# UNIVERSITÄT DES SAARLANDES 

FR 6.2 - Informatik
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Lecture "Automated Reasoning"
(Winter Term 2016/2017)
Midterm Examination

Name: $\qquad$

Student Number: $\qquad$
Some notes:

- Things to do at the beginning:

Put your student card and identity card (or passport) on the table.
Switch off mobile phones.
Whenever you use a new sheet of paper (including scratch paper), first write your name and student number on it.

- Things to do at the end:

Mark every problem that you have solved in the table below.
Stay at your seat and wait until a supervisor staples and takes your examination text.
Note: Sheets that are accidentally taken out of the lecture room are invalid.

| Problem | 1 | 2 a | 2 b | 2 c | 3 | 4 | 5 | 6 a | 6 b | 6 c | 7 | $\Sigma$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Answered? |  |  |  |  |  |  |  |  |  |  |  |  |
| Points |  |  |  |  |  |  |  |  |  |  |  |  |

Refute the following clause set by superposition where you may apply the reduction rules Condensation, and Subsumption Resolution. Use the ordering $P_{4} \succ P_{3} \succ P_{2} \succ P_{1}$. You may also make use of a selection function.
$1 \quad P_{2} \vee P_{4}$
$2 \quad P_{1} \vee P_{4}$
$4 \quad P_{1} \vee \neg P_{3}$
$5 \quad P_{4} \vee \neg P_{3}$
$3 \quad \neg P_{2} \vee P_{1}$
$7 \quad \neg P_{1} \vee P_{3} \quad 8 \quad \neg P_{4} \vee P_{3}$
$6 \neg P_{3} \vee P_{2}$
$10 \quad \neg P_{2} \vee \neg P_{3} \quad 11 \quad \neg P_{4} \vee \neg P_{2} \vee P_{3}$

Problem 2 (Superposition Model Building)
Consider the below clause set with atom ordering $P_{5} \succ P_{4} \succ P_{3} \succ P_{2} \succ P_{1}$.

| 1 | $P_{1} \vee P_{2} \vee P_{2}$ | 2 | $\neg P_{1} \vee \neg P_{2}$ | 3 | $\neg P_{2} \vee \neg P_{3}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 4 | $P_{1} \vee P_{3}$ | 5 | $P_{4} \vee P_{5} \vee P_{1}$ | 6 | $\neg P_{4} \vee P_{1}$ |
| 7 | $\neg P_{4} \vee P_{2}$ | 8 | $\neg P_{5} \vee P_{2}$ | 9 | $\neg P_{5} \vee \neg P_{3}$ |

(a) Compute $N_{\mathcal{I}}$.
(b) Determine the minimal false clause in $N_{\mathcal{I}}$. Perform the respective superposition inference on the clause. Add the derived clause to $N$ resulting in $N^{\prime}$ and compute $N_{\mathcal{I}}^{\prime}$.
(c) Determine the minimal false clause in $N_{\mathcal{I}}^{\prime}$. Perform the respective superposition inference on the clause. Add the derived clause to $N^{\prime}$ resulting in $N^{\prime \prime}$ and compute $N_{\mathcal{I}}^{\prime \prime}$.

Check via CDCL whether the below clause set is satisfiable.

| 1 | $P 11 \vee P 12$ | 2 | $P 21 \vee P 22$ | 3 | $P 31 \vee P 32$ |
| ---: | :--- | ---: | :--- | ---: | :--- |
| 4 | $P 41 \vee P 42$ | 5 | $\neg P 11 \vee P 42$ | 6 | $\neg P 42 \vee P 11$ |
| 7 | $\neg P 11 \vee \neg P 21$ | 8 | $\neg P 11 \vee \neg P 31$ | 9 | $\neg P 31 \vee \neg P 41$ |
| 10 | $\neg P 12 \vee \neg P 22$ | 11 | $\neg P 32 \vee \neg P 42$ | 12 | $\neg P 12 \vee \neg P 32$ |

Problem 4 (CNF)
Transform the formula

$$
(P \vee((Q \leftrightarrow \top) \wedge \neg R)) \vee(P \leftrightarrow(Q \leftrightarrow \perp))
$$

into CNF using $\Rightarrow$ ACNF .

Prove that the formula

$$
((\neg P \vee \neg R) \rightarrow Q) \rightarrow(\neg Q \rightarrow(P \wedge R))
$$

is valid using tableau. You may use a tree representation of the tableau.

Which of the following statements are true or false? Provide a proof or a counter example.

1. If $N_{\mathcal{I}} \models N$ then $N$ is saturated up to redundancy.
2. If all clauses in $N$ have at most one positive literal and the CDCL rule Propagate is not applicable to the state $(\epsilon ; N ; \emptyset ; 0 ; \top)$ then $N$ is satisfiable.
3. If all clauses in $N$ have at most one positive literal and there is no clause in $N$ having only negative literals then $N_{\mathcal{I}} \models N$.

Consider a reasonable CDCL run

$$
(\epsilon ; N ; \emptyset ; 0 ; \top) \Rightarrow{ }_{\mathrm{CDCL}}^{*}\left(L_{1} \ldots L_{n} ; N ; \emptyset ; k ; D\right)
$$

where the last applied rule was Conflict and hence $D \notin\{\top, \perp\}$. Consider the atom ordering atom $\left(L_{1}\right) \prec \operatorname{atom}\left(L_{2}\right) \prec \ldots \prec \operatorname{atom}\left(L_{n}\right)$. Prove that any of the subsequent CDCL Resolve steps until backtracking is a Superposition Left inference with respect to $\prec$, where clauses are always condensed.

