



— Randomized Algorithms and Probabilistic Analysis of Algorithms, Exercise Sheet 6 —

<https://www.mpi-inf.mpg.de/departments/algorithms-complexity/teaching/winter22/random>

Total Points: 100

Due: Wednesday, **January 25, 2023**

You are allowed to collaborate on the exercise sheets. Justify your answers. Cite all external sources that you use (books, websites, research papers, etc.). You need to collect at least 50% of all points on exercise sheets to be admitted to the exam. Please hand in your solutions before the lecture on the day of the deadline.

— **Exercise 1** — 25 points —

Write  $G = (V, E)$  for a connected, undirected (non-bipartite) graph with maximum degree  $\Delta := \max_{v \in V} |N(v)|$ . Prove that for any distribution  $\pi : V \rightarrow [0, 1]$ ,  $\sum_{v \in V} \pi_v = 1$ , and any  $M \geq \delta$ , the following Markov chain on  $V$  has  $\pi$  as its stationary distribution:

$$P_{u,v} := \begin{cases} (1/M) \min\{1, \pi_v/\pi_u\}, & \text{if } u \neq v \text{ and } \{u, v\} \in E, \\ 0, & \text{if } u \neq v \text{ and } \{u, v\} \notin E, \\ 1 - \sum_{w \neq u} P_{u,w}, & \text{if } u = v. \end{cases}$$

— **Exercise 2** — 10 + 5 + 10 points —

One may shuffle a deck of  $n$  cards by repeatedly picking a card uniformly and independently at random and putting the chosen card on top of the deck. Thinking of all permutations of the cards as the states, we can view this process as a Markov chain, which we call the card-shuffling chain.

- a) Show that the card-shuffling chain is finite, irreducible, and aperiodic and compute its stationary distribution.

Consider the following coupling  $Z_t = (X_t, Y_t)$  for the card-shuffling chain. To obtain  $X_{t+1}$  from  $X_t$ , we proceed as before, that is, we pick a card  $C$  uniformly and independently at random from among the  $n$  cards and move it to the top. To obtain  $Y_{t+1}$  from  $Y_t$ , we pick the card with value  $C$  and place it at the top of the chain.

- b) Prove that  $Z_t$  is indeed a valid coupling for the card-shuffling chain.  
c) Using the coupling  $Z_t$ , prove that the card-shuffling chain is rapidly mixing, that is, its mixing time is polynomial in  $\log(1/\varepsilon)$  and  $n$ . (Hint: Recall the Coupon Collector's Problem.)

— **Exercise 3** — 10 + 10 points —

- a) Prove that any sequence  $Z_0, Z_1, \dots, Z_n$  that is a martingale with respect to  $X_0, X_1, \dots, X_n$  is also a martingale with respect to itself.  
b) For a sequence of random variables  $X_0, X_1, \dots$ , write  $S_i := \sum_{j=0}^i X_j$ . Show that if  $S_0, S_1, \dots$  is a martingale with respect to  $X_0, X_1, \dots$ , then for all  $i \neq j$ , we have  $\mathbb{E}[X_i X_j] = 0$ .

— **Exercise 4** — 30 points —

Prove the following generalization of the Azuma-Hoeffding inequality seen in class.

**Theorem.** For any martingale  $X_0, \dots, X_n$ , constants  $d_k$  and random variables  $B_k$  (that may be functions in  $X_0, \dots, X_{k-1}$ ) that satisfy

$$B_k \leq X_k - X_{k-1} \leq B_k + d_k,$$

we have for all  $t \geq 0$  and any  $\varepsilon > 0$  that

$$\Pr[|X_t - X_0| \geq \varepsilon] \leq 2e^{-2\varepsilon^2 / (\sum_{k=1}^t d_k^2)}.$$