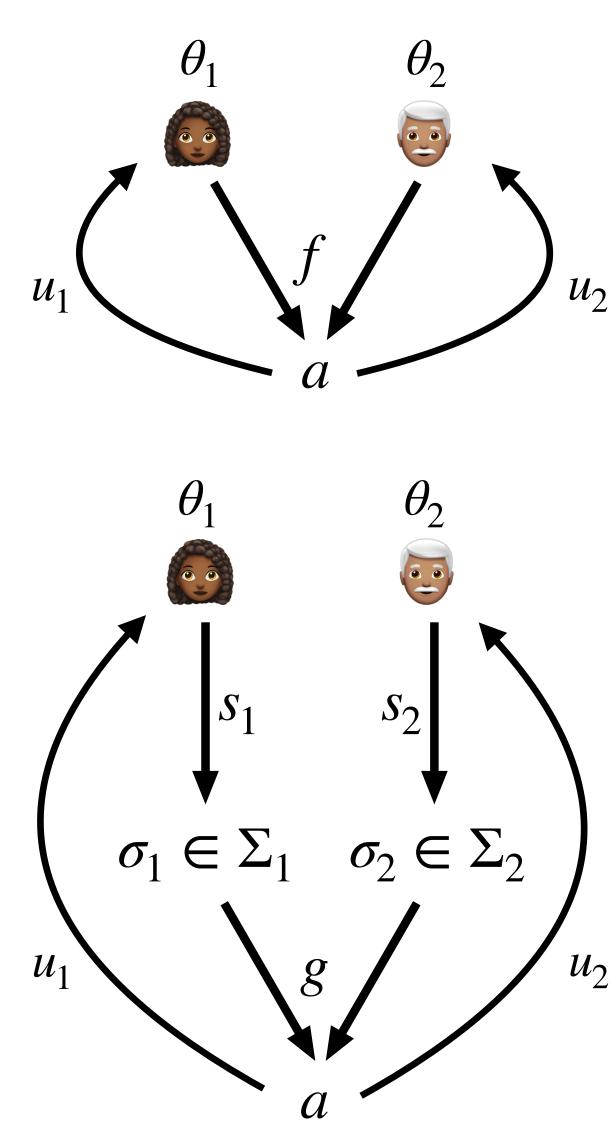








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 - $u_i(g(s_i(\theta_i), \sigma_i), \theta_i) \ge u_i(g(\sigma), \theta_i)$ for all $i \in V$ and $\sigma \in \Sigma$ S_i is a dominant strategy for every player i









• a mechanism is **direct** if $\Sigma = \Theta$

The Revelation Principle







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Revelation Principle: any implementable social choice function can be obtained from a strategyproof direct mechanism

Takeaways







- Revelation Principle: any implementable social choice function can be obtained from a strategyproof direct mechanism
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G-S: when $|A| \ge 3$, $\Theta = \mathscr{L}(A)^n$, and $u_i \colon A \times \Theta_i \to \mathbb{R}$ is s.t. $u_i(a, \theta_i) > u_i(b, \theta_i) \iff a \succ_i b$,







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The G-S theorem seems to quash any hope of designing incentive-compatible social-choice functions. The whole field of Mechanism Design attempts escaping from this impossibility result using various modifications in the model. [Nisan '07]







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