

Exercise 2: Flirting with Synchrony and Asynchrony

Task 1: Growing Balls

Denote by $B(v, r)$ the ball of radius r around v , i.e., $B(v, r) = \{u \in V : \text{dist}(u, v) \leq r\}$. Consider the following partitioning algorithm.

Algorithm 1 Cluster construction. $\rho \geq 2$ is a given parameter.

```
1: while there are unprocessed nodes do
2:   select an arbitrary unprocessed node  $v$ ;
3:    $r := 0$ ;
4:   while  $|B(v, r + 1)| > \rho|B(v, r)|$  do
5:      $r := r + 1$ 
6:   end while
7:   makeCluster( $B(v, r)$ )           // all nodes in  $B(v, r)$  are now processed
8:   remove all cluster nodes from the current graph
9: end while
10: select intercluster edges
```

- Show that Algorithm 1 constructs clusters of radius at most $\log_\rho n$.
- Show that Algorithm 1 produces at most ρn intercluster edges.
- For $k \in \{1, \dots, \lceil \log n \rceil\}$, determine an appropriate choice $\rho(k)$ and use it to prove Corollary 2.14!

Task 2: Showing Dijkstra, and Bellman & Ford the Ropes

- Show that if the asynchronous Bellman-Ford algorithm from the lecture is executed synchronously, it sends only $\mathcal{O}(|E|)$ messages.
- Use this to construct an asynchronous BFS tree construction algorithm of time complexity $\mathcal{O}(D)$ that uses $\mathcal{O}(|E|D)$ messages and terminates. You may assume that D is known here.
- Can you give an asynchronous Bellman-Ford-based algorithm that sends $\mathcal{O}(|E| + nD)$ messages and runs for $\mathcal{O}(D^2)$ rounds? (Hint: Either answer is feasible, provided it is backed up by appropriate reasoning!)

Task 3*: Liaison with Leslie Lamport

- Look up what Lamport causality, Lamport clocks, and Lamport vector clocks are.
- Contemplate their relation to synchronizers and what you've learned in the lecture.
- Discuss your findings in the exercise session!